ENSC 835: High-Performance Networks Spring 2008

Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Network

Chih-Hao Howard Chang howardc@sfu.ca

Final Project Demo School of Engineering Science Simon Fraser University

Roadmap



- Project Description
- Modified MobileNode in ns-2
- AODV Routing in ns-2
- Modified AODV with Interface Switching Capability
- Simulation
 - Configuration
 - Interface Switching
 - Sample ns-2 Output
 - O Average Throughputs
 - Observations

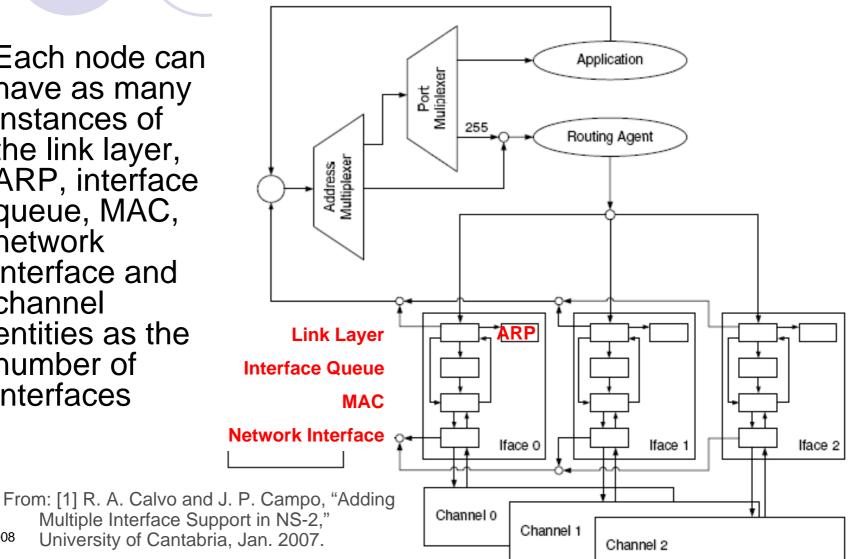
Project Description

- Extend the Network Simulator (ns-2.32) to support multiple channels and multiple interfaces by referring to [1]
 - O modify ns-2's MobileNode library to support multiple interfaces
 - preserve the legacy operations of IEEE 802.11 interfaces
- Implement the interface switching protocol proposed in [2]
 integrate this algorithm in the existing AODV routing agent in ns-2
- Simulate a multi-channel multi-interface ad-hoc wireless network (in chain topologies) using the modified ns-2
 - demonstrate the effectiveness of interface switching and the improvement in the network throughput
- AODV Ad-hoc On-demand Distance Vector
- [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2," University of Cantabria, Jan. 2007 (User Guide).
- [2] P. Kyasanur and N. H. Vaidya, "Routing and Link-layer Protocols for Multi-Channel Multi-Interface Ad Hoc Wireless Networks," SIGMOBILE Mobile Computing and Communications Review, vol. 10, no. 1, pp. 31-43, Jan. 2006.

Modified MobileNode in ns-2

Each node can have as many instances of the link layer, ARP, interface queue, MAC, network interface and channel entities as the number of interfaces

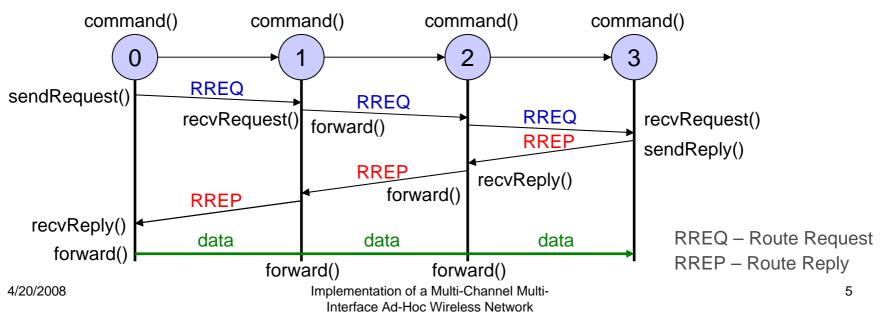
4/20/2008



AODV Routing in ns-2

Route establishment:

- Source node broadcasts a RREQ to find a route to destination node
- O each node receiving the RREQ forwards to the next node
- a route is determined when the RREQ reaches a node that offers accessibility to destination node
- the route is established by sending a RREP back to source node



Modified AODV with Interface Switching Capability

command()

- initially, the node chooses a random channel for its fixed interface and switchable interface
- add the fixed channel used by this node to its NeighbourTable
- O update the node's ChannelUsageList with its fixed channel
- sendRequest(), sendReply(), sendHello()
 - add the fixed channel used by this node and its NeighbourTable to the outgoing RREQ, RREP, or Hello packet
- recvRequest(), recvReply(), recvHello()
 - when the node receives a RREQ, RREP, or Hello packet from a neighbour, it updates:
 - the node's NeighbourTable with the fixed channel of that neighbour
 - the node's ChannelUsageList using the NeighbourTable of its neighbour.

Modified AODV with Interface Switching Capability

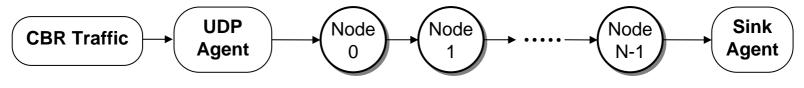
forward()

- RREQ or RREP:
 - add the fixed channel used by this node and its NeighbourTable to the outgoing packet
- O data:
 - consult the node's ChannelUsageList, find the channel with the largest usage
 - if the node's fixed channel has the largest usage:
 - with a probability of 0.4 (from paper [2]), the node:
 - reverses its ChannelUsageList about the fixed channel previously used
 - changes its fixed channel to a less used channel
 - transmits a new Hello packet informing neighbours of its new fixed channel → sendHello()
 - if the usage of the node's fixed channel is ok:
 - look up the fixed channel of the next node in NeighbourTable
 - assign this fixed channel to the node's switchable interface

Simulation: Configuration

Using the modified ns-2.32:

 test the effectiveness of interface switching and throughput in multi-channel and multi-interface ad-hoc wireless networks



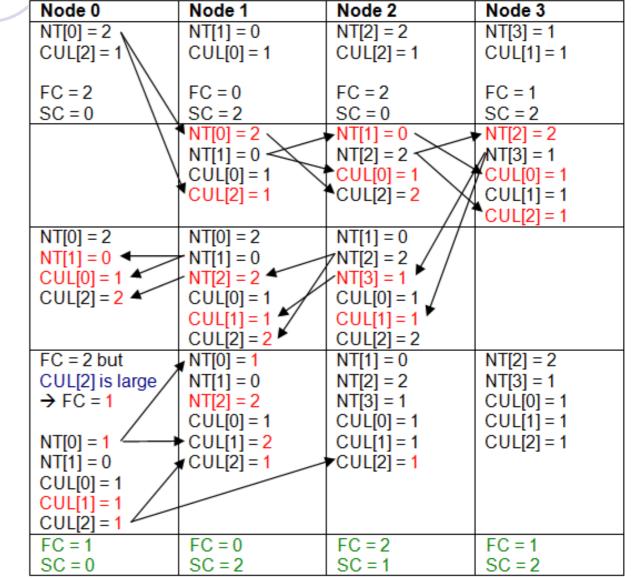
Scenarios:

- Simple chain topologies of 2-11 stationary nodes (single route)
- 2-4 channels per node
- 2 interfaces per node (1 fixed, 1 switchable)
- simulation duration: 60 s (actual simulation time is slightly longer)
- O channel capacity: 5.4 Mbps
- constant bit-rate (CBR) traffic flow from Node 0 to Node N-1
 - transmitted over UDP (no flow and congestion controls)
 - 1000 bytes per packet, sent every 1.4 ms

Simulation: Interface Switching

- Example: 4 nodes,
 3 channels [0, 1, 2],
 2 interfaces [0, 1]
- NT[n] is the fixed channel used by node n
- CUL[c] is the number of nodes using c as their fixed channel
- Fixed Channel (FC)
- Switchable Channel (SC)

4/20/2008



Simulation: Sample ns-2 Output

4 nodes, 3 Channels, 2 interfaces:
 CBR traffic from Node 0 to Node 3
 Average throughput is 2586.36 kbps

```
/usr/ns-allinone-2.32/mcmi test
                                                                        nerry@BENQ-JOYBOOK /usr/ns-allinone-2.32/mcmi_test
 ns wireless_chain_mcmi.tcl
Ad-Hoc Wireless Network in Chain Topologies - 4 Nodes, 3 Channels, 2 Interfaces
nerry@BENQ-JOYBOOK /usr/ns-allinone-2.32/mcmi_test
          flowID:
                   И
     trafficType:
                   cbr
         srcNode:
                   И
        destNode:
                   з
    startTime[s]:
                   1
    stopTime[s]:
                   61
    duration[s]:
                   6Й
    receivedPkts: 19468
   avgTput[kbps]:
                   2586.36
```

Simulation: Sample nam Output

Network Animator (nam) in ns-2.32:

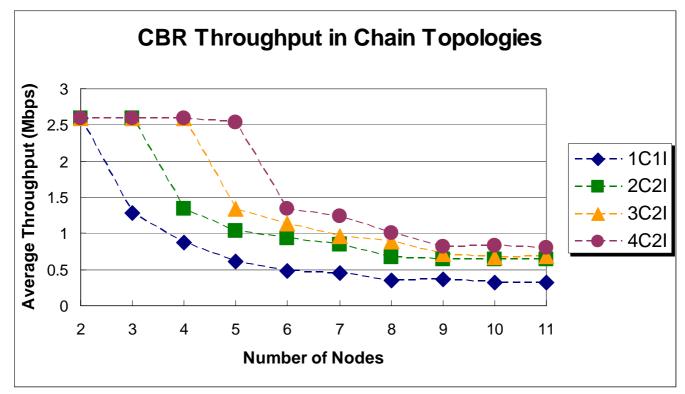
- O only nodes can currently be seen
- dumping of traffic data and thus visualization of data packet movements for wireless scenarios is still not supported

The following is a chain topology with 4 wireless nodes:

χ nam: wireless_cha	in_n4_c3_i2.nam	Aude manent	And And	Gurrantly See as		
<u>File Views Anal</u>	lysis	wireless_chain_n4_c3_i2.nam				
			•		0.00000	Step: 2.0ms
↔ §		\bigcirc		ଡ		3
		U		e		
					1 1 1 1	

Simulation: Results

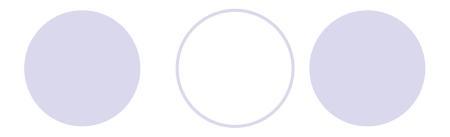
 Average throughputs while varying the number of channels, n → nC2I vs. 1C1I



Simulation: Observations

- The throughput of 1C1I networks degrades as the number of nodes increases by 1 each time
 - intermediate nodes cannot send and receive data at the same time
 - interference within the carrier sense range
- Higher throughput with multiple channels and 2 interfaces on each node
 - interface switching assigns the fixed channel of successive nodes to different channels
 - intermediate nodes can send data to the next node using its switchable interface, while receiving data on its fixed interface
- Smaller throughput improvement when the number of nodes > number of channels + 1
 - \bigcirc some nodes will be on some common channels \rightarrow interference
 - O however, generally still higher than the case of 1C11

Questions?



Thank you!