



# More Efficient Routing Algorithm for Ad Hoc Network

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ENSC 835: HIGH-PERFORMANCE  
NETWORKS

INSTRUCTOR: Dr. Ljiljana Trajkovic

Mark Wang [mrw@sfu.ca](mailto:mrw@sfu.ca)

Carl Qian [chungq@sfu.ca](mailto:chungq@sfu.ca)



# Outline

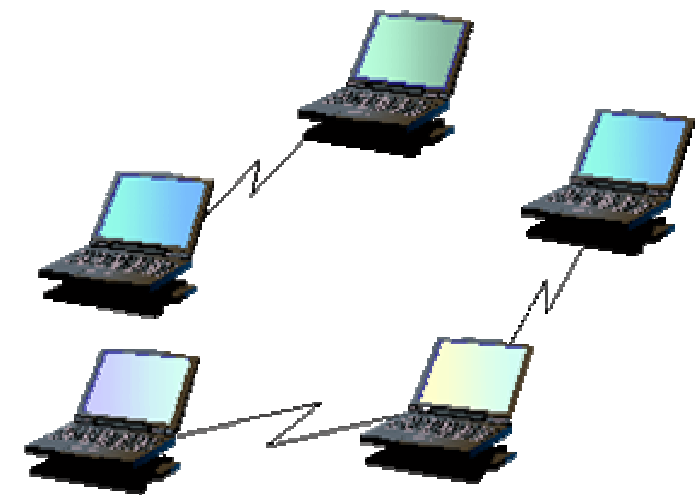
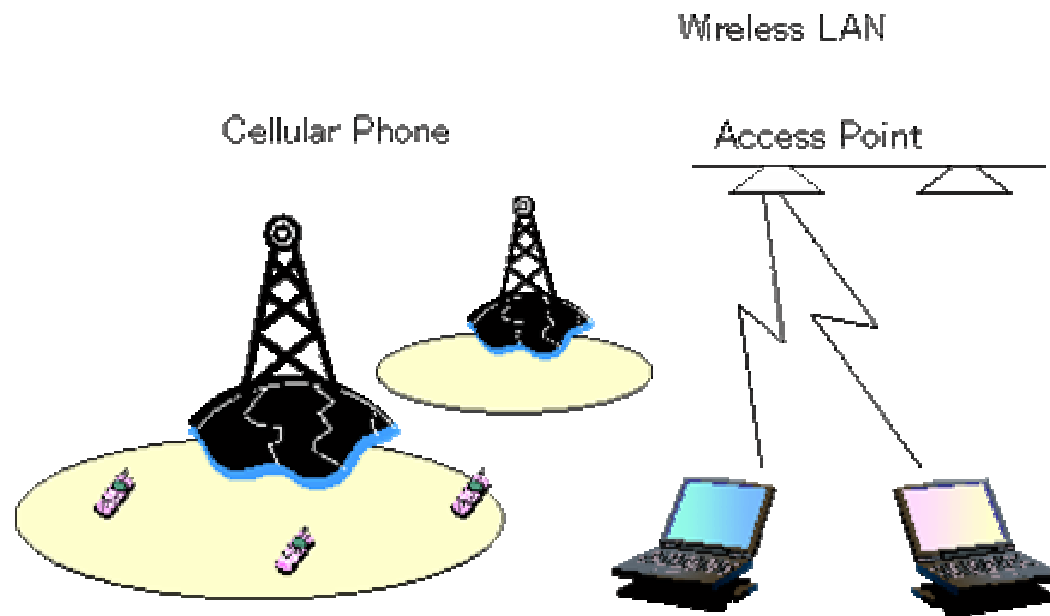
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- Quick Overview of Ad hoc Networks
- AODV Routing Protocols
- Motivation
- Multipoint Relays Select techniques
- Implementation and Challenges
- NS2 Simulation Environment and Results
- Conclusion and Future Works

# Wireless Ad Hoc Networks

Present Mobile Communication

Wireless Ad Hoc Network

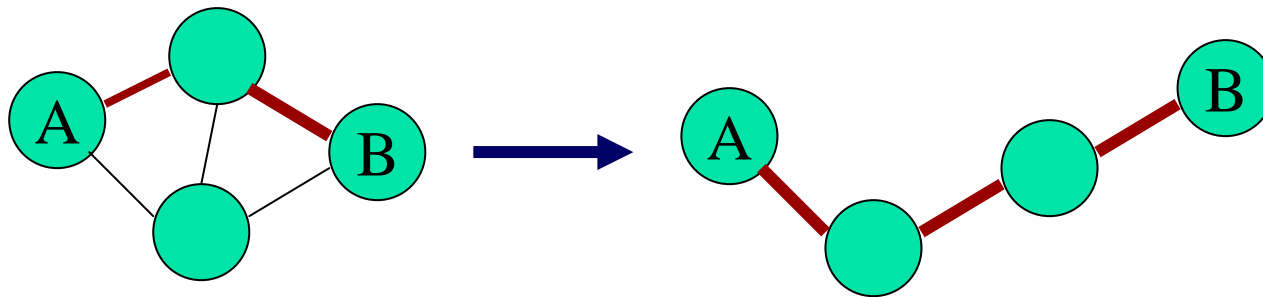


Base station or access point recognizes terminal location and decides communication route.

- ◆ No infrastructure (base station, access point)
  - Network anywhere (disaster stricken area, stadium)
- ◆ Key technologies
  - Routing algorithm
  - Adaptation to network topology change
  - Efficiency in frequency and power

# Mobile Ad Hoc Networks (MANET)

- Host movement frequent
- Topology change frequent



- No cellular infrastructure. Multi-hop wireless links may need to traverse multiple links to reach destination
- Data must be routed via intermediate nodes.



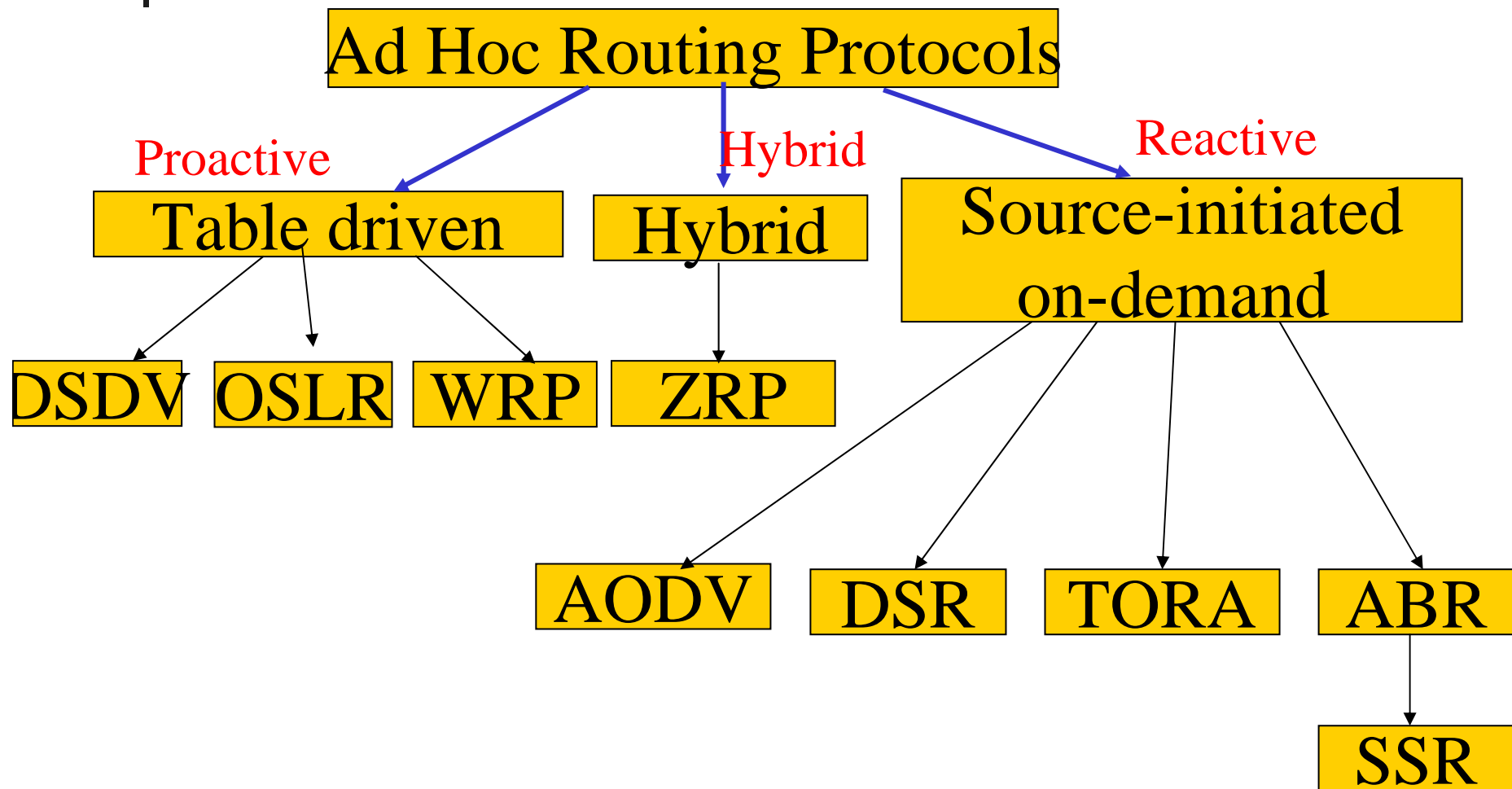
# Unicast Routing Protocols

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- Many protocols have been proposed
- Some specifically invented for MANET
- Others adapted from protocols for wired networks
- No single protocol works well in all environments
  - some attempts made to develop adaptive/hybrid protocols
- Standardization efforts in IETF
  - MANET, MobileIP working groups
  - <http://www.ietf.org>

MANE: Mobile Ad Hoc Networks

# Existing Ad Hoc Routing Protocols



# Routing Protocols



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- **Proactive protocols**
  - Traditional distributed shortest-path protocols
  - Maintain routes between every host pair at all times
  - Based on periodic updates; High routing overhead
  - Example: DSDV (destination sequenced distance vector)
- **Reactive protocols**
  - Determine route if and when needed
  - Source initiates route discovery
  - Example: DSR (dynamic source routing)
- **Hybrid protocols**
  - Adaptive; Combination of proactive and reactive
  - Example : ZRP (zone routing protocol)

# Protocol Trade-offs



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- **Proactive protocols**

- Always maintain routes
- Little or no delay for route determination
- Consume bandwidth to keep routes up-to-date
- Maintain routes which may never be used

- **Reactive protocols**

- Lower overhead since routes are determined on demand
- Significant delay in route determination
- Employ flooding (global search)
- Control traffic may be bursty

- Which approach achieves a better trade-off depends on the traffic and mobility patterns



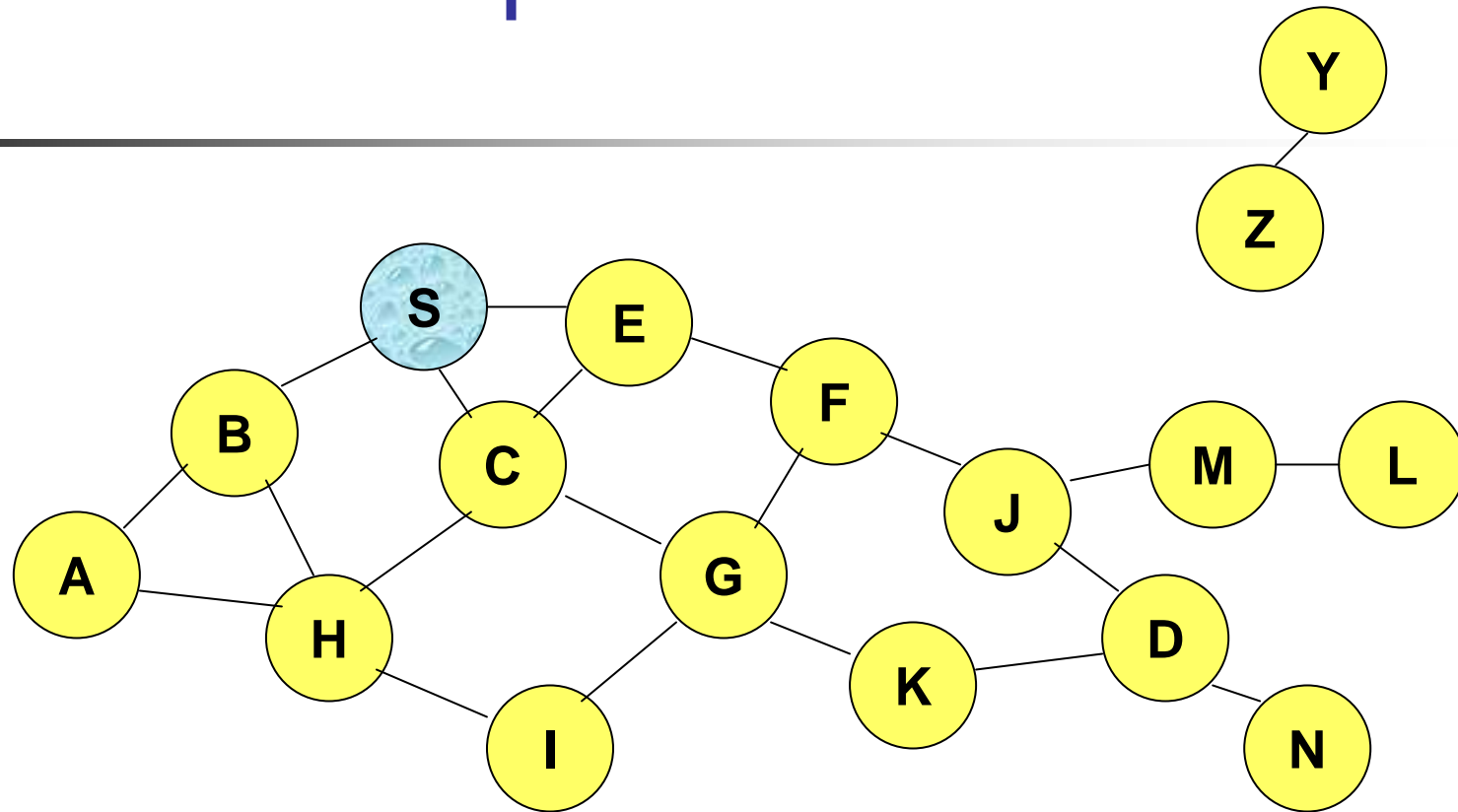
# Ad Hoc On-Demand Distance Vector Routing (AODV)



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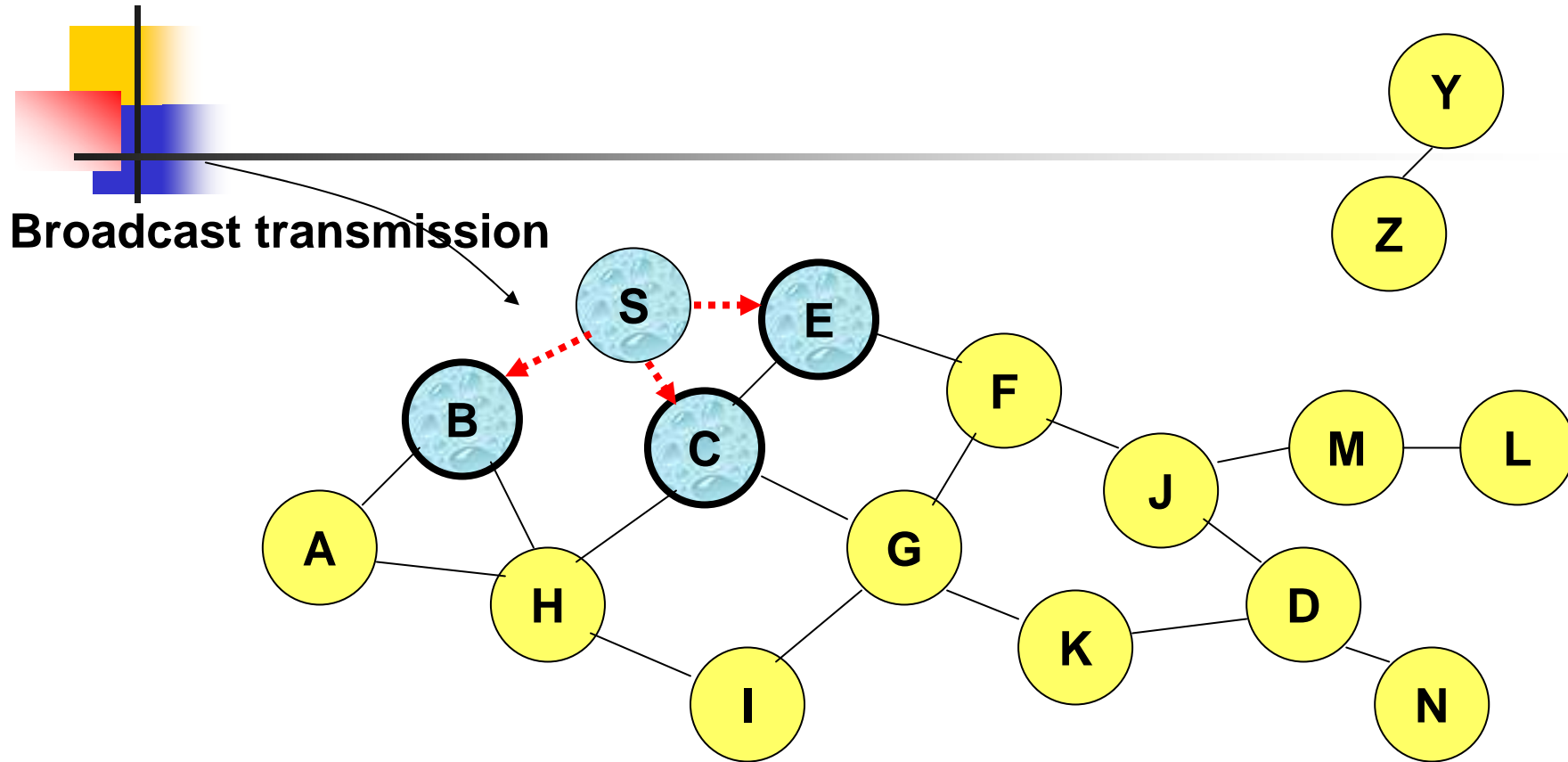
- **Route Requests (RREQ)** are flooded through entire network searching for destination
- When a node re-broadcasts a Route Request, it sets up a reverse path pointing towards the source
  - AODV assumes symmetric (bi-directional) links
- When the intended destination receives a Route Request, it replies by sending a **Route Reply (RREP)**
- Route Reply travels along the reverse path set-up when Route Request is forwarded

# Route Requests in AODV



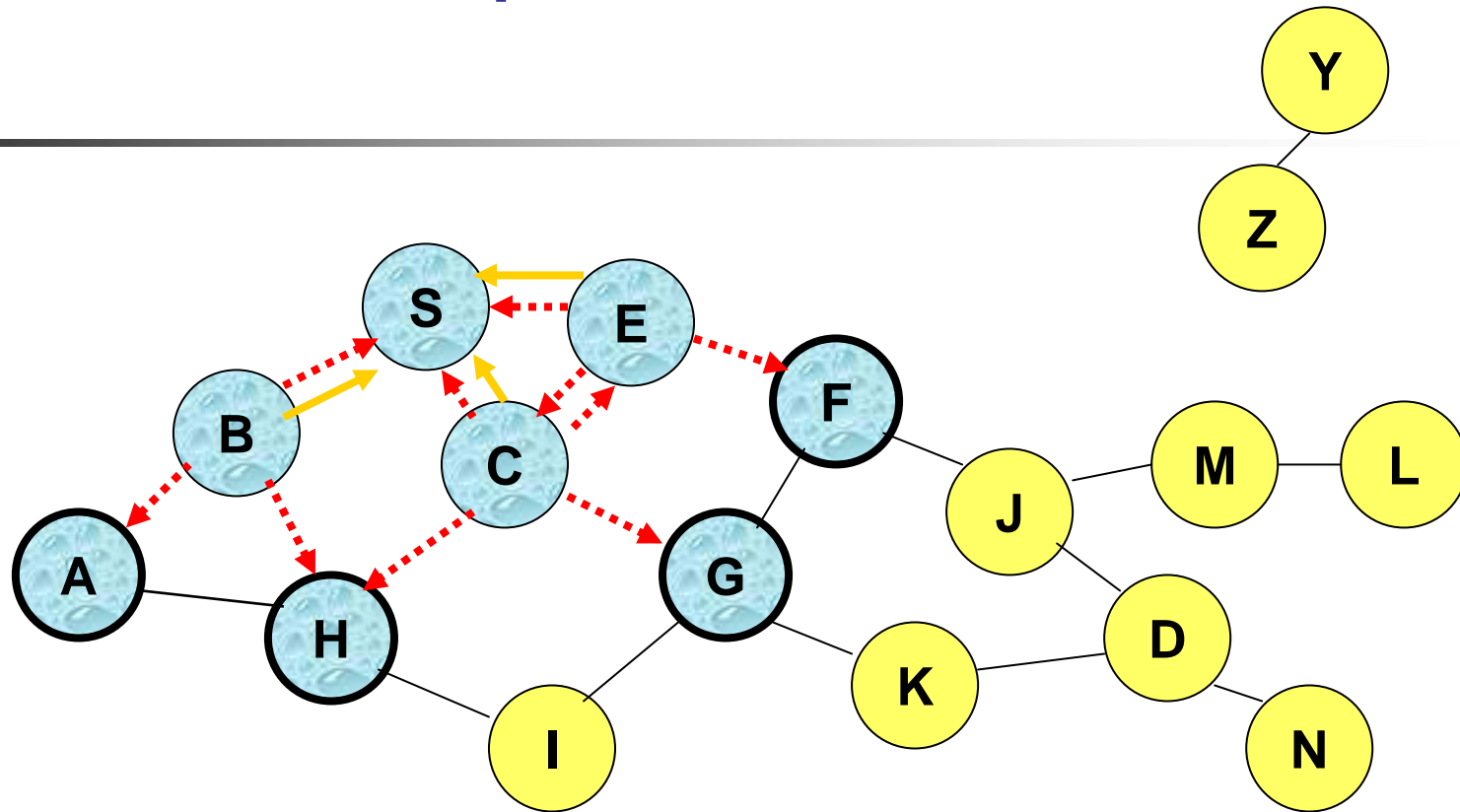
**Represents a node that has received RREQ for D from S**

# Route Requests in AODV



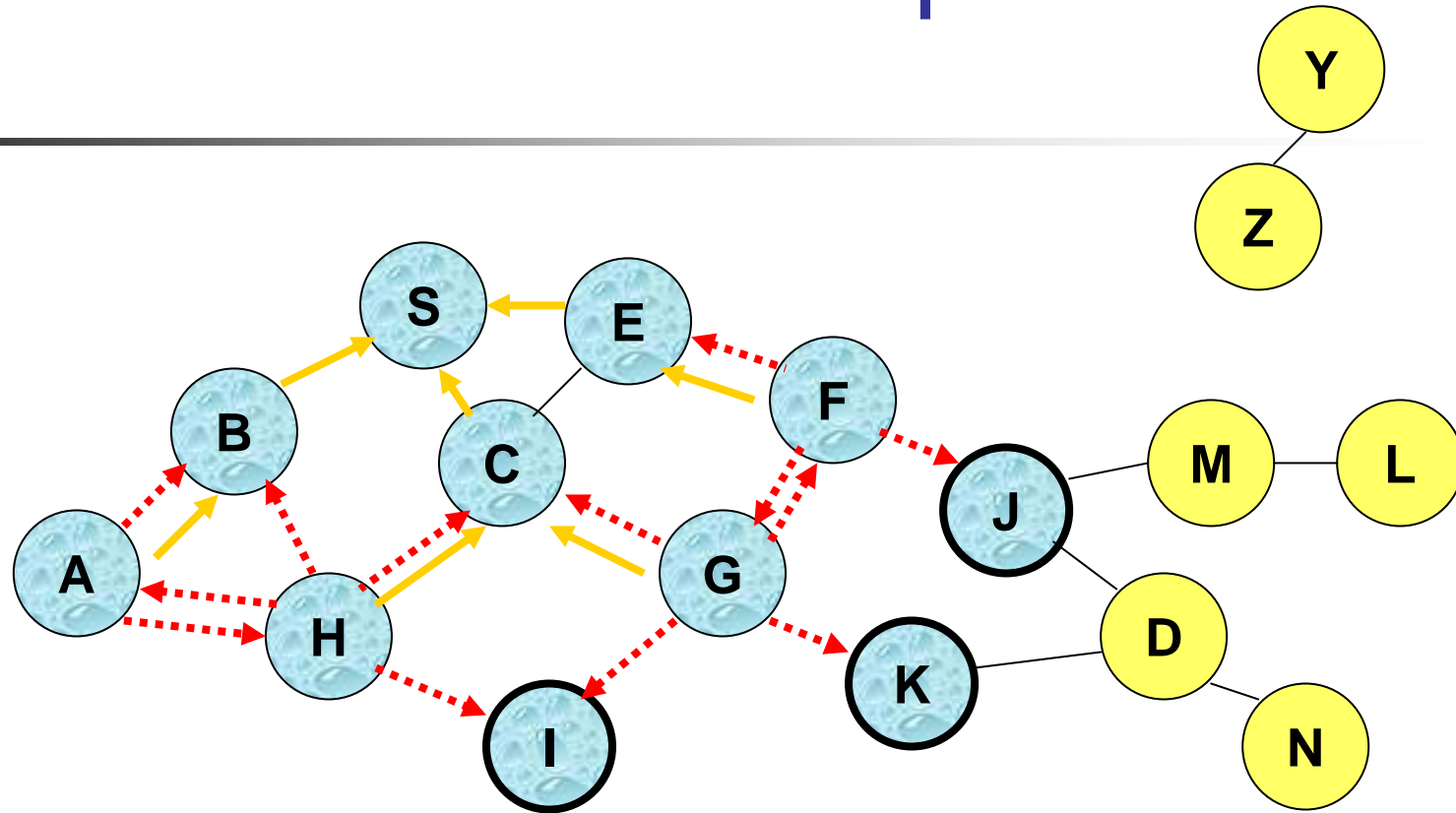
.....→ Represents transmission of RREQ

# Route Requests in AODV



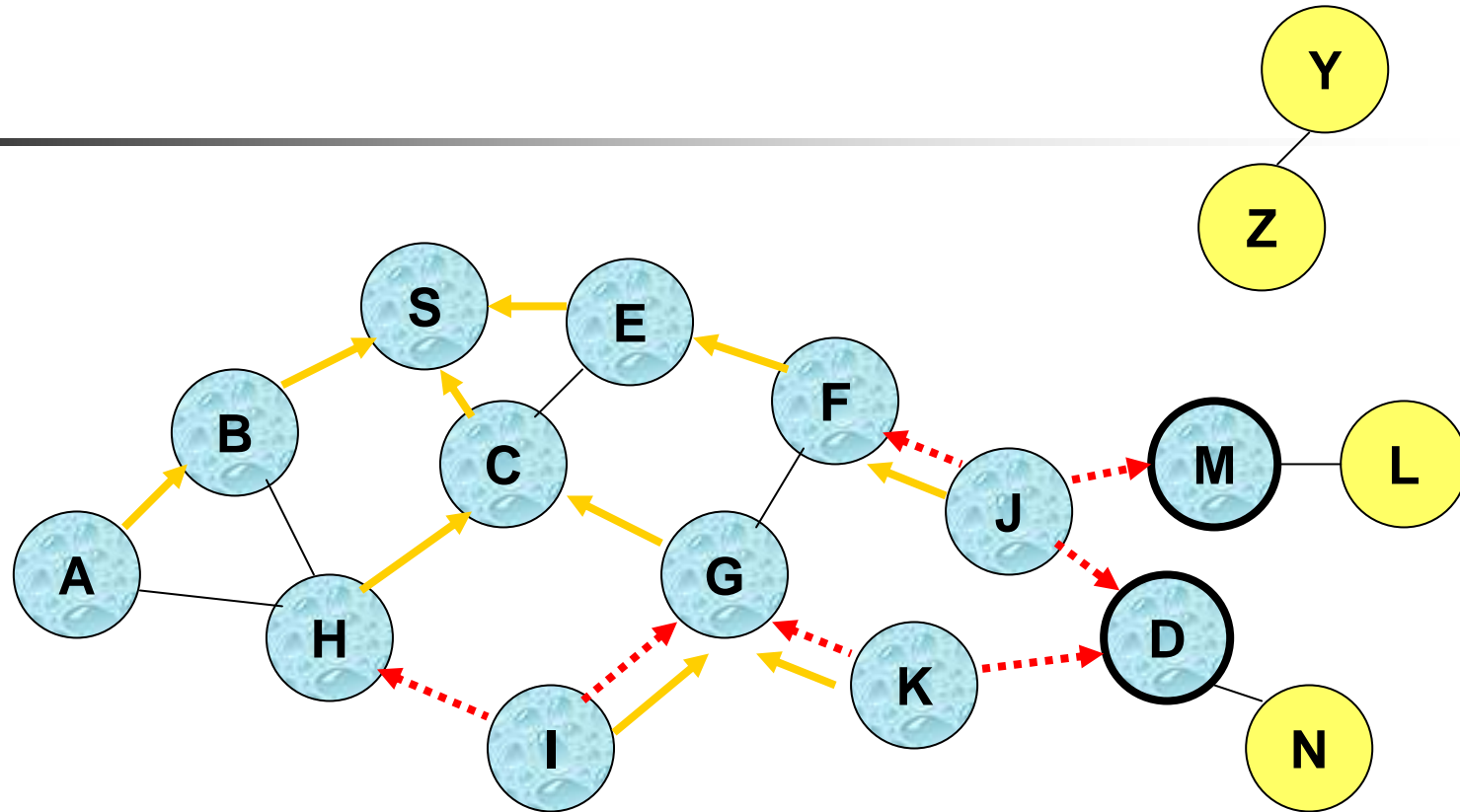
← Represents links on Reverse Path

# Reverse Path Setup in AODV

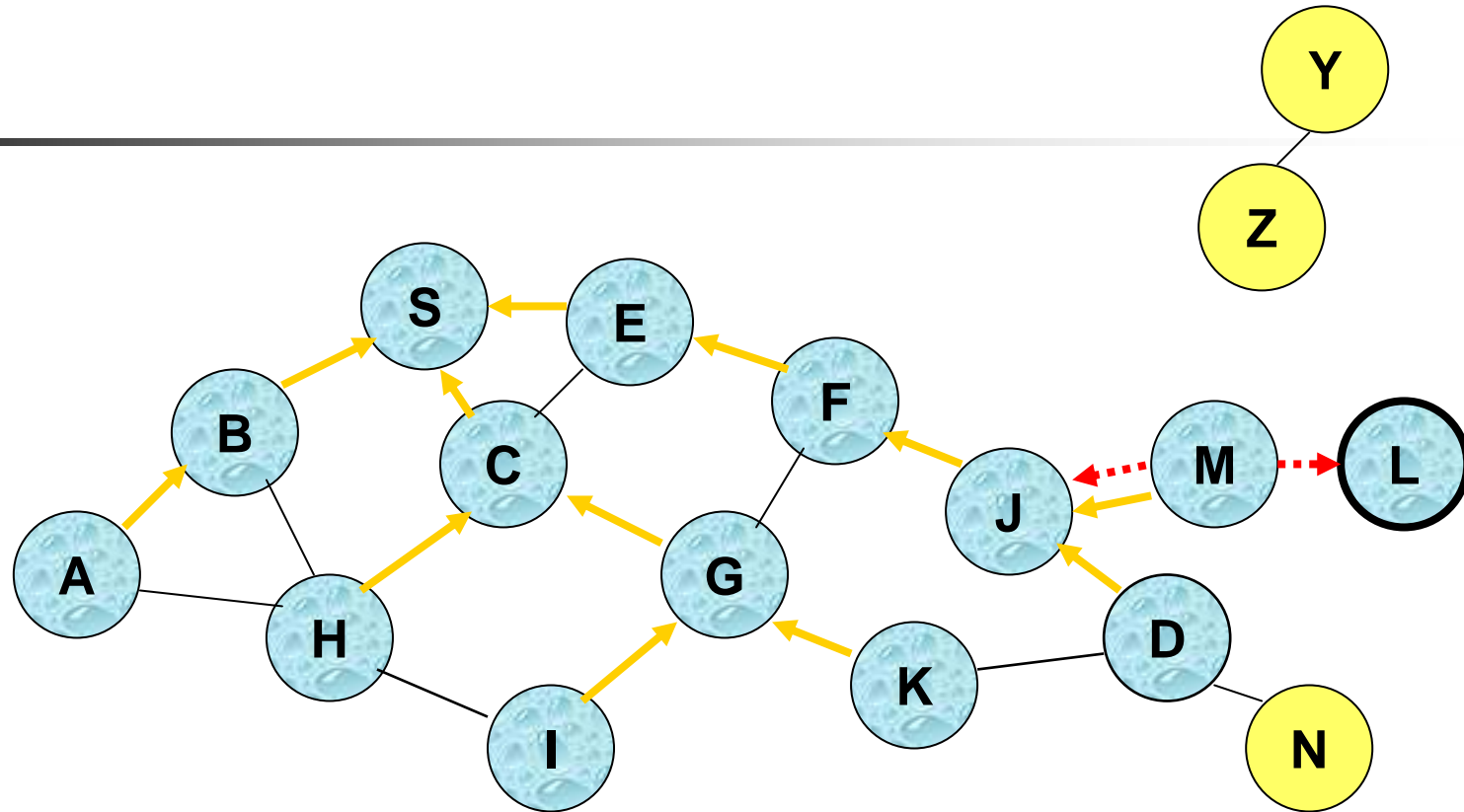


- Node C receives RREQ from G and H, but does not forward it again, because node C has **already forwarded RREQ** once

# Reverse Path Setup in AODV



# Reverse Path Setup in AODV



- Node D **does not forward** RREQ, because node D is the **intended target** of the RREQ







# Motivations

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- The Lack of Scalability of AODV:
  - As the number of source-destination pairs increases
    - Major control overhead of AODV is caused by “Route Query” flood packets
    - Routing overhead is proportional to the number of route queries
  - As the given traffic becomes heavy
    - Heavy routing overhead causes significant effective throughput degradation



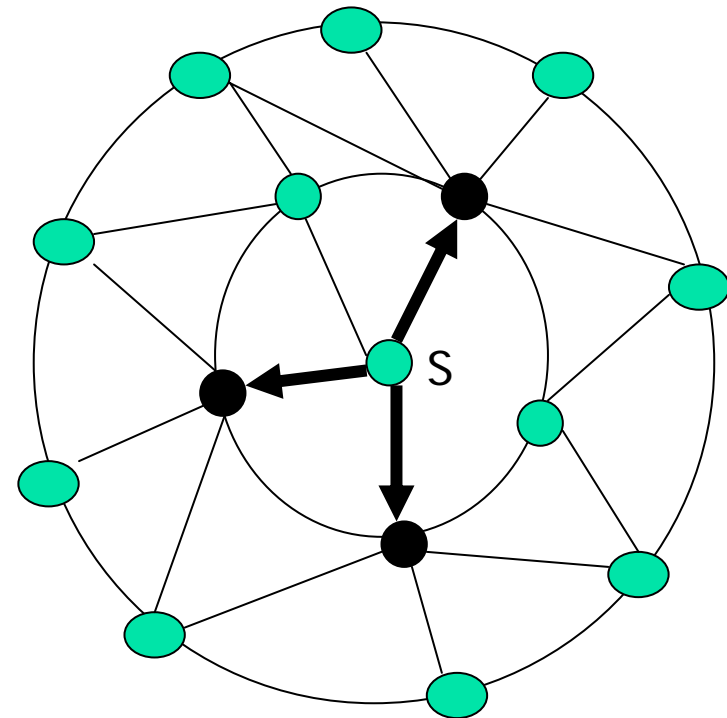
# Proposed Modification

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- Reduce routing overhead of AODV using **Efficient Flooding (Selective Flooding)**
- What's efficient flooding?
  - Only a subset of nodes (dominating nodes) forwards a Route Query flood packet
  - In contrast, in blind flooding all nodes relay each packet at most once
- How to choose dominant nodes?
  - Multipoint Relay Sets (MPRs)

# Multipoint Relay (MPR)

- The Concept of MPR is to reduce the number of duplicated retransmissions while forwarding a broadcast packet
- Multipoint relay set (MPRs): subset of a node's 1-hop neighbors, such that each of its 2-hop neighbors is a 1-hop neighbor of a node in the MPR set



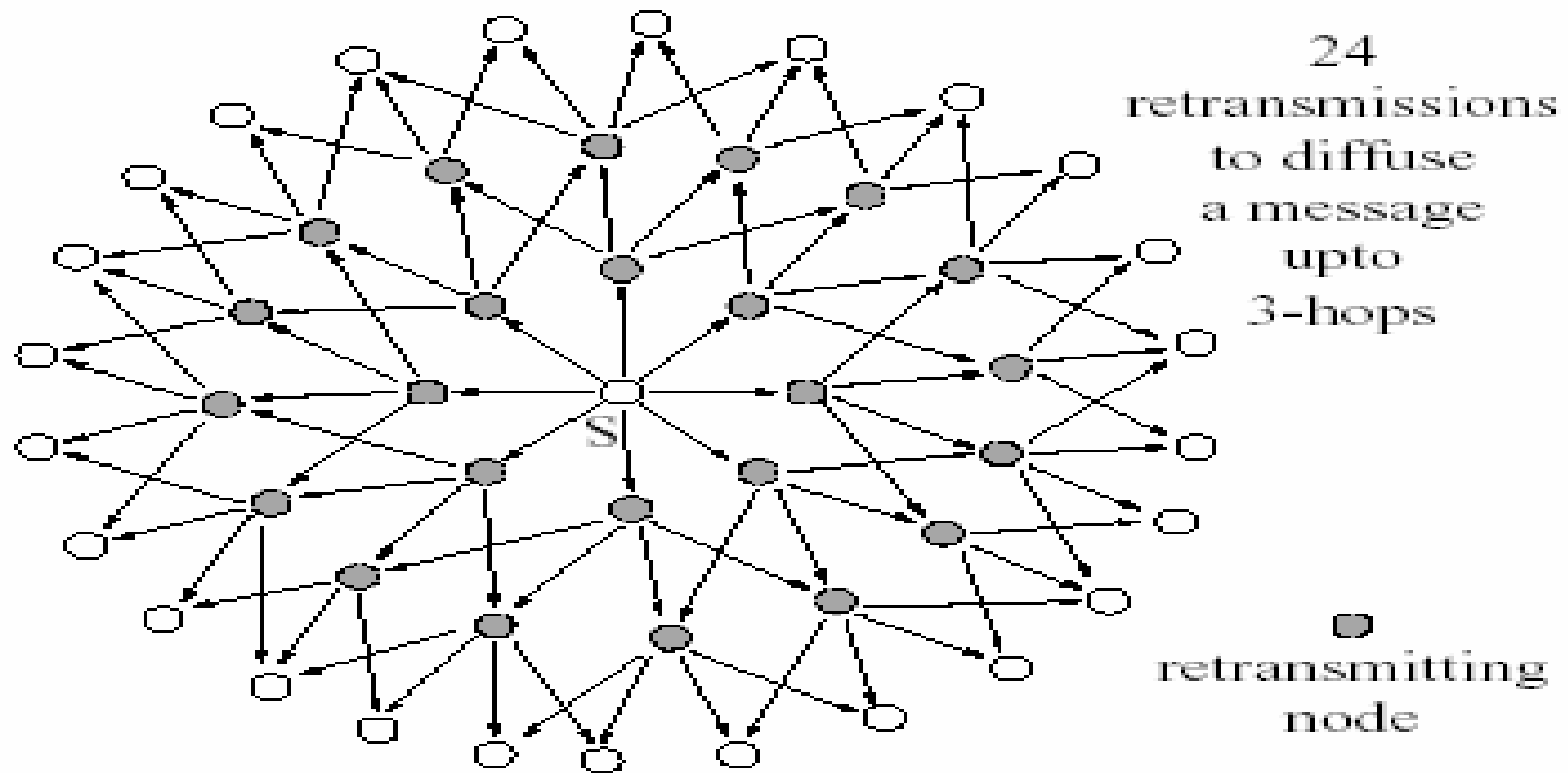


# Multipoint Relay

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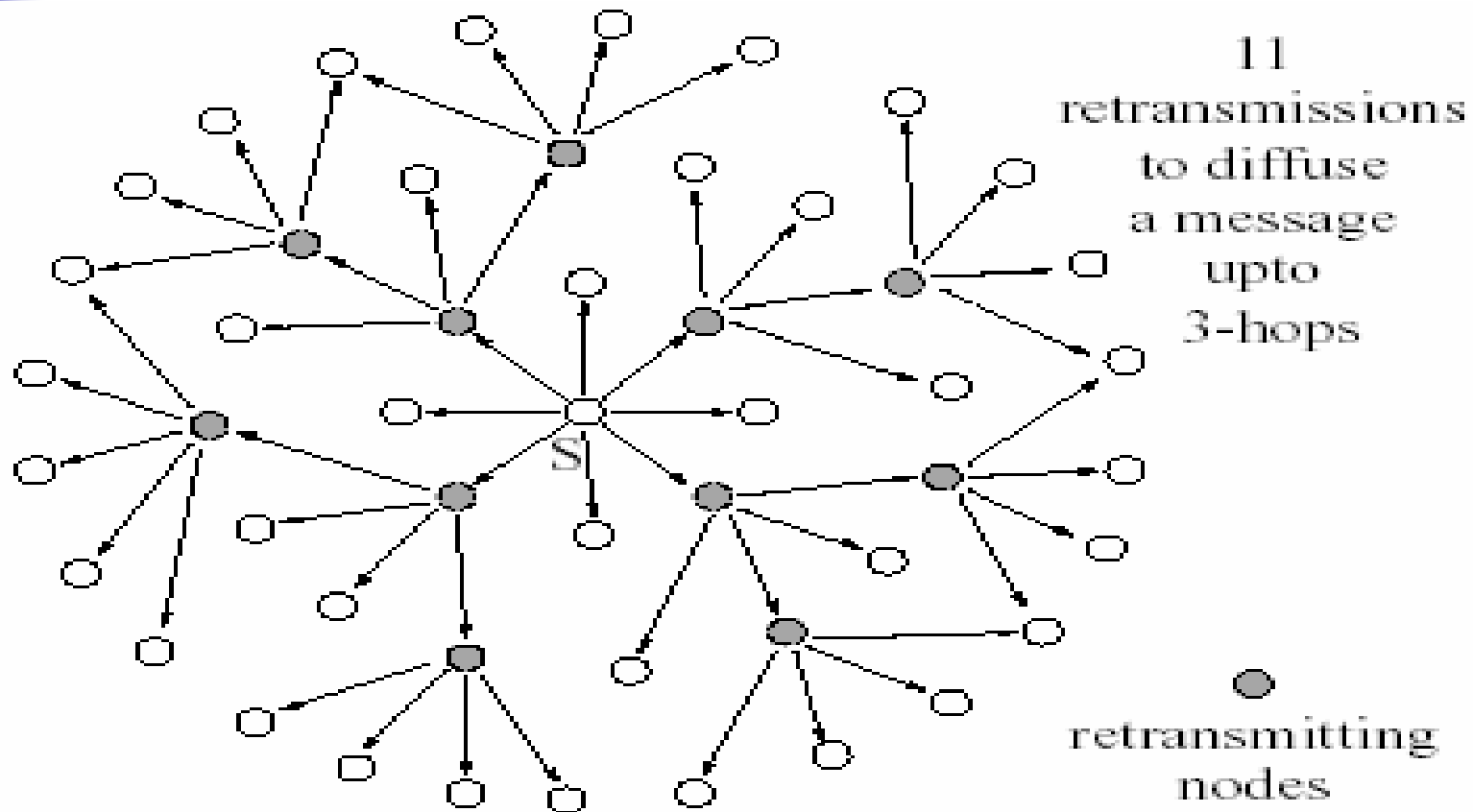
- A node selects its Multipoint relays with two rules:
  - Any 2-hop neighbors must be covered by at least one multipoint relay
  - Try to minimize the multipoint relay set
- Note that each node independently determines its own MPR set (no global “network MPR set”)
- A node forward a flooding packet with the following rules:
  - The packet has not yet been received.
  - The node is multipoint relay of last emitter

# Multipoint Relay



Diffusion of broadcast message using pure flooding

# Multipoint Relay



Diffusion of broadcast message using multipoint relays



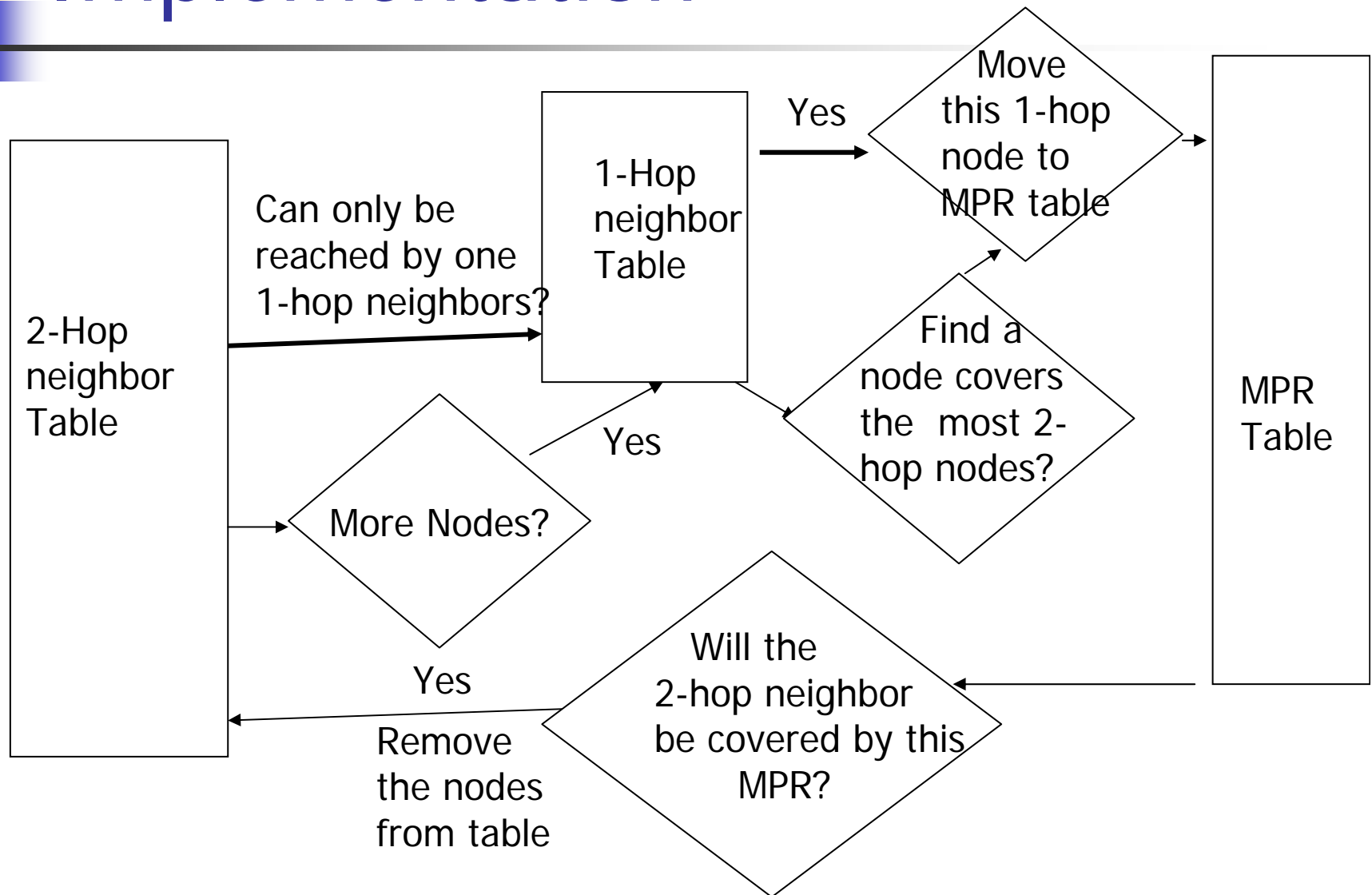
# Implementation

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**The algorithm for calculating the multipoint relay table is shown below:**

1. Find all 2-hop neighbors that can only be reached by one 1-hop neighbor. Assign those 1-hop neighbors as MPRs.
2. Determine the resultant cover set (i.e., the set of 2-hop neighbors that will receive the packet from the current MPR set).
3. From the remaining 1-hop neighbors not yet in the MPR set, find the one that would cover the most 2-hop neighbors not in the cover set.
4. Repeat from step 2 until all 2-hop neighbors are covered.

# Implementation







# Challenges

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- Because of the nature of Ad Hoc network, nodes are moving constantly. We have to keep updating each node's movement and their neighbors.
- Each node must have the 1-hop and 2-hop neighbor information at any given time.
- This information can only be obtained by exchanging message periodically

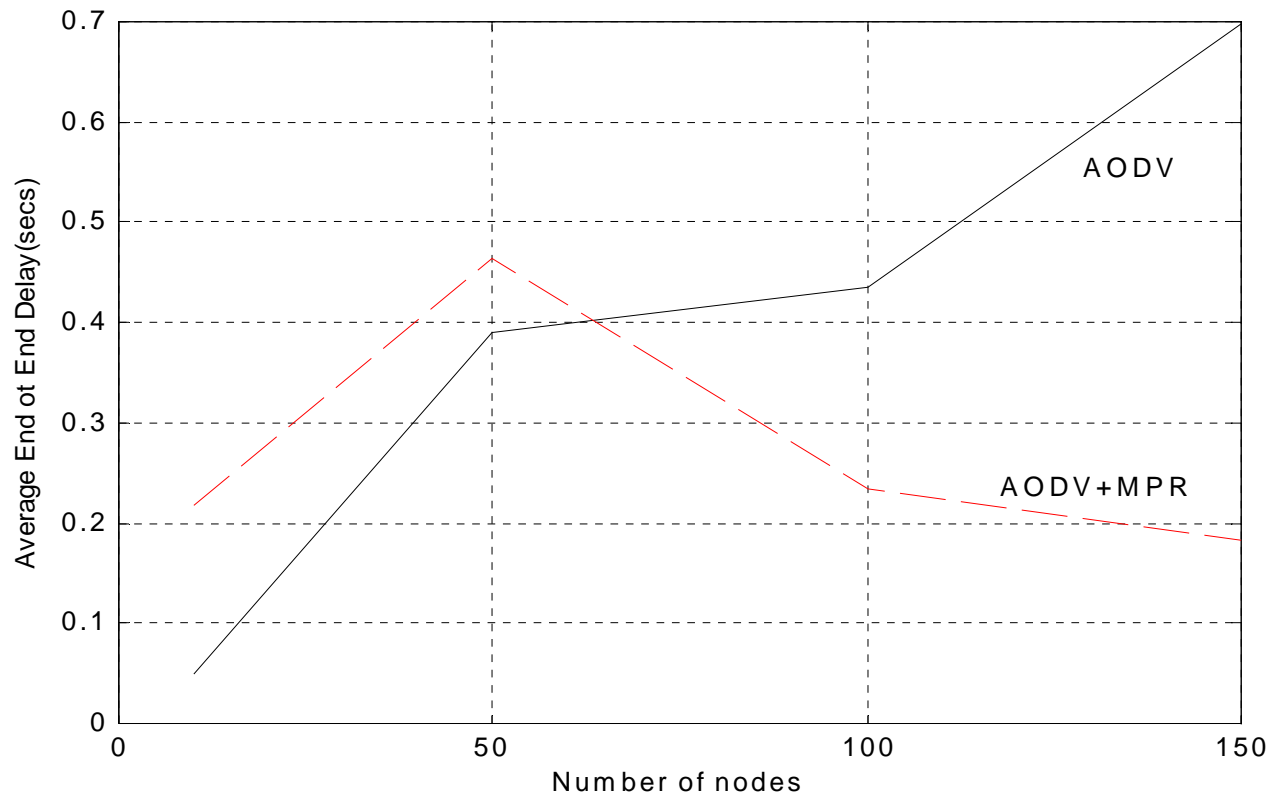


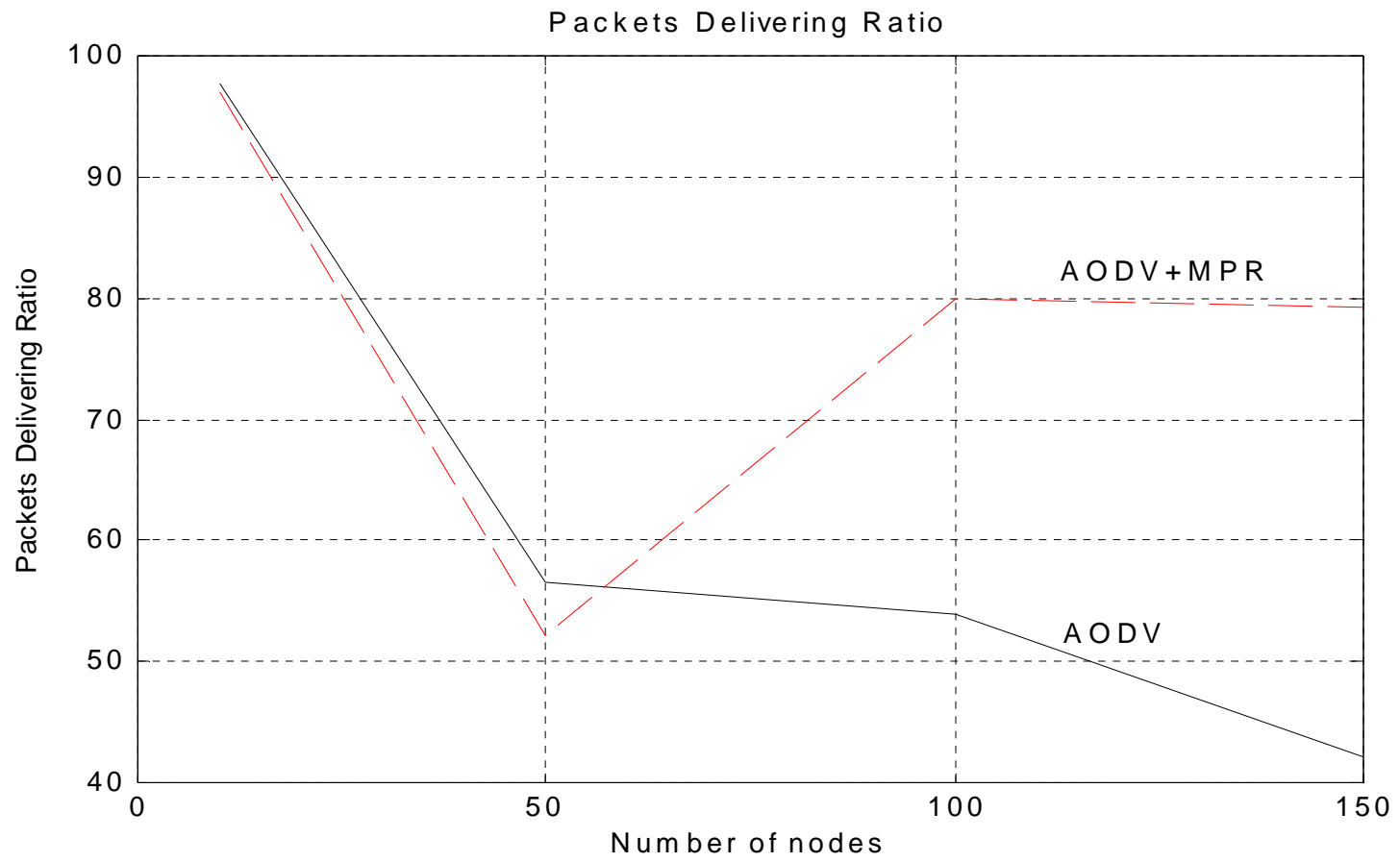
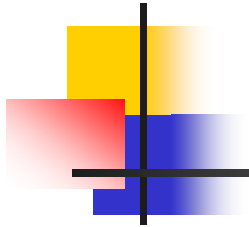
# NS2 Simulation Environment

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- Simulator: NS2-2.26
- Operating System: Linux
- Network Area: 900 \* 900 meters
- Number of nodes simulated: 10, 50, 100, 150
- Max. pause time: 10s
- Max. speed: 20m/s

# Results







# Conclusion

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- Our simulations show the MPR technique improves AODV protocol significantly by reducing the overhead and delay in dense node networks
- With this technique, AODV can achieve better package delivery ratio

MPR: Multipoint Relay

AODV: Ad Hoc On-Demand Distance Vector Routing



# Future Works

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- The AODV protocol can be further optimized by applying other techniques such as probability based methods or location based methods



# References

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Thank you!

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Any Questions And Comments?