

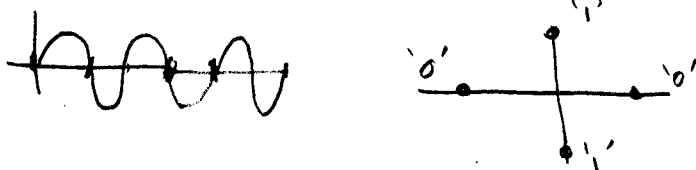
6.5 Detecting Signals With Memory - the Viterbi Algorithm

6.5.1

7.4.2

- When we detect signals with memory, we have two choices:

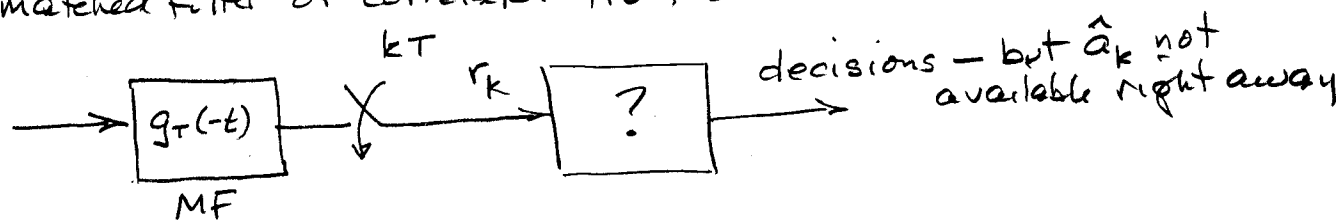
- Ignore the memory and take a hit on BER vs SNR. We did this for AMI & bipolar in previous section. Similarly, treat this $h = \frac{1}{2}$ 2FSK as biorthogonal with duplicates



- Account for the memory. More complex, but significant performance improvement. That's what we learn in the present section, using the celebrated Viterbi Algorithm.

- When we account for memory, we use the facts that some sequences cannot occur, and the remainder have, in principle, a new pdf. Most of the applications have all allowable sequences equiprobable. Since we are now detecting entire sequences, the approach is MLSE (max likelihood sequence estimation).

- The detector in white noise is still constructed with a matched filter or correlator front end:



- In each symbol interval, we still have conditional probabilities $p(r|s_1)$, $p(r|s_2)$. For example, binary antipodal:

$$s_1 = \sqrt{E_b}, \quad s_2 = -\sqrt{E_b}$$

$$p(r|s_1) = \frac{1}{\sqrt{2\pi}\sigma_n} e^{-\frac{1}{2}\left(\frac{r-\sqrt{E_b}}{\sigma_n}\right)^2} \quad p(r|s_2) = \frac{1}{\sqrt{2\pi}\sigma_n} e^{-\frac{1}{2}\left(\frac{r+\sqrt{E_b}}{\sigma_n}\right)^2}$$

and for AMI

$$s_1 = \sqrt{2E_b} \quad s_2 = 0 \quad s_3 = -\sqrt{2E_b}$$

$$p(r|s_1) = \frac{1}{\sqrt{2\pi}\sigma_n} \exp\left(-\frac{1}{2}\left(\frac{r-\sqrt{2E_b}}{\sigma_n}\right)^2\right) \quad p(r|s_2) = \frac{1}{\sqrt{2\pi}\sigma_n} \exp\left(-\frac{1}{2}\left(\frac{r}{\sigma_n}\right)^2\right)$$

$$p(r|s_3) = \frac{1}{\sqrt{2\pi}\sigma_n} \exp\left(-\frac{1}{2}\left(\frac{r+\sqrt{2E_b}}{\sigma_n}\right)^2\right)$$

- Sequence estimation. Define $\underline{r}_k = (r_1, r_2, \dots, r_k)$ and $\underline{\sigma}_k = (\sigma_1, \sigma_2, \dots, \sigma_k)$ where σ_i is the state at time i . Since the state sequence determines the data sequence, and the state sequences are equiprobable, we do this:

$$\arg \max_{\underline{\sigma}_k} p(\underline{r}_k | \underline{\sigma}_k) \quad \text{for length-} k \text{ sequences.}$$

This is ML. Suppose there are 2 states. It looks as though we have to run through all 2^k sequences $\underline{\sigma}$, compute $p(\underline{r}_k | \underline{\sigma}_k)$ for each one and find the largest. Is this realistic?

- A recursive formulation is our first step out of the swamp of exponential growth.

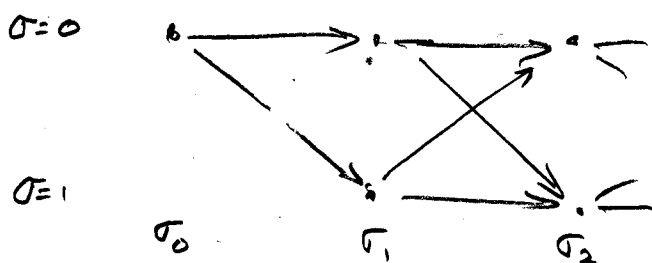
$$\begin{aligned}
 p(\Gamma_K | \Sigma_K) &= p(\Gamma_K, \Sigma_{K-1} | \Sigma_K) = p(\Gamma_K | \Sigma_{K-1}, \Sigma_K) p(\Sigma_{K-1} | \Sigma_K) \text{ why?} \\
 &= p(\Gamma_K | \Sigma_K) p(\Sigma_{K-1} | \Sigma_{K-1}) \text{ why?} \\
 &= p(\Gamma_K | \sigma_K, \sigma_{K-1}) p(\Sigma_{K-1} | \Sigma_{K-1}) \text{ why?} \\
 &= p(\Gamma_K | \sigma_K, \sigma_{K-1}) p(\Gamma_{K-1} | \sigma_{K-1}, \sigma_{K-2}) \cdots p(\Gamma_2 | \sigma_2, \sigma_1) p(\Gamma_1 | \sigma_1, \sigma_0)
 \end{aligned}$$

and the logarithm converts it to an additive form

$$\ln(p(\Gamma_K | \Sigma_K)) = \ln(p(\Gamma_K | \sigma_K, \sigma_{K-1})) + \ln(p(\Gamma_{K-1} | \sigma_{K-1}, \sigma_{K-2})) + \cdots + \ln(p(\Gamma_1 | \sigma_1, \sigma_0))$$

from previous page $\ln(p(\Gamma_i | \sigma_i, \sigma_{i-1})) \sim -(\Gamma_i - \underbrace{s(\sigma_i, \sigma_{i-1})}_{\text{a function that gives } s \text{ from the state pair}})^2$

$$\text{so } \underset{\Sigma_K}{\text{argmax}} \ln(p(\Gamma_K | \Sigma_K)) = \underset{\Sigma_K}{\text{argmin}} \sum_{i=1}^K (\Gamma_i - s(\sigma_i, \sigma_{i-1}))^2$$



We minimize the

Euclidean metric.

branch metric

$$\mu(\Gamma, \sigma_{\text{now}}, \sigma_{\text{then}}) = (\Gamma - s(\sigma_{\text{now}}, \sigma_{\text{then}}))^2$$

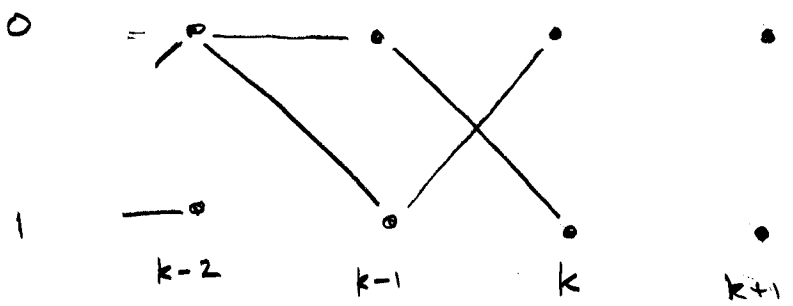
cumulative metric

$$M(\Sigma) = \sum_{i=1}^K \mu(\Gamma_i, \sigma_i, \sigma_{i-1})$$

Not much progress so far.

The "dynamic programming" recursion is the next step. Suppose we know at time k the best paths terminating in each state $0, 1$ (denote $\tilde{\sigma}_k(0), \tilde{\sigma}_k(1)$) and the corresponding metrics $M(\tilde{\sigma}_k(0)), M(\tilde{\sigma}_k(1))$.

$\tilde{\sigma}_k(0), \tilde{\sigma}_k(1)$ are the "survivor" paths at time k



To extend to time $k+1$, note that the best path terminating in $\sigma_{k+1} = 0$ has either $\tilde{\sigma}_k(0)$ or $\tilde{\sigma}_k(1)$ as a prefix. Same with best path terminating in $\sigma_{k+1} = 1$. That is,

$$\tilde{\sigma}_{k+1}(0) = [\tilde{\sigma}_k(0), 0] \text{ or } [\tilde{\sigma}_k(1), 0]$$

$$\tilde{\sigma}_{k+1}(1) = [\tilde{\sigma}_k(0), 1] \text{ or } [\tilde{\sigma}_k(1), 1]$$

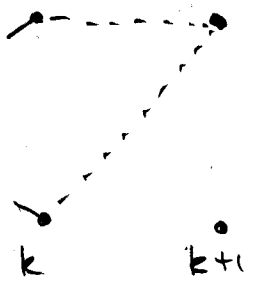
For state 0:

Choose prefix that minimizes accumulated metric

$$M(\tilde{\sigma}_k(0)) + \mu(r_{k+1} | 0, 0)$$

or

$$M(\tilde{\sigma}_k(1)) + \mu(r_{k+1} | 1, 0)$$



The smaller value becomes $M(\tilde{\sigma}_{k+1}(0))$.

For state 1:

Choose prefix that minimizes accumulated metric

$$M(\tilde{\sigma}_k(0)) + \mu(r_{k+1}, 0, 1)$$

or

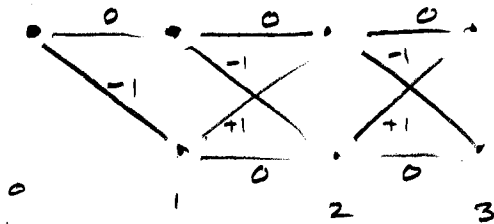
$$M(\tilde{\sigma}_k(0)) + \mu(r_{k+1}, 1, 1)$$



The smaller value becomes $M(\tilde{\sigma}_{k+1}(0))$

Example

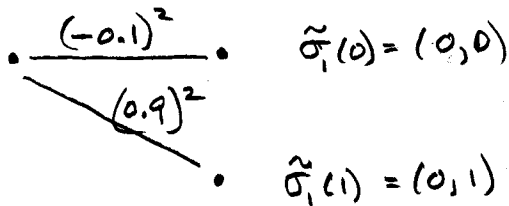
AMI, initialized in state 0, $M(\sigma_0=0) = 0$
normalized to amplitude 1, so $E_b = \frac{1}{2}$



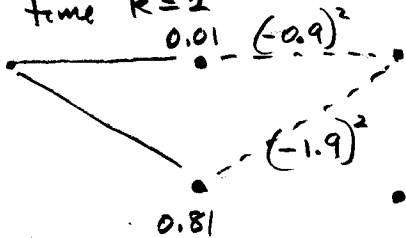
We receive

$$r_1 = -0.1, r_2 = -0.9, r_3 = 1.1$$

- time $k=1$

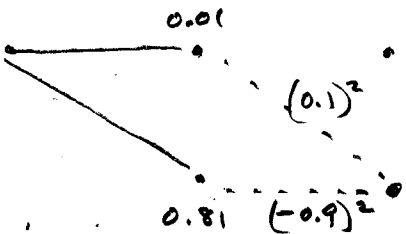


- time $k=2$



smaller of 0.82 and 4.42 makes 0,0 the better prefix.

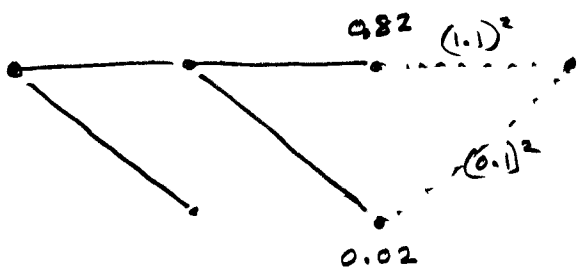
$$\tilde{\sigma}_2(0) = (0, 0, 0) \quad M(\tilde{\sigma}_2(0)) = 0.82$$



smaller of 0.02 and 1.62 makes 0,0 the better prefix

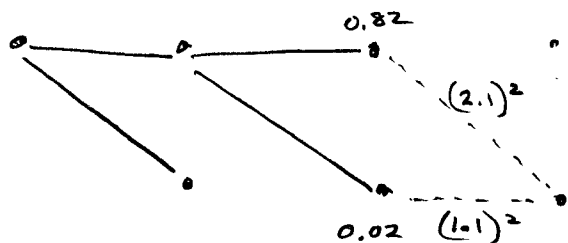
$$\tilde{\sigma}_2(1) = (0, 0, 1) \quad M(\tilde{\sigma}_2(1)) = 0.02$$

- time $k=3$



smaller of 2.03 and 0.03 makes 0,0,1
the better prefix

$$\hat{\sigma}_3(0) = (0,0,1,0) \quad M(\hat{\sigma}_3(0)) = 0.03$$

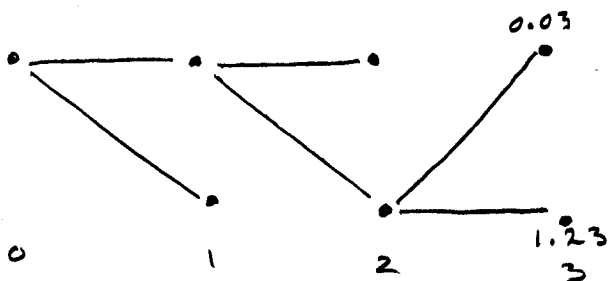


smaller of 5.23 and 1.23 makes 0,0,1
the better prefix

$$\hat{\sigma}_3(1) = (0,0,1,1) \quad M(\hat{\sigma}_3(1)) = 1.23$$

All possible future paths have

0,0,1 as a prefix. Release them as decisions. Termed a "merge" of paths



This particular trellis and constellation can be simplified further, but this example should be enough to establish:

- the method
- the fact that computation is proportional to k not 2^k (!)