

FINAL PROJECT PRESENTATION Spring 2001

## ATM Traffic Control Based on Cell Loss Priority and Performance Analysis

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# Introduction

- Overview of ATM (Asynchronous Transfer Mode)
  - Ultimate solution of BISDN (Broadband ISDN)
  - Connection-oriented service
  - 53-byte fixed packet called cell
  - Transmission rate : 25Mbps ~ 2.5Gbps
  - Five service categories
    - : CBR, rt-VBR, nrt-VBR, ABR, UBR
  - QoS parameters
    - : Cell loss, source traffic rate, delay, delay jitter etc.

### ATM cell structure



GFC : Generic Flow Control VCI : Virtual Channel Identifier VPI : Virtual Path Identifier CLP : Cell Loss Priority PT : Payload Type HEC : Header Error Control

- CLR 1 bit set to 0 for high priority cell or 1 for low priority

cell, which applies to several buffer priority schemes

- CLP based queuing schemes
  - Priority queuing
  - Appropriate in cases where WAN (Wide Area Network) links are congested from time to time, but unnecessary otherwise because of extra processes required and performance degradations for low priority traffic.
  - Schemes
  - Push-out, partial buffer sharing, buffer separation, hybrid and so on



- Push-out
- Diagram



- If the buffer is full and a high priority cell (CLP = 0) arrives, the last low priority cell, which already resides in the buffer, will be pushed out and lost. All incoming low priority cells arriving during congestion will be discarded.
- Otherwise, the queue operates based on FCFS.

Flowchart



- Partial buffer sharing
  - Diagram



- Once the threshold is met, any incoming low priority cells are discarded. On the other hand, high priority cells can access the buffer unless it remains full.
- Otherwise, the queue operates based on FCFS.

• Flowchart





- Two separate buffers, one of which is for high priority cells and the other for low priority cells.
- The high priority queue is always emptied before the low priority queue is served.



- Implementation
  - Topology



#### Simulation environment

- Topology remains the same through simulations.
- Switch processes incoming cells at 10613 cells/sec (? 4.5 Mbps =  $3 \times DS1$ )
- QoS to be secured is 0.0075 for avg. CLR and 0.01 secs for max. queuing delay. - Loads
- Intended load
- CBR at 3537 cells/sec (? DS1) and nrt-VBR at 3537 cells/sec (average)
- rt-VBR at 2358 cells/sec (average) for off-duration and at 7075 cells/sec (average) for on-duration
- Total of 70151 cells expected during a simulation
- Actual load (collected during simulations)
- CBR at 3537 cells/sec and nrt-VBR at 3510 cells/sec (average)
- rt-VBR at 2355 cells/sec (average) for off-duration and at 4800 cells/sec (average) for on-duration
- Total of 67953 cells generated during a simulation

- Results
  - Push-out



- Graphs show the CLR of CBR and the queuing delay by queue size.

#### • Partial buffer sharing (queue size = 100)



- Graphs show the CLR of CBR and the queuing delay by threshold.

#### • Buffer separation (queue size = 100)



- Graphs show the CLR of CBR and the queuing delay by queue size ratio.

• Graphs for performance comparison (queue size = 100)



• Table for cell loss comparison (queue size = 100, total load = 67953)

| T        | С | r | n         | Т  |
|----------|---|---|-----------|----|
| r        | В | t | r         | 0  |
| ٩N       | Ŗ | 5 | <u>\$</u> | ¢, |
| 6        | 0 | Ø | 7         | ð  |
| β        | 8 | B | N A       | \$ |
| 8        | 0 | 2 | b b       | 6  |
|          | 8 | 2 | 8         | Z  |
| <b>V</b> | 6 | 2 | 9         | 9  |
| <b>F</b> | 9 | 2 | Q         | 7  |
|          | 3 | 0 | 8         | 8  |
| Ī        |   | 9 | 2         | 4  |

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# Conclusion and Discussion

- Conclusion
  - As expected, all the three implemented schemes improve the CLR of the high priority traffic by sacrificing the low priority traffic.
- Some queuing schemes may bring an improvement in the CLR of the whole traffic.
- Queuing delay could vary with queuing schemes used though the queue size is fixed (especially in partial buffer sharing).
- Discussion
- Difficulties
- Time-consuming OPNET debugging process, determination of simulation scale for better comparison, clear understanding of relevant existing models required to create user-defined models or attributes

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