ENSC 835-3: HIGH-PERFORMANCE NETWORKS

DYNAMIC RIGHT-SIZING A TCP FLOW-CONTROL INVESTIGATION IN OPNET

by

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ROADMAP

- □ Introduction
- □ Flow Control in TCP
- OPNET Investigation of Dynamic Right Sizing
- □ Results
- □ References

INTRODUCTION

- Networking researchers typically manually optimizing TCP buffer sizes
- □ To keep the network pipe full
- Achieve acceptable performance over the wide-area network.

INTRODUCTION

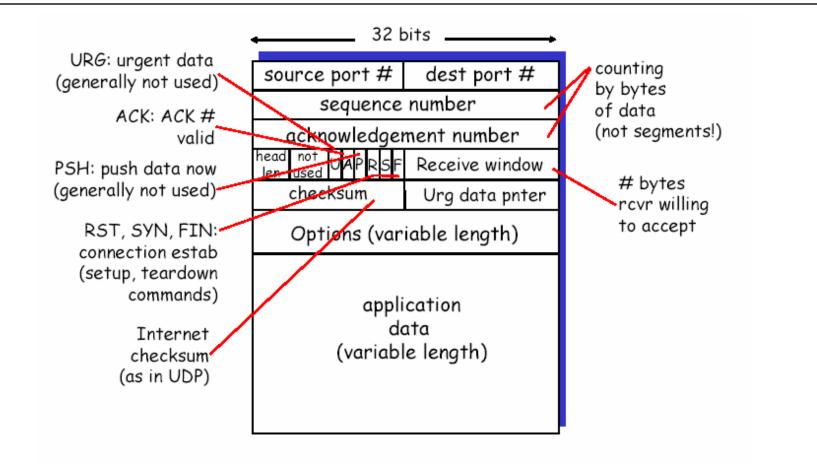
□ Cumbersome

- Sub-par performance for connections with larger delay-bandwidth products
- Misappropriation of scarce resources to connections with smaller delay-bandwidth products

INTRODUCTION

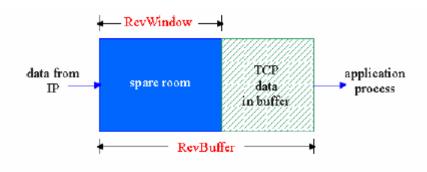
- Originally implemented in ns
- Dynamic Right sizing lets the receiver estimate the sender's congestion window size
- □ Uses this estimate to dynamically change the size of the receiver's window advertisements.
- As a result, the sender will be congestionwindow-limited rather than flow-controlwindow-limited

TCP SEGMENT STRUCTURE



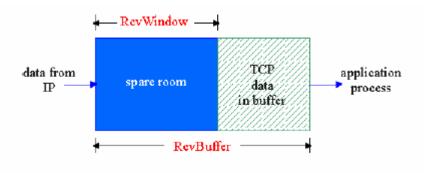
FLOW CONTROL IN TCP

- Receive side of TCP connection has a receive buffer
- □ App process may be slow reading from buffer
- Won't overflow receiver's buffer by transmitting too much, too fast

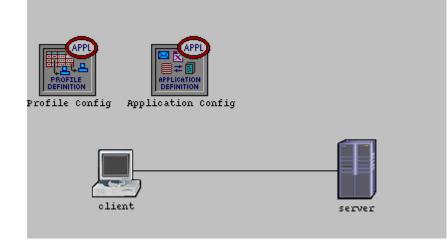


FLOW CONTROL IN TCP

- RcvWindow = RcvBuffer-[LastByteRcvd -LastByteRead]
- Rcvr advertises spare room by including value of **RcvWindow** in segments
- Sender limits unACKed data to RcvWindow guarantees receive buffer doesn't overflow



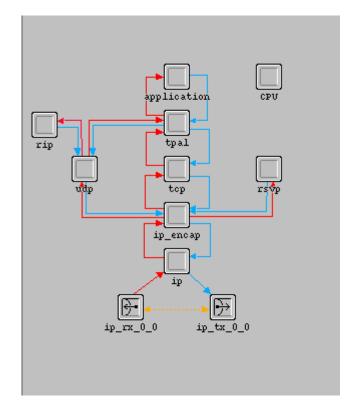
OPNET SETUP



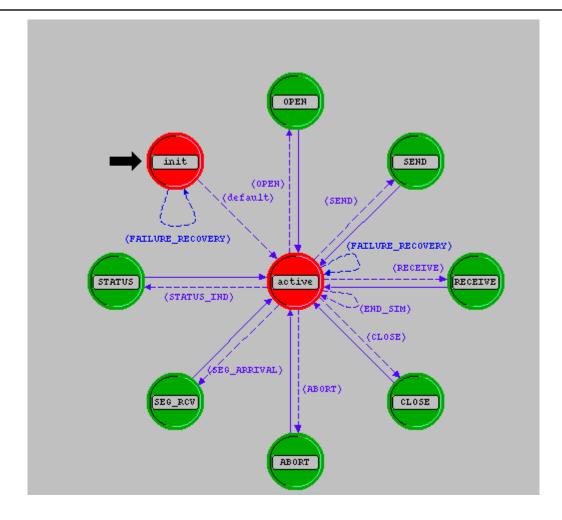
OPNET SETUP

(Application Config) Attributes Type: utility			
Attribute Application Definitions Application Applicatio	Value Application Config Application Config () 1 File Transfer (custom) () Off Off	★ (Ftp) Table Attribute Command Mix (Get/Total) Inter-Request Time (seconds) File Size (bytes) Symbolic Server Name Type of Service RSVP Parameters Back-End Custom Application Details Promote	Value 100% constant (3600) constant (1600000) FTP Server Best Effort (0) None Not Used
Voice Encoder Schemes Apply changes to selected objects Eind Next	All Schemes		

TCP NODE MODEL



TCP PROCESS MODEL



TCP MODEL PROCESSES

- \Box tcp_manager_v3.c
- □ Represents the root process of the "tcp" module.
- □ It manages a set of TCP connections by invoking the appropriate api processes

IMPLEMENTATION STRATEGY

- Receiver can determine the sender size by below equation
- □ d is bytes of data received
- □ round trip times between Nmin and Nmax

$$\frac{d}{n_{max}} \le w \le \frac{d}{n_{min}}$$

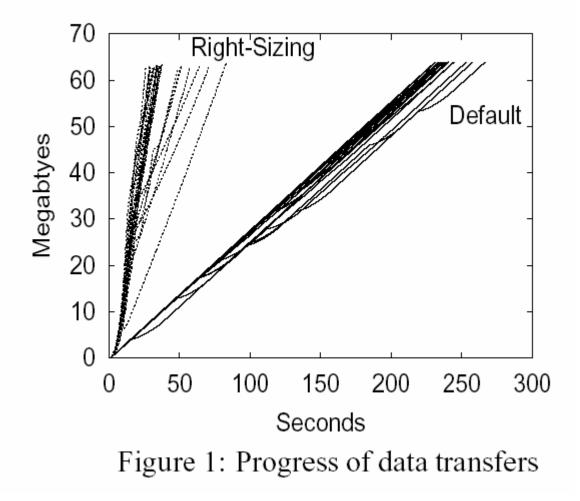
IMPLEMENTATION STRATEGY

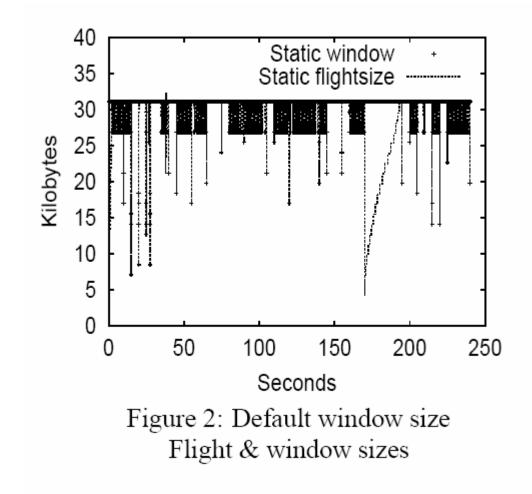
- Observing the Congestion Window variable change
- Altering the receiver buffer window accordingly

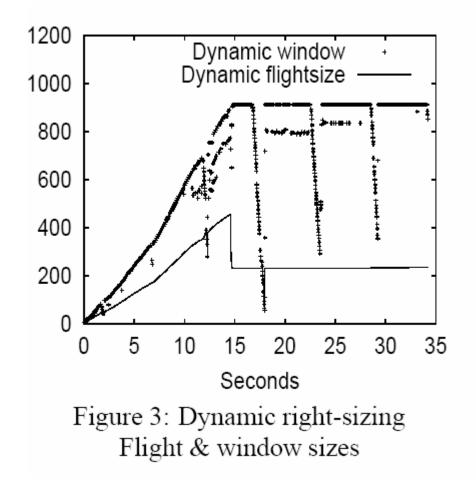
STATISTICS OBSERVED

- □ Delay
- □ Receiver Window Size
- □ Congestion Window Size
- Flight Size (sent but unacked data in sender buffer)
- □ Traffic Received

- The first 25 transfers used the default window sizes of 64 kilobytes for both the sender and receiver.
- The second 25 transfers, shown in dotted lines, used dynamically sized windows







REFERENCES

□ Mike Fisk and Wu-chun Feng, Dynamic Right-Sizing: TCP Flow-ControlAdaptation, in Proceedings of the 14th Annual ACM/IEEE SC2001 Conference,November 2001. <u>http://public.lanl.gov/radiant/pubs/drs/sc2001-poster.pdf</u>

□ RFC1323 – TCP Window Scale Option

□ Kurose and Ross – Chapter 3 presentation slides pg 72, 73 Computer Networking: A Top Down Approach Featuring the Internet, 3rd edition. Jim Kurose, Keith Ross Addison-Wesley, July 2004.

□ Van Jacobson, Congestion Avoidance and Control, in Proceedings, SIGCOMM 1988 Workshop. ACM SIGCOMM, Aug. 1988, pp. 314-329, ACM Press, Stanford, CA.

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