

ENSC 835: HIGH-PERFORMANCE NETWORKS

Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Network

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Final Project Report

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Abstract

The IEEE 802.11 standard defines multiple non-overlapping channels at the physical layer in the 2.4 GHz and 5 GHz spectrums. However, most ad-hoc wireless networks today are configured to operate under a single channel in order to ensure connectivity of all their nodes. Hence, the aggregate bandwidth provided by the radio spectrums is not fully-utilized. In order to meet the high throughput demand, it is essential to use all available spectrums. Several past research proposals have exploited multiple channels and multiple interfaces to increase the network capacity, by having multiple simultaneous transmissions without interference.

In this project, we implement and evaluate the technique of *interface switching* to utilize multiple channels and multiple interfaces using the Network Simulator (ns-2). This includes incorporating the multiple channel and multiple interface supports to the core of ns-2, which are not available in the current version of the simulator. The simulation results demonstrate the effectiveness of the approach in improving the network capacity.

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1. Introduction

In recent years, wireless local area networks (WLAN) have become popular, as they provide us the mobility to move around within a broad coverage area and still be connected with each other. The increasing popularity of WLANs is primarily due to their convenience, cost efficiency, and ease of integration with other networks and network components. More electronic and computer devices sold to consumers nowadays come pre-equipped with a wireless network interface card. This trend, along with reducing wireless hardware costs, has accelerated the use of WLANs, increasing the throughput demand.

IEEE 802.11 is a widely used set of standards for WLANs. Endpoint devices (or hosts) designed for this wireless technology can be operated in two modes. In infrastructure mode, hosts are connected via a base station, also known as an access point, which provides network services such as address assignment and routing. Another mode is ad-hoc, whereas wireless hosts have no such infrastructure with which to connect; the hosts must provide for these network services. Figure 1 illustrates the fundamental differences between these two types of wireless networks.

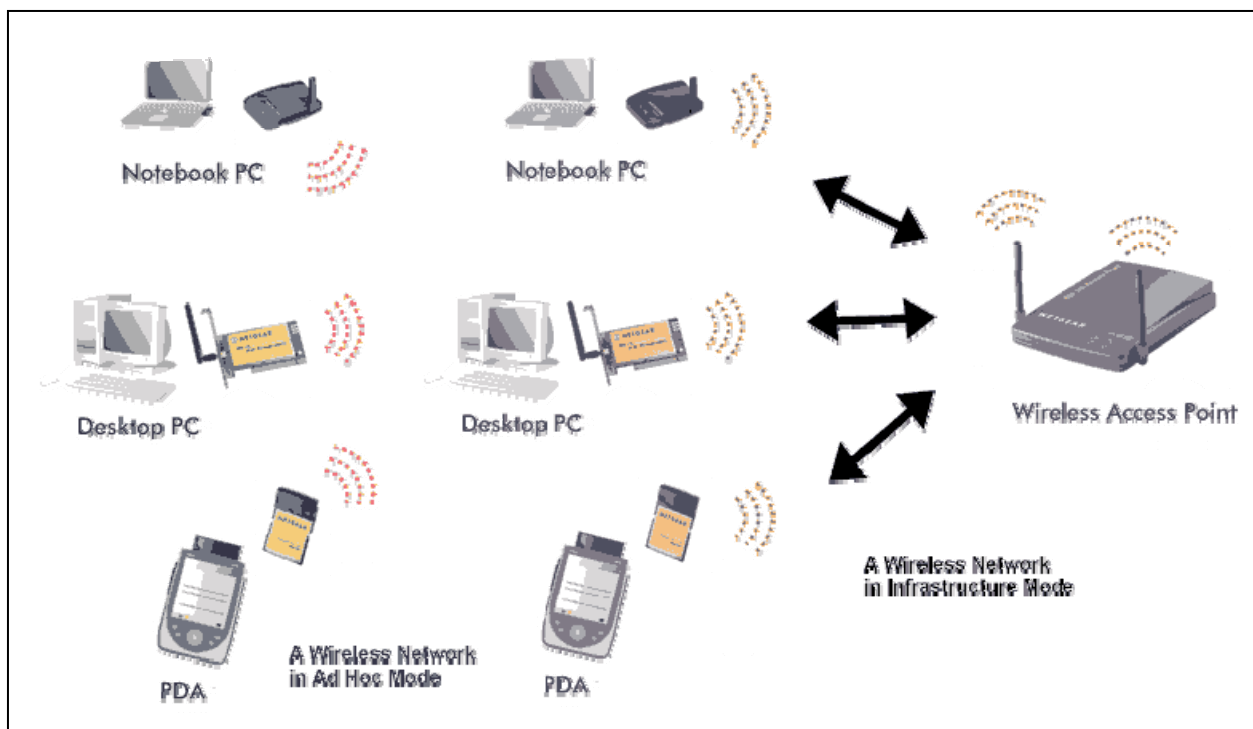


Figure 1: Comparison between wireless networks in ad-hoc and infrastructure modes¹

¹ From: "Selecting Between Infrastructure and Ad Hoc Wireless Modes," http://kbserver.netgear.com/kb_web_files.

Multiple channels² have been utilized in infrastructure-based wireless networks by assigning distinctive channels to adjacent access points. However, typical ad-hoc wireless networks are configured to use only a single channel to ensure connectivity of all their nodes; thus, the aggregate bandwidth provided by the available radio spectrums is not fully-utilized. Such a dilemma has inspired many researchers to propose the exploitation of multiple channels and multiple interfaces to increase the network capacity. Many IEEE 802.11 interfaces³ can be switched from one channel to another, allowing an interface to access multiple channels.

1.1. Project Scope

In this project, we start with extending the Network Simulator (ns-2) to support multiple channels and multiple interfaces by referring to the guideline written by R. A. Calvo and J. P. Campo in [1]. We then explore the use of multiple channels and multiple interfaces in ad-hoc wireless networks by implementing the interface assignment strategy proposed by P. Kyasanur and N. H. Vaidya in [2] using ns-2. Two key benefits of the approach in this paper are that the number of available interfaces can be less than the number of channels since equipping a node with one interface per channel is expensive, and that it can be implemented over existing 802.11 devices without any hardware modifications. Finally, we simulate a simple multi-channel multi-interface ad-hoc wireless network using the modified ns-2 to demonstrate the effectiveness of *interface switching* and the improvement in network throughput.

² A channel is a wireless spectrum with a specified bandwidth.

³ An interface is a network interface card equipped with a half-duplex radio transceiver.

2. Wireless Ad-hoc Network

2.1. Overview

Ad-hoc networks are infrastructure-less. The interconnected nodes coordinate the transmissions with other nodes and they may have to relay messages among several other nodes in order to reach a destination. In other words, an ad-hoc network is a self-configuring network of wireless links connecting mobile nodes, where these nodes may be routers and/or hosts as shown in Figure 2. The decision of which nodes forward data is made dynamically based on the network connectivity. This mechanism is contrary to wired networks in which routers perform the task of forwarding. Ad-hoc networks are also in contrast to the infrastructure-based wireless networks, which require an access point to manage the connection among other nodes.

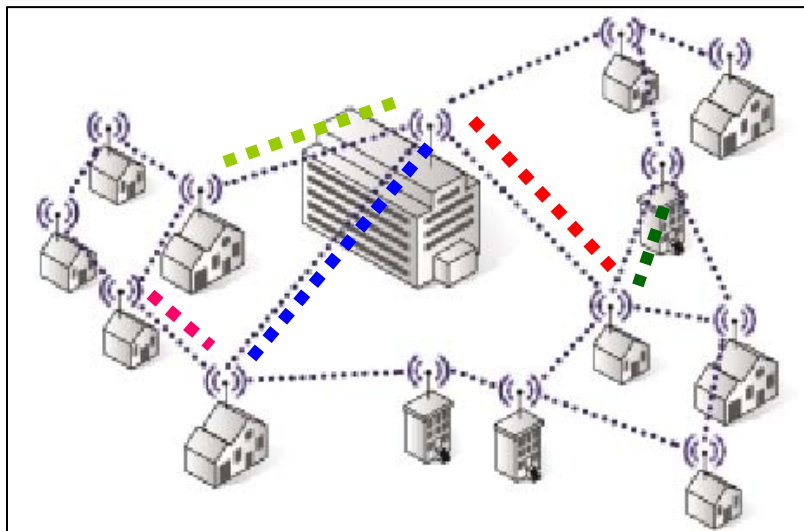


Figure 2: A wireless ad-hoc network

Nodes in a wireless ad-hoc network relies on each other to establish a connection, as there is no central access point which co-ordinates address assignment and routing among the nodes. Thus, a packet can travel from a source to a destination indirectly through some intermediate nodes in the neighbourhood. To maintain a persistent packet forwarding path, a channel is commonly shared with all nodes in the ad-hoc network, due to the fact that two adjacent nodes can only communicate if they are on the same channel. Unfortunately, utilizing one channel among all nodes may eventually lead to traffic congestion as the number of nodes in the network increases.

2.2. Multiple Channels and Multiple Interfaces

IEEE 802.11 offers several non-overlapping channels in the 2.4 GHz and 5 GHz spectrums that are separated in frequency. For example, the 802.11b/g standards operating in the 2.4 GHz bands have 3 non-overlapping channels (1, 6, 11) for use in Canada and the United States. On the other hand, 802.11a operating in 5 GHz offers 12 non-overlapping channels. The large number of available channels in 802.11a is an example in which equipping a node with one interface per channel may not be feasible.

In infrastructure-based wireless networks, multiple channels have already been utilized by assigning different channels to adjacent access points to minimize interference. The reason that most wireless ad-hoc networks today are configured to use a single channel is to avoid the need for node co-ordination. Due to the increasing throughput demand, the idea of exploiting multiple channels is appealing in ad-hoc wireless networks, thus motivating many research communities to develop new algorithms for multi-channel operation.

When multiple wireless channels are available, having more than one interface on a node allows two different nodes communicating in parallel on different channels. An interface has a capability to dynamically switch to different channels over time. Two adjacent nodes can communicate with each other if they have at least one interface on a common channel. In a wireless ad-hoc network consisting of multiple hops, if each node only uses a single interface, packets may be delayed at some hops if their next hops are not on the same channel. The reason is that node A needs to wait after node B's interface has switched to listen to the same channel as node A is using before transmitting a packet to node B, as shown in the figure below.

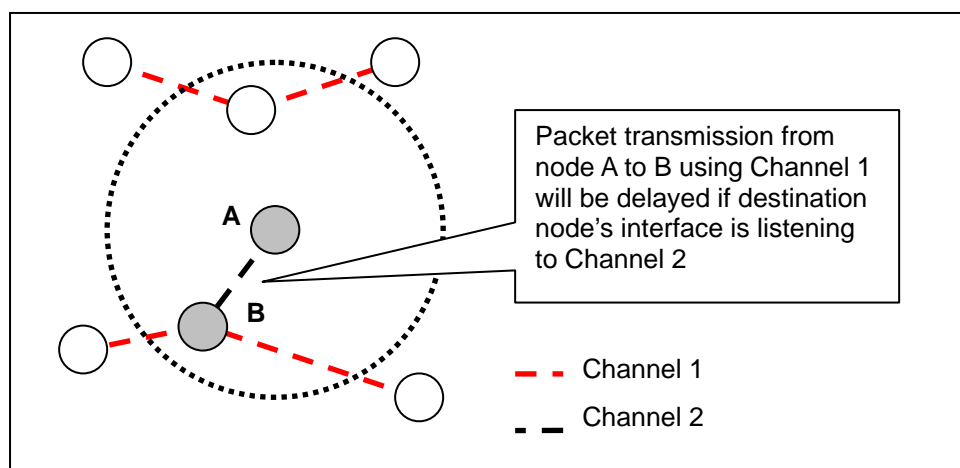


Figure 3: In the presence of only 1 interface, destination node has to listen to the same channel

In other words, the end-to-end delay increases if the traversed hops are on different channels. One might be tempted to solve this problem by assigning each node to be on the same channel. Unfortunately, transmissions on consecutive nodes could interfere with each other, therefore degrading the transmission capacity, as shown in Figure 4.

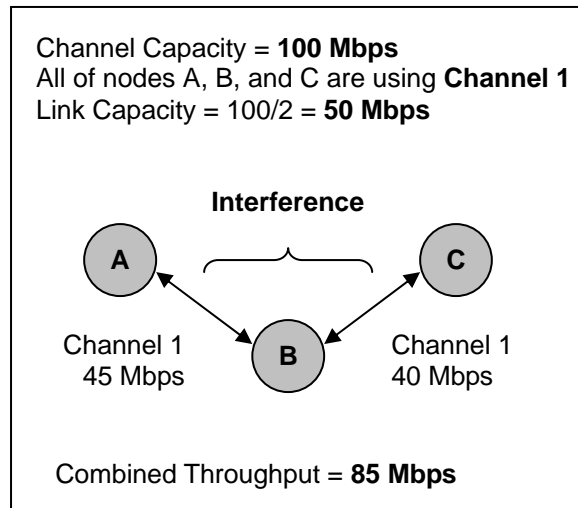


Figure 4: Transmission capacity is degraded when adjacent nodes are on the same channel

Using multiple interfaces, a node is able to transmit and receive data simultaneously. Such an action cannot be carried out when only a single wireless interface is used, as it is half-duplex. When two interfaces and multiple channels are available, we can have one interface transmitting data on one channel while the other interface is receiving data on another channel. The maximum achievable throughput can in turn be nearly doubled as illustrated in the following diagram.

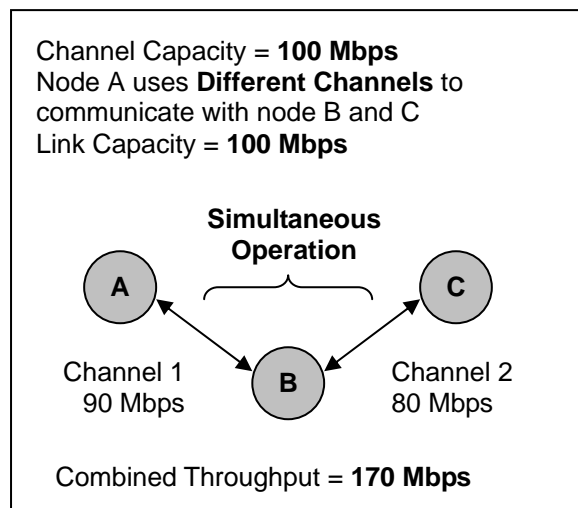


Figure 5: Transmission capacity is doubled when adjacent nodes are on different channels

3. Related Work

Many researchers have proposed the use of multiple channels and multiple interfaces for ad-hoc wireless networks. However, we have selected P. Kyasanur's and N. H. Vaidya's interface assignment scheme in [2] for our implementation, as this approach is more flexible and versatile among others. Below we present a brief comparison between various related work and [2].

For instance, A. Raniwala, K. Gopalan, and T. Chiueh propose the centralized channel assignment and routing solution in [3]. Unlike this solution that is designed for use in static networks where traffic is directed towards specific gateway nodes, the approach in [2] allows any node to communicate with any other. Additionally, the approach in [3] expects stationary nodes and traffic load on every link, since this information is used to assign interfaces and compute routes. Oppositely, the solution in [2] does not need this information and thus is more suitable for ad-hoc networks involving mobile nodes.

Next, S. Wu, C. Lin, Y. Tseng, and J. Sheu propose a MAC layer solution in [4] that requires two interfaces: one for control, therefore assigned to a common channel; the other for data exchange, thus is switched among the remaining channels. Also, R. Maheshwari, H. Gupta, and S. R. Das propose new MAC protocols for multi-channel operation in wireless ad-hoc and mesh networks in [5]. The aforementioned two proposals all require changes to the existing IEEE 802.11 standard. On the contrary, the approach in [2] can be implemented using existing IEEE 802.11 wireless network interface cards.

To address this compatibility issue, P. Bahl, R. Chandra, and J. Dunagan propose Slotted Seeded Channel Hopping (SSCH) in [6], which is a link-layer solution that uses a single interface and can run over unmodified IEEE 802.11. However, SSCH does not support multiple interfaces.

Finally, paper [7] proposes a new routing metric called Weighted Cumulative Expected Transmission Time (WCETT), for multi-channel ad-hoc networks. WCETT ensures channel diverse routes are selected by assuming the number of interfaces per node is equal to the number of channels used by the network. In contrast, the *interface switching* technique proposed in [2] can allow the number of interfaces to be smaller than the number of available channels, while still manages to utilize all the channels.

4. Multi-Channel Multi-Interface Solution

In paper [2], P. Kyasanur and N. H. Vaidya have proposed the following designs for routing and interface assignment in multi-channel multi-interface ad-hoc wireless networks:

1. A multi-interface solution for exploiting multiple channels that can be implemented on existing IEEE 802.11 hardware.
2. An interface assignment strategy using the technique of *interface switching*, that simplifies coordination among nodes while utilizing multiple available channels.
3. A routing protocol, namely the Multiple-Channel Routing (MCR), that selects routes with the highest throughput by accounting for channel diversity and *interface switching* cost.

For the scope of this project, we have limited our implementation in ns-2 to only realize and evaluate the first two goals. Point 1 suggests that existing behaviours of the data-link and physical layers in ns-2 constituting the IEEE 802.11 standard should not be altered. In the next section, we will show how our modifications to ns-2's *MobileNode* to add the support for multiple interfaces preserve the legacy operations of IEEE 802.11 interfaces. We will also incorporate the interface assignment algorithm from point 2; however, since point 3 will not be implemented, point 2 will be added to the existing ad-hoc on-demand distance vector (AODV) routing agent in ns-2 as an extended feature. One shortcoming with AODV is that it selects the shortest-path route (the route that has the least number of hops through other nodes), which may not utilize all the available channels [2].

4.1. Interface Switching

In order to maximize the utilization of all available channels, interfaces have to be switched from one channel to another. A switching protocol is required to assign interfaces to specific channels, ensuring that the neighbour nodes of node X can communicate with itself on-demand. In a sense, all neighbour nodes of node X must at least have knowledge about the channel being used by one channel of X. We describe the proposed algorithm from paper [2] in the following sub-sections.

4.1.1. Fixed and Switchable Interfaces

In a multi-channel multi-interface ad-hoc wireless network, M interfaces available at each node are divided into the two groups:

1. **Fixed Interface:** Some K of the M interfaces at each node are assigned for long intervals of time to some K channels. The corresponding channels are regarded as fixed channels. Fixed interfaces are used to receive data and are to be switched based on the number of nodes using a channel.
2. **Switchable Interface:** The remaining $M-K$ interfaces are dynamically assigned to any of the remaining $M-K$ channels over short time scales based on data traffic. The corresponding channels are designated as switchable channels. A switchable interface enables node X to transmit to node Y in its neighbourhood by switching to the fixed channel used by Y .

The value of M and K can vary in different nodes; however, for simplicity, let's consider a simple scenario in which $M = 2$ and $K = 1$ for all nodes (one fixed and one switchable) and there are three non-overlapping channels. In the following figure, nodes A, B and C have the fixed interfaces on channels 1, 2, and 3 and the switchable interfaces as shown initially. Suppose that the routing path is $A \rightarrow B \rightarrow C$ such that node A wants to send a packet to node C. In the first step, node A switches its switchable interface from 3 to 2 before transmitting the packet to node B since the fixed interface (channel) of node B is 2. Similarly, in the next step, node B switches its switchable interface from 1 to 3 to forward the packet to node C as the fixed interface (channel) of node C is 3. Once the interface has been appropriately set up during a flow initiation, subsequent flow of packets will not need to switch the interfaces for a long time.

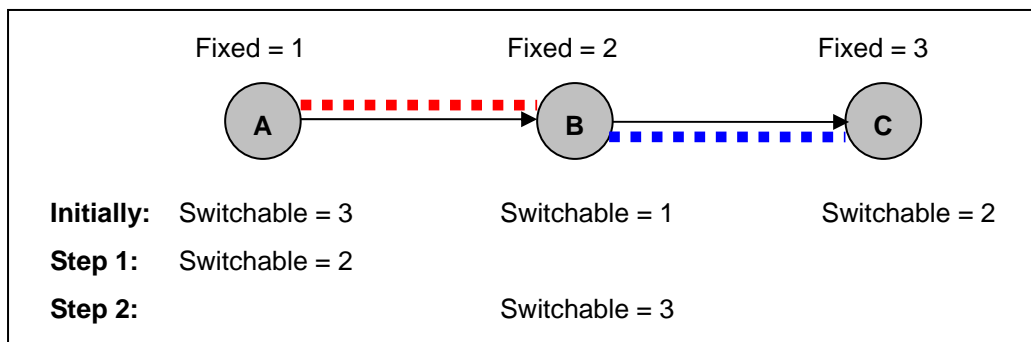


Figure 6: Interface switching with 3 channels and 2 interfaces [2]

In general, when node X has to communicate with node Y over channel c , the switchable interface of X is switched to channel c for communicating with Y . If the fixed channel being used by X is already c , then no switching will be required. In other words, the transmission between two nodes is receiver-based. The sender adapts to the receiver by changing its switchable interface to the receiver's fixed interface, which is the channel used by the receiver.

4.1.2. Fixed Interface Assignment

As illustrated in Figure 6, nodes A, B, and C use channels 1, 2 and 3, respectively; all transmissions directed to A, B, and C will be on channel 1, 2, and 3, respectively. The purposes of the fixed interface assignment are to choose the channel to be assigned to a fixed interface and to inform the neighbour nodes about the channel being used by the fixed interface. First of all, each node maintains two tables:

- *NeighbourTable (NT)*:
 - Contains the fixed channels being used by its neighbours.
- *ChannelUsageList (CUL)*:
 - Keeps the number of nodes using each channel as their fixed channel, but it only tracks the nodes present within its communication range.

Below we present the localized assignment algorithm proposed in [2] along with a visual example of the approach for an ad-hoc wireless network with 3 nodes, each has 3 channels and 2 interfaces (one fixed and switchable). Please note that the node's table we present in each of the following figures is actually an implicit combination of *NeighbourTable* and *ChannelUsageList*.

1. Initially, each node chooses a random channel for its fixed interface.

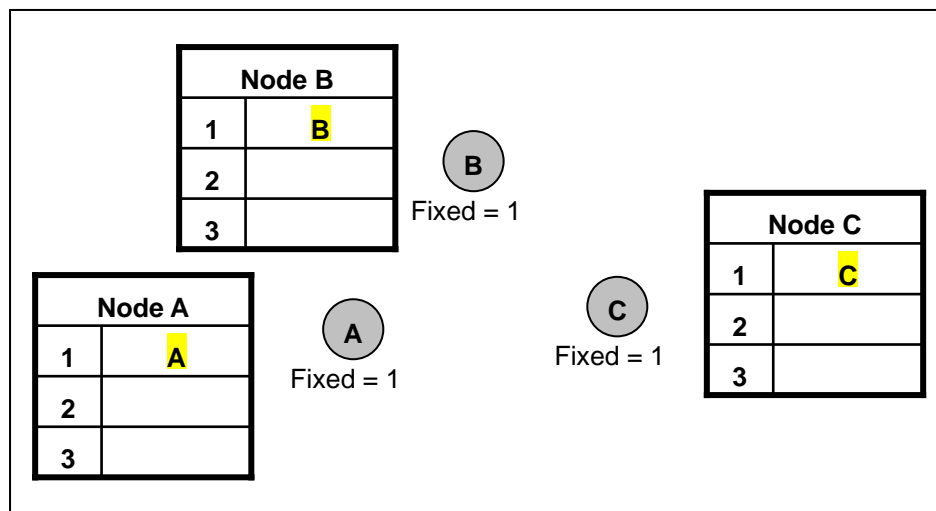


Figure 7: Each node chooses a random channel for its fixed interface

2. Periodically, each node broadcasts a Hello or Route Discovery packet on every channel, which contains the fixed channel being used by the node and its current *NeighbourTable*.
 - a. Many routing protocols such as AODV broadcast a Hello or Route Discovery packet to offer connectivity information.

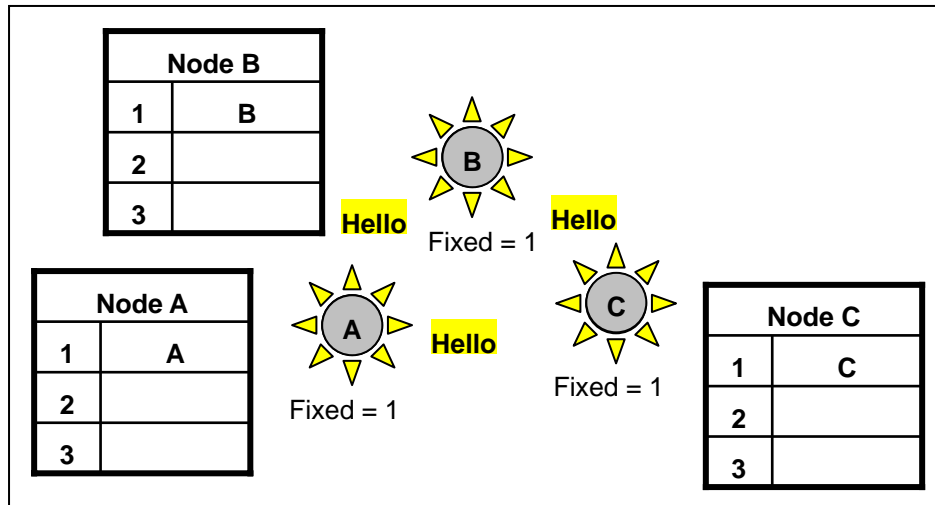


Figure 8: Each node periodically broadcasts on every channel its current fixed channel

3. When a node receives a Hello or Route Discovery packet from a neighbour, it updates its:
 - a. *NeighbourTable* with the fixed channel of that neighbour.
 - b. *ChannelUsageList* using the *NeighbourTable* of its neighbour so *ChannelUsageList* will contain two-hop usage information.

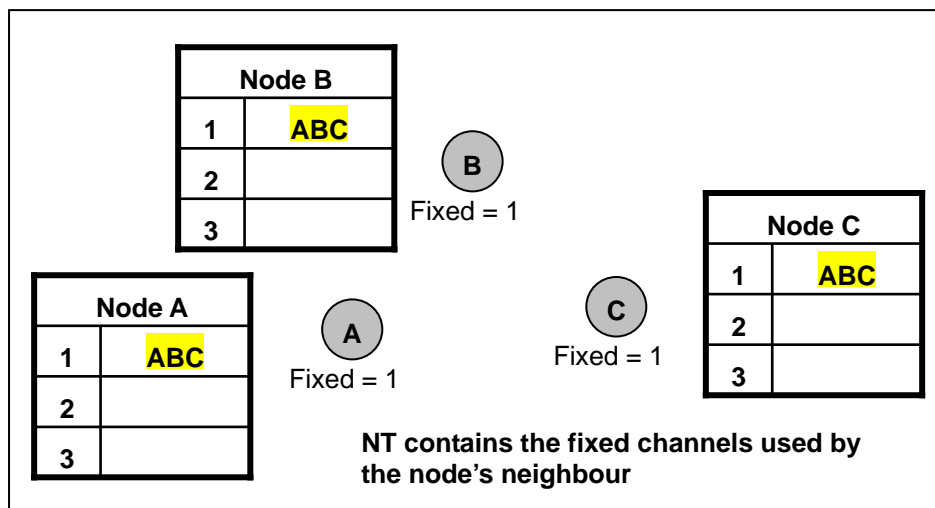


Figure 9: On receiving a Hello packet, a node updates its NeighborTable and ChannelUsageList

4. Each node periodically consults its *ChannelUsageList*. The period chosen is large relatively to packet transmission time.
 - a. If the number of other nodes using the same fixed channel as itself is large, then a node with some probability p changes its fixed channel to a less used channel.
 - i. Then the node transmits a Hello or Route Discovery packet to inform its neighbors of the new fixed channel it is using.

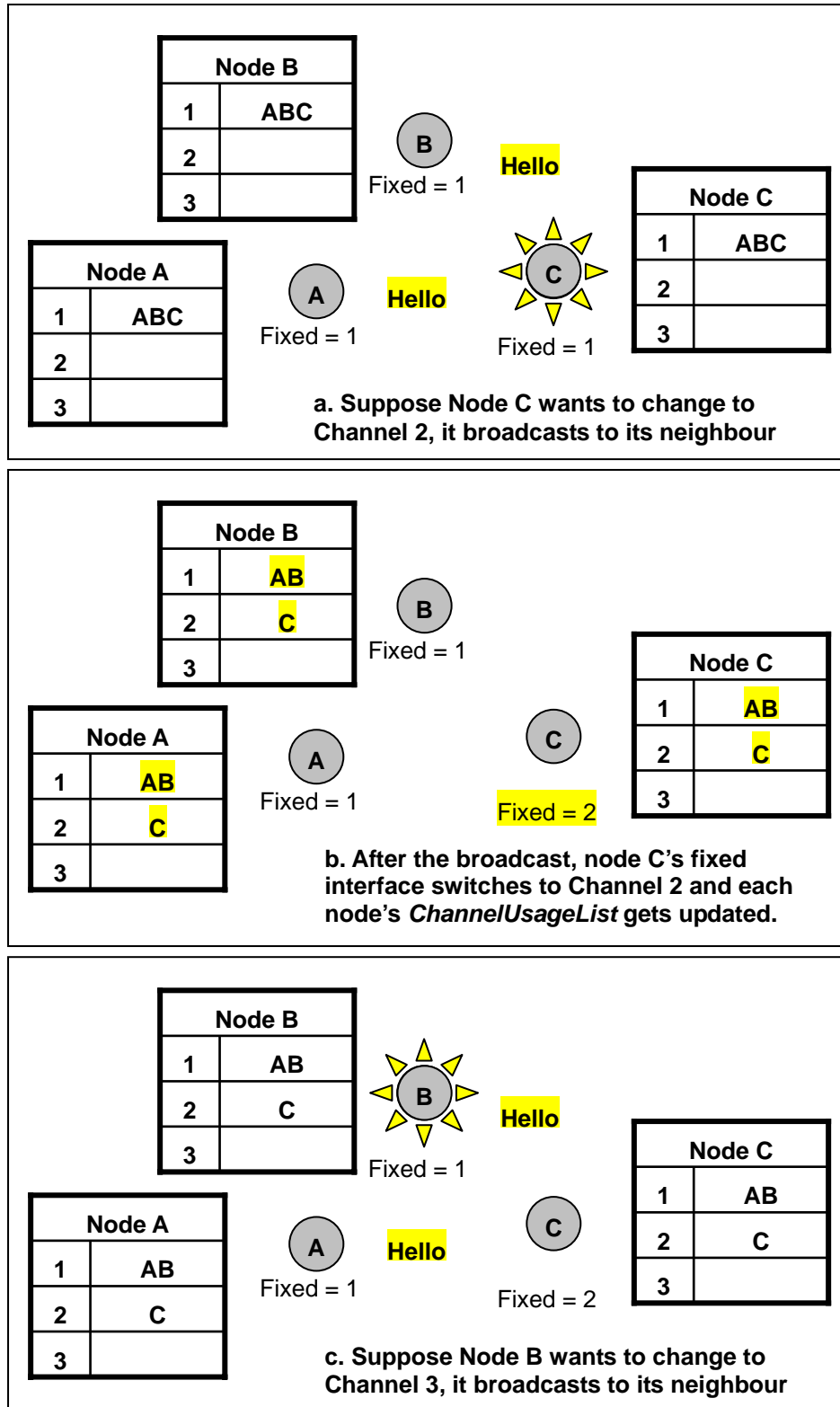


Figure 10: Each node periodically consults its ChannelUsageList

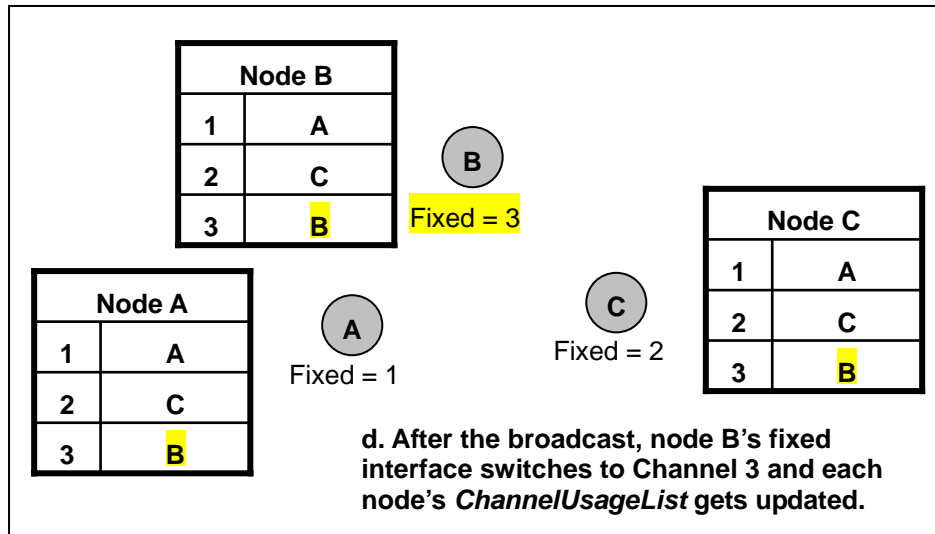


Figure 9: Each node periodically consults its *ChannelUsageList* (continued)

According to paper [2], the period of consulting the *ChannelUsageList* chosen is large relatively to packet transmission time. Also, the probability, p , of switching the fixed interface is set to **0.4**. The enforcement of the switching probability and the decision of switching based on the number of nodes using as channel avoid frequent change of fixed channels. The only scenario that switching will be required is when the network topology changes significantly.

4.1.3. Switchable Interface Assignment

As mentioned earlier, switchable interfaces are used to transmit data from node X whenever the fixed channel of the destination is different from the fixed channel of X. Each node maintains a separate packet queue for each channel as depicted in the figure below:

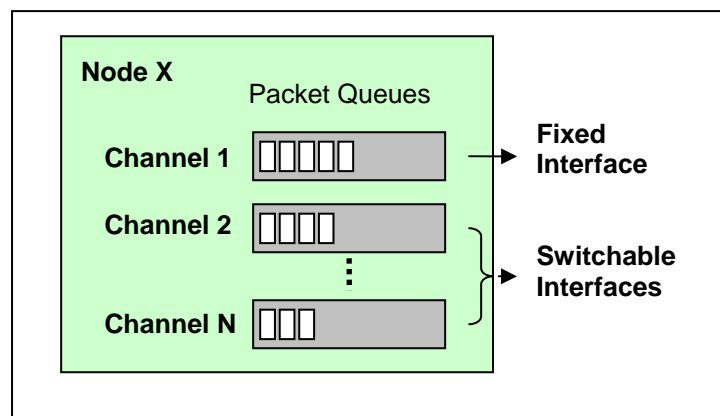


Figure 11: Packet queue maintained by each node for each channel [2]

Below, we present the rules for assigning the switchable interface of a node proposed in [2]:

1. When an unicast packet is received at the link layer for transmission,
 - a. The node looks up the fixed channel of the destination of the packet in the *NeighbourTable*. If the sender has the same fixed channel as the receiver, enqueue the packet to the fixed channel. Otherwise, enqueue to the switchable channel.
 - b. For broadcasting the packet, the node copies it to each channel's queue. The packet will be sent out when that channel is scheduled for transmission.
2. The switchable interface changes channels if there are packets queued for another channel.

5. Adding Multi-Channel Multi-Interface Support to ns-2

The latest version (at the time of the project) of the ns-2 simulator is 2.32; it does not support wireless network simulation involving multiple channels and multiple interfaces. As a result, before attempting to evaluate the exploitation of multiple channels and multiple interfaces in wireless ad-hoc networks, we must add this capability to the current framework of ns-2.

5.1. Overview of ns-2

ns-2 is an open-source discrete event network simulator that supports a variety of different network protocols, producing simulation results for both wired and wireless networks. In this project, we choose ns-2 to implement a multi-channel multi-interface ad-hoc wireless network, primarily due to the flexible extensibility of the simulator. ns-2 is developed in C++ but it also provides a simulation interface through an object-oriented tool command language (OTCL).

When performing a network simulation, users first write an OTCL script to describe the relevant network topology, and to define traffic sources and when to start and stop transmitting packets through an event scheduler. ns-2 has an OTCL interpreter that translates the event scheduler and the network component OTCL objects and member functions into their corresponding C++ counterparts. Such a mechanism provides a linkage between the OTCL and C++ realms, making the control functions and the configurable variables interoperable. Next, ns-2's main program simulates the topology with user-specified parameters. Eventually, the simulation is finished; ns-2 generates an output trace file that contains detailed simulation data. The data can be either post-processed to create simulation graphs or input to a graphical simulation tool called Network Animator (nam) for later viewing. The architecture of ns-2 is shown in the figure below.

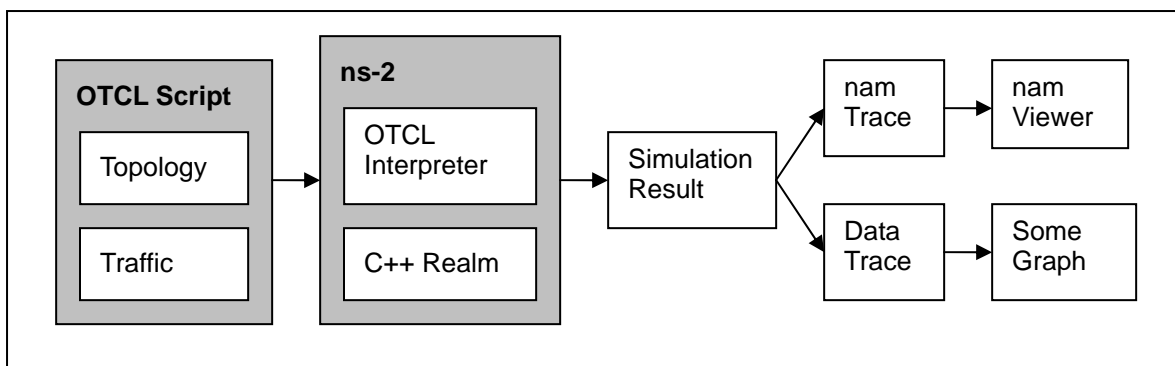


Figure 12: Overall Architecture of ns-2

5.2. Multiple Interface Node Model

The reason that the guideline in [1] is chosen for our implementation is that the multiple interface node model it proposes offer the following advantages which are suitable for the approach in [2]:

- The number of channels is modifiable.
- The number of interfaces per node is variable; it needs not to be the same for all nodes.
- Each node can be connected to different number of channels.
- Legacy operations of ns-2 are still preserved. This backward compatibility is particularly important as we will use the modified simulator to simulate wireless ad-hoc networks without the exploitation of multiple channels and multiple interfaces.

As mentioned by the guideline, extending ns-2 to support multiple interfaces requires modifying its existing *MobileNode* model whose architecture is depicted in Figure 13. *MobileNode* is an extension of the *Node* object in ns-2 with additional functionalities such as mobility, ability to transmit and receive on a channel that allows it to be used in mobile and wireless network simulations. Some main component objects that a *MobileNode* object is consisted of are:

- **Routing Agent**
 - Routes data packets to the next-hop node on the *MobileNode*'s behalf. Currently, ns-2 supports four ad-hoc routing protocols as follows:
 - Ad-hoc On-demand Distance Vector (AODV)
 - Destination-Sequenced Distance-Vector Routing (DSDV)
 - Dynamic Source Routing (DSR)
 - Temporally Ordered Routing Algorithm (TORA)
- **Link Layer**
 - Sets the MAC destination address in the MAC packet header by finding the IP address of the next-hop-node directed by the routing agent and resolving this address into the corresponding MAC address by ARP.
- **Address Resolution Protocol (ARP)**
 - Receives queries from the link layer to resolve the associated hardware address by referring to an ARP table. If the address can be found locally, it is written to the MAC header of the packet; otherwise, an ARP query is broadcasted. Once the MAC address is retrieved, the packet is inserted into the interface queue.
- **Interface Queue**
 - Gives priority to the stored routing protocol packets.
- **Media Access Control (MAC)**
 - Processes data packets received from or to be sent to the link layer. When

sending, it adds the MAC header and transmits the packet onto the channel. Or, it asynchronously receives packets from the classifier of the physical layer.

- **Network Interface**
 - Serves as a hardware interface for *MobileNode* to access the channel.
- **Channel**
 - Simulates the effect of the real wireless channel on the transmitted signal.

Each of the objects emulates a real-life entity in the physical and data-link layers of a wireless network. Together, these components allow simulation of WLANs or multi-hop ad-hoc networks.

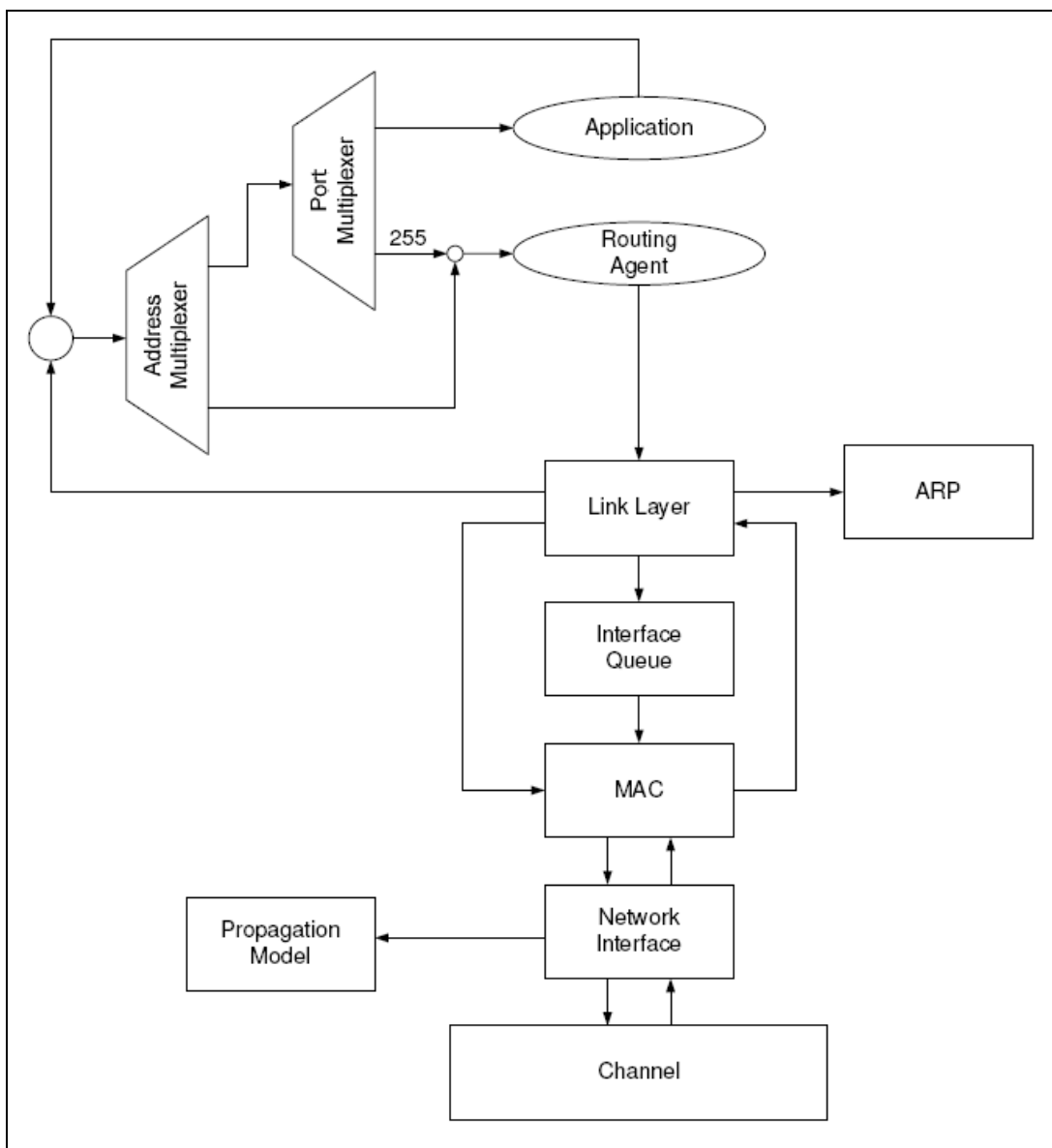


Figure 13: Original MobileNode Architecture in ns-2 [1]

Figure 14 presents a high-level architecture of the modified *MobileNode* object proposed in [1]. Each node can have as many instances of the link layer, ARP, interface queue, MAC, network interface and channel entities as the number of interfaces it has. One can imagine that each instance actually represents a wireless network interface. Thus, this design scheme emulates the fact that our multi-channel multi-interface ad-hoc network implementation will not require any modification to existing IEEE 802.11 hardware.

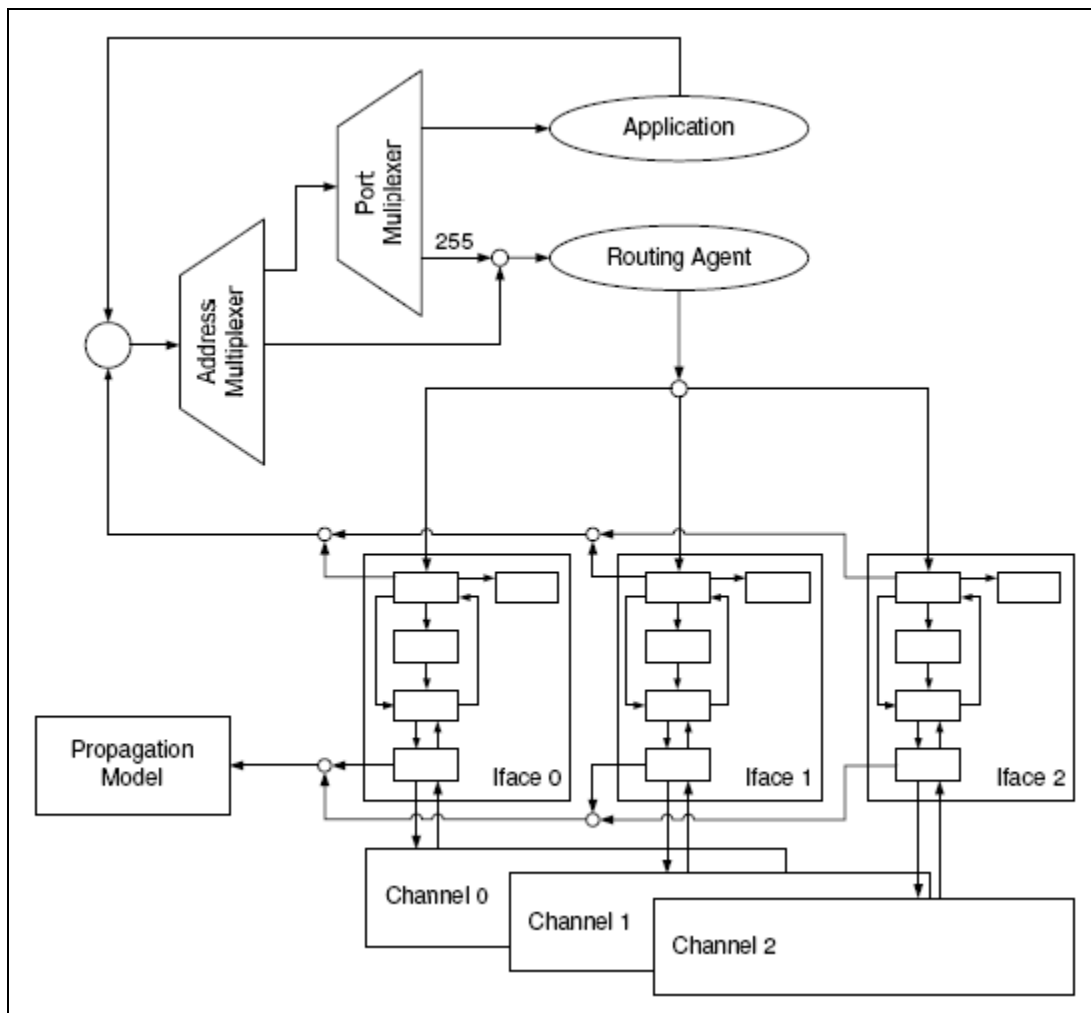


Figure 14: Modified MobileNode architecture, with the support for multiple interfaces [1]

As can be observed, most legacy operations of ns-2 are still preserved. Incoming traffic arrives through the corresponding channel and travel through the different components in ascending order then eventually merges to a single point at the address multiplexer. For outgoing traffic, the determination of selecting which interface to pass the data packets is to be handled by the routing agent. In other words, modifications will be required in implementing the routing agent to add the intelligence of selecting the appropriate interface, as will be discussed in Section 5.4.

5.3. Extending ns-2 with Multi-Interface Support

This section provides a description of the various changes that the guideline in [1] proposes to add the multiple interface support in ns-2. As can be deduced from the ns-2 architecture mentioned in earlier sections, the modifications involves both the OTCL and C++ realms. The changed lines of code in each modified file can be found in the Appendix section, marked with a commented tag, “MCMI”.

5.3.1. Changes on OTCL Code

Four new OTCL procedures as follows are created in `/ns-2.32/tcl/lib/ns-lib.tcl`:

- **change-numifs** {*newnumifs*}
 - Allows users to specify *newnumifs* interfaces per node. This procedure is to be called prior to creating a wireless node in the scenario script. Once called, it will affect all subsequent nodes until another invocation of the procedure is issued.
- **add-channel** {*indexch ch*}
 - Adds an interface (channel) to a node. This procedure takes two arguments: *indexch* is the index of the channel to be added; *ch* references to the channel object previously created.
- **get-numifs** {}
 - Retrieves the number of interfaces currently defined.
- **ifNum** {*val*}
 - Adds multiple interfaces as an argument to `node-config`, which is an existing ns-2 command used to configure a *MobileNode* object, by setting a value for the local variable, `numifs_`.

Secondly, two of the existing OTCL procedures in `/ns-2.32/tcl/lib/ns-lib.tcl` are modified:

- **node-config args** {}
 - Adds the support for multiple channels.
 - For backward compatibility, initializes `chan` as a single variable if normal operation is used, or as an array if multiple interfaces are defined.
 - Adds the `numifs_` variable as a new member in the argument list, `args`, that is passed to the procedure.
- **create-wireless-node** {}
 - Takes in the number of interfaces specified in `numifs_` and iteratively calls

`add-interface` as many times as the number of interfaces that the node has.

- `add-interface` is an existing ns-2 procedure that adds an interface to a previously created *MobileNode* object.

Next, several modifications are performed on the following existing ns-2 procedures in `/ns-2.32/tcl/lib/ns-mobilenode.tcl`:

- **`add-target {agent port}`**
 - This procedure attaches a routing agent to a *MobileNode* object, picks a port and binds the agent to the port number.
 - Gets the number of interfaces via calling `get-numifs`; doing so allows the procedure to determine whether multiple interfaces are present. If the number of interfaces is non-zero (implying that the multiple interface extension is used), the procedure iteratively calls the `if-queue` command as many times as the value returned by `get-numifs`.
- **`add-target-rtagent {agent port}`**
 - This procedure, called by `add-target`, adds a target routing agent.
 - It first gets the number of interfaces that a node has from `get-numifs`. If the number of interfaces is non-zero, the procedure associates the routing agent with the corresponding link layer `target` entity as many times as the number of interfaces for both of the sending and receiving targets.
- **`add-interface {}`**
 - This procedure adds an interface to a *MobileNode* object.
 - Originally, it creates one ARP table (for address resolution) per node. We modify it to be one ARP table per interface although having one ARP table per node resembles the real-life case. The reason for such a walk-around is that, if a node is using one interface to communicate with another one, the current design of *MobileNode* in ns-2 will not allow the node to use another interface since the request to the ARP entity will still be serving the previous interface.
- **`init args {}` and `reset {}`**
 - Due to the above change of assigning the ARP table, these two procedures are modified to initialize and reset the ARP table of a *MobileNode* object per the number of interfaces defined.

5.3.2. Changes on C++ Code

After adding the multiple interface support in the OTCL realm, relevant changes have to be made in the C++ code for the *MobileNode* object, the channel entity, and the MAC layer model.

/ns-2.32/common/mobilenode.h

- We define a new declaration of the *MobileNode* lists to replace the existing ones.
 - ns-2 controls each instance of the *MobileNode* objects which are associated with a channel by means of a linked-list. Two lists are managed; one references the previous node, `prevX_`, while the other references the next node, `nextX_`.
 - The original format of the list is simply a pointer to a node. In order to support multiple channels, the list is modified to be an array of pointers with the size of the array being the maximum of number of channels:
 - `nextX_[MAX_CHANNELS]`
 - `prevX_[MAX_CHANNELS]`
 - The index for referring to a node is the channel number.
- We remove the inline declaration of the `getLoc()` function. Due to the above changes on the *MobileNode* lists, the original declaration has been found to always return a zero distance, which leads to wrong packet receptions.

/ns-2.32/common/mobilenode.cc

- We add the `getLoc()` method definition which retrieves the location of a node.

/ns-2.32/mac/channel.cc

- Due to the changes on the *MobileNode* lists, we modify accessing each node entry when attaching, removing, and updating a new node to a channel to refer to the corresponding channel number. This number can be accessed by `this->index()`, where `this` is the current instance of the channel class object. In other words, whenever `nextX_` and `prevX_` appear in **channel.cc**, they need to be replaced by:
 - `nextX_[this->index()]`
 - `prevX_[this->index()]`
- In the `affectedNodes()` function, we add a new condition to check which of the interfaces of the destination node is connected to the same channel before transmitting the packet to another interface. The original design of ns-2 does not consider the case with multiple channels; it simply sends the packet to all of the destination interfaces. Accessing the channel of the destination interface is carried out by:

- o `rifp->channel()`, where `rifp` is a pointer to the receiving interface. The `channel()` member returns a channel object that has the same type as `this`.

/ns-2.32/mac/mac-802_11.cc

- For correct handling of multiple interfaces by the routing agent, we modify the `recv()` method in the MAC 802.11 class to register the correct MAC receiving interface in the MAC header. The hardware address of the interface can be access by:
 - o `hdr->iface() = addr()`

5.4. Changes on the AODV Routing Protocol Code

As mentioned in an earlier section, the MCR protocol proposed in [2] is not implemented. For simplicity, we have chosen the exiting AODV routing protocol for multi-hop networks in ns-2. AODV is a reactive routing protocol, defined in RFC 3561, for ad-hoc wireless networks, as it establishes a route to a destination only when required [8]. Although the RFC defines the AODV protocol to be capable of supporting multiple wireless interfaces at each node, the original ns-2 design does not incorporate this capability. Hence, the enhancements consist of two parts:

- Support to make use of the multi-interface model (the modified *MobileNode* by [1]).
- Implementation of the *interface switching* algorithm in [2].

First of all, we must understand how routing is performed in ns-2’s AODV agent. As an example, we examine how AODV establish a route from the first to the last node in a chain topology with four nodes. AODV builds routes using a Route Request (RREQ)/Route Reply (RREP) query cycle, as illustrated in the figure below.

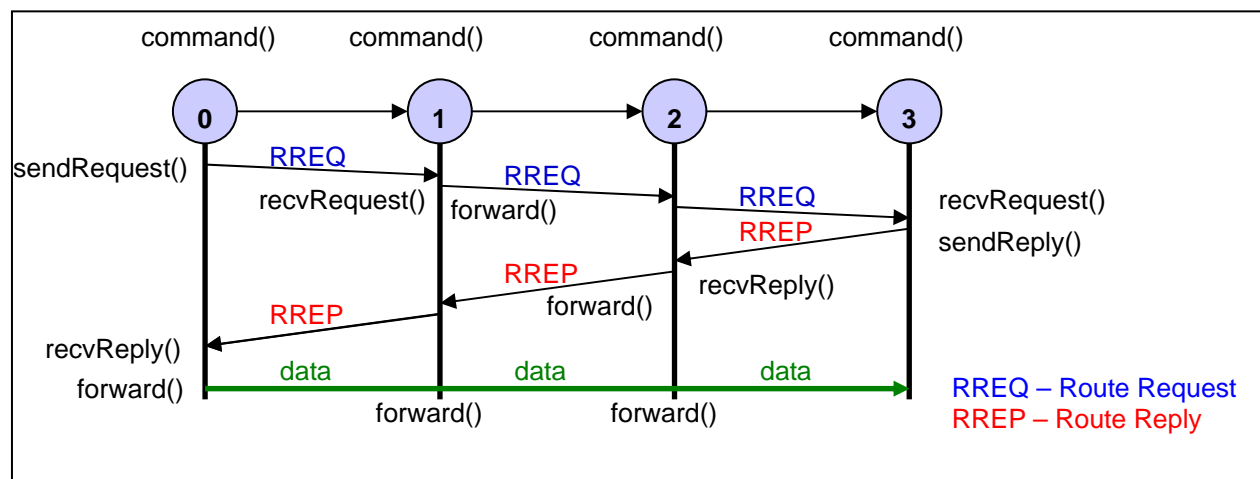


Figure 15: AODV route establishment for a chain topology with 4 nodes

When the nodes in an ad-hoc wireless network are created, function `command()` in ns-2's AODV library is executed by each node to initialize itself in the AODV routing agent. When a source desires a route to a destination node, it broadcasts a RREQ packet across the network by calling `sendRequest()`. Intermediate nodes receiving this packet update their information for the source node and set up a backward pointer to the source node in their routing tables. The packet is also forwarded to the next node via `forward()`. Once a node receiving the RREQ is the desired destination, it unicasts a RREP packet back to the source by calling `sendReply()`. As the RREP propagates back to the source node, the intermediate nodes receiving the packet via `recvReply()` set up a forward pointer to the destination and forward it to the previous node by `forward()`. As soon as the source node receives the RREP, an active route has been established; it may start to forward data packets to the destination by calling `forward()`. The aforementioned steps are known as the Route Discovery state.

Next, the Route Maintenance state (not illustrated in Figure 15 for lack of space) is for updating the route previously established. Occasionally, a node of an active route may offer connectivity information by calling `sendHello()` to broadcast local Hello messages. Whenever a node receives a Hello message from a neighbor via `recvHello()`, the node makes sure that it has an active route to the neighbor, and creates one if necessary. The current node can now begin using this route to forward data packets. Furthermore, when a link breakage is detected by a node, the node will call `sendError()` to notify the source node with a Route Error (RERR) packet. The source node will initialize a new Route Discovery state as mentioned before.

5.4.1. Support to Make Use of the Modified MobileNode

The goal is to utilize the multi-interface model provided by the modified *MobileNode*. The relevant files to be modified are described in the following paragraphs.

/ns-2.32/aodv/aodv.h

- We add a new member to the AODV routing agent class, `nIfaces`, which keeps track of the number of interfaces that the agent is managing.
- We define a constant, `MAX_IF = 12`, for the maximum number of allowable interfaces.
- The routing agent needs to decide which one of the interfaces the outgoing packets should be routed to. Since there are multiple interfaces. The originally used `ifqueue` and `target` pointers need to be modified to two arrays, `ifqueuelist` and `targetlist`. In a sense, these two containers resemble the packet queues maintained by each node, as mentioned in [2].

- o `NSObject *targetlist[MAX_IF]`
 - `ifqueuelist` stores the link-layer modules for all the interfaces a particular node has.
- o `PriQueue *ifqueuelist[MAX_IF]`
 - `targetlist` keeps the corresponding queue of each individual node.

/ns-2.32/aodv/aodv.cc

- We initialize `nIfaces` to be 0 in the constructor of the AODV agent.
- We modify the `command()` function of the AODV routing agent to make use of the various variables from the following OTCL procedure calls in **/ns-2.32/tcl/lib/ns-mobilenode.tcl**, and to store the corresponding entries in `ifqueuelist` and `targetlist`. Note that all of the OTCL statements highlighted in bold below have four arguments and begin with the routing agent object, `$agent`. Thus, we add a special case in `command()` to deal with the condition in which four parameters are present. The second, the third and the fourth parameters are of our interest; their meanings and respective actions to be carried out in the `command()` function are also described below.

- o **add-target**

- We have previously added the following lines in `add-target`:


```
for {set i 0} {$i < [$self set nifs_]} {incr i} {
    $agent if-queue $i [$self set ifq_($i)]
}
```
- `if-queue` is a label signifying that the interface queue instance, `ifq_($i)`, should be inserted to index `$i` of `ifqueuelist`.

- o **add-target-rtagent**

- We have also added the following lines in `add-target-rtagent`:


```
for {set i 0} {$i < [$self set nifs_]} {incr i} {
    set sndT [cmu-trace Send "RTR" $self]
    $agent target $i $sndT
    $sndT target [$self set ll_($i)]
}
for {set i 0} {$i < [$self set nifs_]} {incr i} {
    $agent target $i [$self set ll_($i)]
}
```
- `target` is a label signifying that the link layer instance, `ll_($i)`, should be inserted to index `$i` of `targetlist`.

Next, we must add some coupling mechanism within the routing agent implementation so that it has a path to which of the interfaces to transmit the packet. In ad-hoc network routing, broadcast transmission is used during the Route Discovery process. In the presence of multiple interfaces, a broadcast packet, RREQ, needs to be sent through all the interfaces that a node has. As a result, we use a for loop to send a copy of the original packet to each interface whenever the AODV routing agent needs to broadcast a packet, which happens in the following methods:

- `sendRequest()`
- `sendError()`
- `sendHello()`
- `forward()`

Please note that, in order to preserve the original behaviours of ns-2, the above changes are only performed when `nIfaces` is non-zero.

Additionally, for routing the packet via a specific interface to a destination, unicast transmission of AODV is exercised. Just knowing the next node is not enough; the routing agent must also consider which output interface to be used to reach the next hop. In the AODV routing table header file, `/ns-2.32/aodv/aodv_rtable.h`, we define a new member, `rt_interface`, which stores the appropriate interface index. `rt_interface` is used to index an entry in `targetlist` within the following functions for the sake of selecting the intended link-layer entity of the interface that the packet needs to be sent to:

- `sendReply()`
- `forward()`

However, the challenge here is how to associate the appropriate output interface with the `rt_interface` index when a new routing table entry is created. This implies that the respective `rt_interface` needs to be updated so as to be able to keep this index. Thanks to the changes in the `/ns-2.32/mac/mac-802_11.cc` file, the address of the incoming interface is one of the members in the common header of the packet; that is, `ch->iface()`. In the following two functions, we have `rt_interface` to hold the index that will be used for creating and/or updating the route table entry:

- `recvRequest()`
- `recvReply()`

Such an index is determined by the following formula:

```
rt_interface = ch->iface()
               - ((Mac *)ifqueuelist[0]->target())->addr(),
```

where the second term of the subtraction is address of the first interface of the node. If multiple interfaces are not present, we assign `rt_interface` to be -1.

5.4.2. Implementation of Interface Switching

In addition to the various changes in ns-2's AODV that makes use of the multi-interface model from [1], the *interface switching* algorithm from [2] is required to assign interfaces to specific channels, as well as to decide when to switch an interface from one channel to another. We also employ a flag, `MCMI_DEBUG`, for displaying informative debugging messages within each modified function. When building the modified ns-2 with this flag defined, debugging messages will be displayed while a simulation is running; they are useful for validating the functionality of *interface switching* (to be discussed in Section 6.2).

First of all, the following new member variables are created in ns-2's AODV class,

/ns-2.32/aodv/aodv.h:

- **MAX_NT_CUL_ENTRIES:**
 - Maximum number of entries in *NeighbourTable* and *ChannelUsageList*.
- **neighbour_table[MAX_NT_CUL_ENTRIES]:**
 - Contains the fixed channels used by the node's neighbours, indexed by the node IP address.
- **channel_usage_list[MAX_NT_CUL_ENTRIES]:**
 - Contains the number of nodes using each channel as their fixed channel, indexed by the channel number.
- **fixed_interface, switchable_interface**
 - Fixed and switchable interfaces used by this node.

Next, to facilitate the algorithm of *interface switching*, the RREQ, RREP and Hello packets used by the AODV agent during the Route Discovery and Route Maintenance processes need to contain some additional information about a node, such as its *NeighbourTable*. Hello packets share the same header structure as that of RREP in ns-2's AODV class. The following new variables are added to the `hdr_aodv_request` (for RREQ) and `hdr_aodv_reply` (for RREP and Hello) structures in **/ns-2.32/aodv/aodv_packet.h:**

- **rq_fixed_channel_used/rp_fixed_channel_used**
 - Fixed channel used by this node.
- **rq_sender_node_ip/rp_sender_node_ip**
 - Node IP address of the sender that transmits the RREQ/RREP/Hello packet. It

may not be the originator where the packet is originally from.

- ***rq_neighbour_table/*rp_neighbour_table**
 - A pointer to the *NeighbourTable* of this node.

Finally, the various changes in the relevant functions in ns-2's main AODV class code, **/ns-2.32/aodv/aodv.cc**, are briefly described below. For the completed changes, please refer to the Appendix section. Again, in order to preserve the original behaviours of ns-2, the changes are only performed when `nIfaces` is non-zero.

command()

- Initially, the node chooses a random channel for its fixed interface and switchable interface. The following lines of code needs to be performed every time a channel is created by a scenario script and only if the number of interfaces is greater than 1, as the node does not know how many channels it is going to have and there is no need to assignment another interface if the scenario script only intend to simulate an one-interface network. The corresponding lines of code are as follows:

```

    o fixed_interface = rand() % (temp_+1);
      if (nIfaces > 1) {
        do {
          switchable_interface = rand() % (temp_+1);
        } while (switchable_interface == fixed_interface);
      }

```

- Next, the node adds the fixed channel used by itself to its *NeighbourTable* and updates the node's *ChannelUsageList* with its fixed channel. The fixed channel is retrieved by accessing the corresponding channel index of the node's current fixed interface.

```

    o fixed_channel =
      ((Mac *)ifqueuelist[fixed_interface]->target())->netif()
      ->channel()->index();
      neighbour_table[(int)index] = fixed_channel;
      channel_usage_list[fixed_channel]++;

```

sendRequest(), sendReply(), sendHello()

- These function need to add the fixed channel used by this node and its *NeighbourTable* to the outgoing RREQ, RREP, or Hello packet.

```

    o rq->rq_fixed_channel_used = neighbour_table[(int)index];
      rq->rq_sender_node_ip = (int)index;

```

```
rq->rq_neighbour_table = &neighbour_table[0];
```

recvRequest(), recvReply(), recvHello()

- When a node receives a RREQ, RREP, or Hello packet from a neighbour, it updates:
 - The node's *NeighbourTable* with the fixed channel of that neighbour.
 - `neighbour_table[(int)(rq->rq_sender_node_ip)] = rq->rq_fixed_channel_used;`
 - The node's *ChannelUsageList* using the *NeighbourTable* of its neighbour. Doing so ensures the *ChannelUsageList* will contain two-hop channel usage information.
 - ```
for (int i = 0; i < MAX_NT_CUL_ENTRIES; i++) {
 if (i == (int)index) {
 continue;
 }
 if (rq->rq_neighbour_table[i] != -1) {
 channel_usage_list[rq->rq_neighbour_table[i]]++;
 }
}
```
- For `recvReply()` and `recvHello()`, the lines of code are symmetric except for the names of the packet header pointer are `rp` and `rh`, respectively, and the first two letters of the accessed member variables are `rp` in both cases.

### **forward()**

- When forwarding a RREQ, RREP or Hello packet received by the node from its neighbour, the function adds the fixed channel used by this node and its *NeighbourTable* to the outgoing packet, as done by the following lines of code:
  - ```
if (rq->rq_type == AODVTYPE_RREQ) {
    rq->rq_fixed_channel_used = neighbour_table[(int)index];
    rq->rq_sender_node_ip = (int)index;
    rq->rq_neighbour_table = &neighbour_table[0];
}
```
 - ```
if (rp->rp_type == AODVTYPE_RREP
 || rp->rp_type == AODVTYPE_HELLO) {
 rp->rp_fixed_channel_used = neighbour_table[(int)index];
 rp->rp_sender_node_ip = (int)index;
 rp->rp_neighbour_table = &neighbour_table[0];
}
```



- When forwarding a data packet to a node's neighbour:
  - The function consults the node's *ChannelUsageList* to find the current channel with the largest and the lowest usage.
    - ```
for (int i = 0; i < nIfaces; i++) {
    if (channel_usage_list[i] == 0) {
        channel_lowest_usage = i;
        continue;
    }
    if (channel_usage_list[i] >=
        channel_usage_list[channel_largest_usage]) {
        channel_largest_usage = i;
    }
    if (channel_usage_list[i] <=
        channel_usage_list[channel_lowest_usage]) {
        channel_lowest_usage = i;
    }
}
```
 - If the node's fixed channel has the largest usage, with a probability of 0.4 (from paper [2]), the node reverses its *ChannelUsageList* about the fixed channel previously used changes its fixed channel to a less used channel. The node then transmits a new Hello packet informing neighbours of its new fixed channel by calling `sendHello()`.
 - ```
if (channel_usage_list[neighbour_table[(int)index]] ==
 channel_usage_list[channel_largest_usage]) {
 int switch_probability = rand() % 101;

 if (switch_probability <= 40) {
 int fixed_channel =
 ((Mac *)ifqueuelist[fixed_interface]->target())
 ->netif()->channel()->index();
 channel_usage_list[fixed_channel]--;

 fixed_channel = channel_lowest_usage;
 channel_usage_list[fixed_channel]++;
 neighbour_table[(int)index] = fixed_channel;
 }
}
```

```

 for (int i = 0; i < nIfaces; i++) {
 if (((Mac *)ifqueuelist[i]->target()->
netif()->channel()->index() == fixed_channel) {
 fixed_interface = i;
 break;
 }
 }
 sendHello();
 }
}

```

- o If the usage of the node’s fixed channel is ok, the node looks up the fixed channel of the next node in its *NeighbourTable* and assigns this fixed channel to the its switchable interface.

```

 int switchable_channel =
neighbour_table[(int)(rt->rt_nexthop)];
 for (int i = 0; i < MAX_IF; i++) {
 if (((Mac *)ifqueuelist[i]->target()->netif()
->channel()->index()) == switchable_channel) {
 switchable_interface = i;
 break;
 }
 }
}

```

### 5.5. Modified ns-2 with Multi-Interface and Interface Switching Capabilities

In summary, the following chart provides a hierarchical view of the various modified files in ns-2.32 as discussed in earlier sub-sections of Section 5.

| Interface Switching                                                                      | Multi-Interface Support                                              |                                                                           |                                                                             |
|------------------------------------------------------------------------------------------|----------------------------------------------------------------------|---------------------------------------------------------------------------|-----------------------------------------------------------------------------|
| <b>/ns-2.32</b><br><b>/aodv</b><br>/aodv.cc<br>/aodv<br>/aodv_packet.h<br>/aodv_rtable.h | <b>/ns-2.32</b><br><b>/common</b><br>/mobilenode.cc<br>/mobilenode.h | <b>/ns-2.32</b><br><b>/mac</b><br>/channel.cc<br>/mac.h<br>/mac-802_11.cc | <b>/ns-2.32</b><br><b>/tcl</b><br>/lib<br>/ns-lib.tcl<br>/ns-mobilenode.tcl |

**Figure 16: Modified files ns-2 with multi-interface and interface switching capabilities**

## 6. Results and Discussions

We have simulated a single-route multi-channel multi-interface ad-hoc wireless network using the modified ns-2 simulator that incorporates the support for multiple channels and multiple interfaces as well as the approach of *interface switching*. The interface assignment algorithm does not require any modifications to IEEE 802.11's data-link and physical layers. We wish to evaluate the effectiveness of the *interface switching* algorithm implemented in ns-2's existing AODV routing agent, as well as the throughput performance between ad-hoc wireless networks with and without the exploitation of multiple channels and multiple interfaces. Multi-path routing is not tested as the MCR protocol proposed by [2] is not implemented.

### 6.1. Simulation Scenario

To demonstrate the performance of our implemented *interface switching* algorithm in ns-2 and the benefits of using multiple channels and multiple interfaces, ad-hoc wireless networks in simple chain topologies are used for simplicity. The reasons are that only one single route between the source and the destination is involved and that direct communication is only possible between adjacent nodes on the chain, as illustrated in the following block diagram.

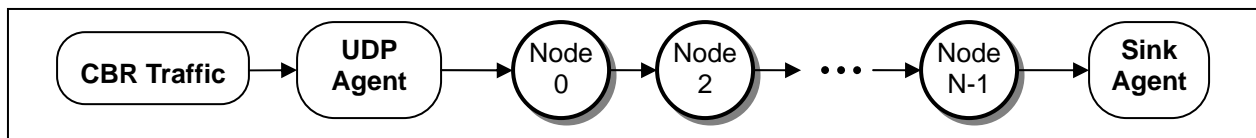


Figure 17: Block diagram of the ns-2 simulation scenario in chain topologies

The following list summarizes the various scenario parameters for our simulations:

- The length of the chain topologies is varied from **2** to **11 nodes**, denoted by  $N$ .
- The nodes are stationary within a confined area of **1000 m x 1000 m**.
- The number of orthogonal channels is varied from **2** to **4**, denoted by  $C$ .
- Each node is always equipped with **two interfaces (1 fixed, 1 switchable)**; thus,  $I = 2$ .
- The duration of each simulation is **60 s**; however, the actual simulation time increases as the length of the chain topology as well as the number of channels increase.
- The data rate or the channel capacity is set to **5.4 Mbps**.
- **Constant Bit-Rate (CBR) traffic** is passed from Node 0 to Node  $N-1$ , as denoted by the ns-2 CBR Traffic application and the Sink Agent in Figure 17.
  - The data packet size is set to **1000 bytes**.

- The CBR data (packetization) rate is set to **1.4 ms**, which is chosen to be large enough to saturate the network according to the following calculation:

$$\frac{(1000\text{bytes})(8\text{bits}/\text{byte})}{1.4\text{ms}} = 5.7\text{Mbps} > 5.4\text{Mbps},$$

where 5.4 Mbps is our specified channel capacity.

- Data packets are transported using the **User Datagram Protocol (UDP)**, as indicated by the UDP Agent in Figure 17. Therefore, no flow and congestion controls will take place that may impair the throughput performance. As well, the packets can be transmitted as fast as the packetization rate.

The performance of the above multi-channel multi-interface scenarios is compared with the performance of AODV when using a single channel single interface ad-hoc wireless network. The complete ns-2 simulation script, **wireless\_chain\_mcmi.tcl**, can be found in the Appendix section. The following highlights some key configurations for the case of 4 nodes in the chain topology, 3 channels and 2 interfaces:

|                                                                                                                                                       |                                                                                                                                                  |
|-------------------------------------------------------------------------------------------------------------------------------------------------------|--------------------------------------------------------------------------------------------------------------------------------------------------|
| <pre>Mac/802_11 set dataRate_ 5.4e6 set val(nn) 4 set val(nc) 3 set val(ni) 2 set pktsize 1000 set pktrate 0.0014</pre>                               | <p>} Set up the channel capacity, the numbers of nodes, channels, and interfaces, packet size and packetization rate.</p>                        |
| <pre>set udp0 [new Agent/UDP] \$ns_ attach-agent \$n(0) \$udp0</pre>                                                                                  | <p>} Define the UDP Agent and attach it to Node 0</p>                                                                                            |
| <pre>set sink0 [new Agent/Null] set last_node_id [expr \$val(nn)-1] \$ns_ attach-agent \$n(\$last_node_id) \$sink0 \$ns_ connect \$udp0 \$sink0</pre> | <p>} Define the Sink Agent and attach it to Node <i>N</i>-1. Finally connect the UDP Agent with the Sink Agent to specify the data flow path</p> |
| <pre>set cbr0 [new Application/Traffic/CBR] \$cbr0 attach-agent \$udp0 \$cbr0 set packetSize_ \$pktsize</pre>                                         | <p>} Define the CBR Traffic application, attach it to the UDP Agent, and specify the CBR packet size and data rate</p>                           |

**Figure 18: Simulation set-up for chain topology with 4 nodes, 3 channels and 2 interfaces**

## 6.2. Interface Switching Validation

In this section, we validate our implemented *interface switching* algorithm in ns-2's AODV routing agent, seeing if the simulated results comply with the proposed algorithm in [2]. A validation example involving 4 nodes in a chain topology, 3 channels (0, 1, 2), and 2 interfaces (0, 1) is presented as a table in Figure 19. This table is manually constructed by examining the debugging messages generated by a simulation run (with the MCMC\_DEBUG flag turned on when building the modified ns-2 simulator).

We will employ the following terminologies in the table:

- NT[N] is the *NeighbourTable* denoting the fixed channel used by Node N.
- CUL[C] is the *ChannelUsageList* denoting the number of nodes using channel C as their fixed channel.
- FC is the fixed channel currently used by the node for receiving data from its neighbour.
- SC is the switchable channel currently used by the node for sending data to its neighbour.

|                  | Node 0                                                                                                            | Node 1                                                                        | Node 2                                                                        | Node 3                                                           |
|------------------|-------------------------------------------------------------------------------------------------------------------|-------------------------------------------------------------------------------|-------------------------------------------------------------------------------|------------------------------------------------------------------|
| <b>Initially</b> | NT[0] = 2<br>CUL[2] = 1<br><br>FC = 2<br>SC = 0                                                                   | NT[1] = 0<br>CUL[0] = 1<br><br>FC = 0<br>SC = 2                               | NT[2] = 2<br>CUL[2] = 1<br><br>FC = 2<br>SC = 0                               | NT[3] = 1<br>CUL[1] = 1<br><br>FC = 1<br>SC = 2                  |
| <b>Step 1</b>    |                                                                                                                   | NT[0] = 2<br>NT[1] = 0<br>CUL[0] = 1<br>CUL[2] = 1                            | NT[1] = 0<br>NT[2] = 2<br>CUL[0] = 1<br>CUL[2] = 2                            | NT[2] = 2<br>NT[3] = 1<br>CUL[0] = 1<br>CUL[1] = 1<br>CUL[2] = 1 |
| <b>Step 2</b>    | NT[0] = 2<br>NT[1] = 0<br>CUL[0] = 1<br>CUL[2] = 2                                                                | NT[0] = 2<br>NT[1] = 0<br>NT[2] = 2<br>CUL[0] = 1<br>CUL[1] = 1<br>CUL[2] = 2 | NT[1] = 0<br>NT[2] = 2<br>NT[3] = 1<br>CUL[0] = 1<br>CUL[1] = 1<br>CUL[2] = 2 |                                                                  |
| <b>Step 3</b>    | FC = 2 but<br>CUL[2] is large<br>→ FC = 1<br><br>NT[0] = 1<br>NT[1] = 0<br>CUL[0] = 1<br>CUL[1] = 1<br>CUL[2] = 1 | NT[0] = 1<br>NT[1] = 0<br>NT[2] = 2<br>CUL[0] = 1<br>CUL[1] = 2<br>CUL[2] = 1 | NT[1] = 0<br>NT[2] = 2<br>NT[3] = 1<br>CUL[0] = 1<br>CUL[1] = 1<br>CUL[2] = 1 | NT[2] = 2<br>NT[3] = 1<br>CUL[0] = 1<br>CUL[1] = 1<br>CUL[2] = 1 |
| <b>Finally</b>   | FC = 1<br>SC = 0                                                                                                  | FC = 0<br>SC = 2                                                              | FC = 2<br>SC = 1                                                              | FC = 1<br>SC = 2                                                 |

Figure 19: Interface switching in chain topology with 4 nodes, 3 channels and 2 interfaces

Changes in the *NeighbourTable* and *ChannelUsageList* are highlighted in red; final results are highlighted in green. As can be seen in Figure 19, all the nodes choose a random channel for its fixed and switchable interfaces initially; their *NeighbourTable* and *ChannelUsageList* are populated accordingly. These tasks are performed in `command()`.

In Step 1, Node 0 calls `sendRequest()` to transmit a RREQ packet to Node 1 which carries the fixed channel used by Node 0 and its *NeighbourTable*. Upon `recvRequest()` receiving the RREQ, Node 1 updates its *NeighbourTable* with the fixed channel of Node 0 as well as its *ChannelUsageList* using Node 0's *NeighbourTable*. Next, Node 1 calls `forward()` to forward the RREQ packet to Node 2, which contains the fixed channel used by Node 1 and its updated *NeighbourTable*. Node 2 then updates its *NeighbourTable* with the fixed channel of Node 1 as well as its *ChannelUsageList* using Node 1's *NeighbourTable*. The similar RREQ packet exchange and *NeighbourTable* update are performed between Node 2 and Node 3 later as well.

In Step 2, Node 3 calls `sendReply()` to send a RREP to Node 2, which includes Node 3's fixed channel and *NeighbourTable*. In `recvReply()`, Node 2 updates its *NeighbourTable* with the fixed channel of Node 3 and its *ChannelUsageList* using Node 3's *NeighbourTable*. Node 2 then calls `forward()` to forward the RREP packet to Node 1, which contains the updated information. The similar actions are also performed between Node 1 and Node 0.

In Step 3, since the destination (Node 3) and the intermediate hops (Nodes 1, 2) have replied to the source (Node 0), a route has been established. CBR data traffic is now ready to be transported along the chain. However, while executing `forward()`, Node 0 realizes that its fixed interface is on Channel 2, which has a usage of 2 among all the available channels. With 2 being the largest usage currently, Node 0 reverses its *ChannelUsageList* about the fixed channel previously used and changes its fixed channel to Channel 1 since  $CUL[1] = 0$ . The randomly determined switching tendency at the time of simulation is 0.14, which is below the switching probability of 0.4; thus, the channel switching occurs. Node 0 then calls `sendHello()` to broadcast a Hello packet to inform its neighbours of its new fixed channel. When Node 0 and Node 1 receives the broadcasted Hello packet via `recvHello()`, they further update their *NeighbourTable* with the fixed channel of Node 0 and their *ChannelUsageList* using the *NeighbourTable* of Node 0.

Finally, we have verified that *interface switching* effectively assigns distinctive fixed channels to each successive node. None of the adjacent channels are interfering each other when data packets are transmitted along the chain. As a result, intermediate nodes can send data to the next node using its switchable interface, while receiving data on its fixed interface.

### 6.3. Simulation Output

Figure 20 presents a sample ns-2 simulation output for chain topology with 4 nodes, 3 Channels, and 2 interfaces. As can be observed, CBR traffic is flowed from Node 0 to Node 3. The simulation duration is 60 s and the average end-to-end throughput is 2586.36 kbps.

```

/usr/ns-allinone-2.32/mcmi_test
merry@BENQ-JOYBOOK /usr/ns-allinone-2.32/mcmi_test
$ ns wireless_chain_mcmi.tcl
Ad-Hoc Wireless Network in Chain Topologies - 4 Nodes, 3 Channels, 2 Interfaces
merry@BENQ-JOYBOOK /usr/ns-allinone-2.32/mcmi_test
$
 flowID: 0
 trafficType: cbr
 srcNode: 0
 destNode: 3
 startTime[s]: 1
 stopTime[s]: 61
 duration[s]: 60
 receivedPkts: 19468
 avgTput[kbps]: 2586.36

```

Figure 20: ns-2 output for chain topology with 4 nodes, 3 channels and 2 interfaces

The corresponding network animator (nam) view of the aforementioned simulation scenario is shown in the following figure. Please note that only wireless nodes (MobileNode) can be displayed in nam in the current ns-2. Dumping of traffic data and thus visualization of data packet movements for wireless scenarios is still not supported.

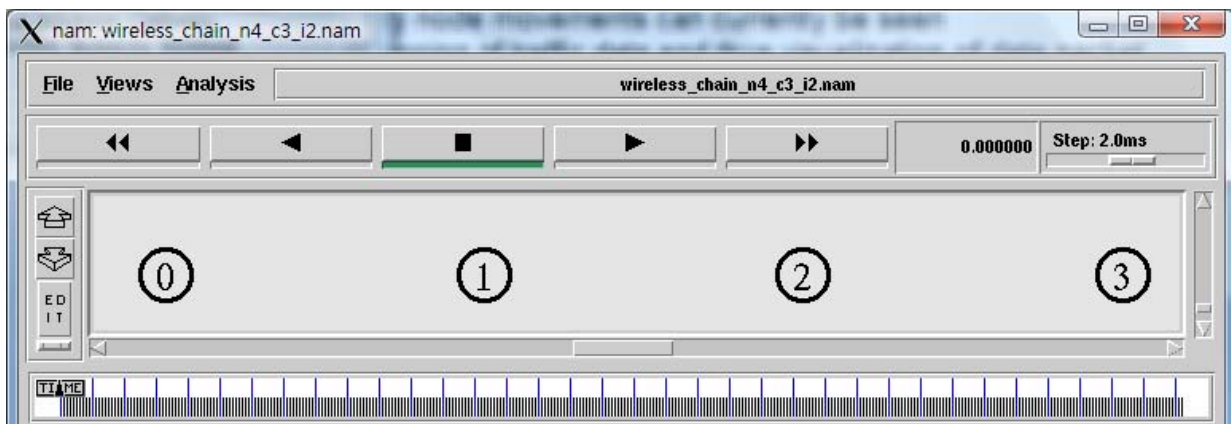


Figure 21: nam output for chain topology with 4 nodes, 3 channels and 2 interfaces

## 6.4. Throughput Performance

We evaluate the end-to-end throughput of the ad-hoc wireless network in simple chain topologies. The throughput is obtained using an AWK script written by Marco Fiore in `avgStats.awk`, which can be found in the Appendix section. When an ns-2 wireless simulation completes, we let the simulation script to call this AWK script automatically to analyze the wireless traffic trace generated by the simulator. While varying the number of channels from 2 to 4 and equipping each node with two IEEE 802.11 interfaces, we designate each scenario setting as  $x\text{C}2\text{I}$ , where  $x$  is the number of channels.

In each scenario, the length of the chain is varied from 2 to 11 nodes. We compare the results of  $x\text{C}2\text{I}$  with chain topologies in single channel and single interface; that is,  $1\text{C}1\text{I}$ . Since channel assignments for each node are carried out on a random basis, for the throughput of each scenario, we perform the simulation four times and use the average value of the four obtained throughputs to plot the comparison graph, as illustrated in the following figure.

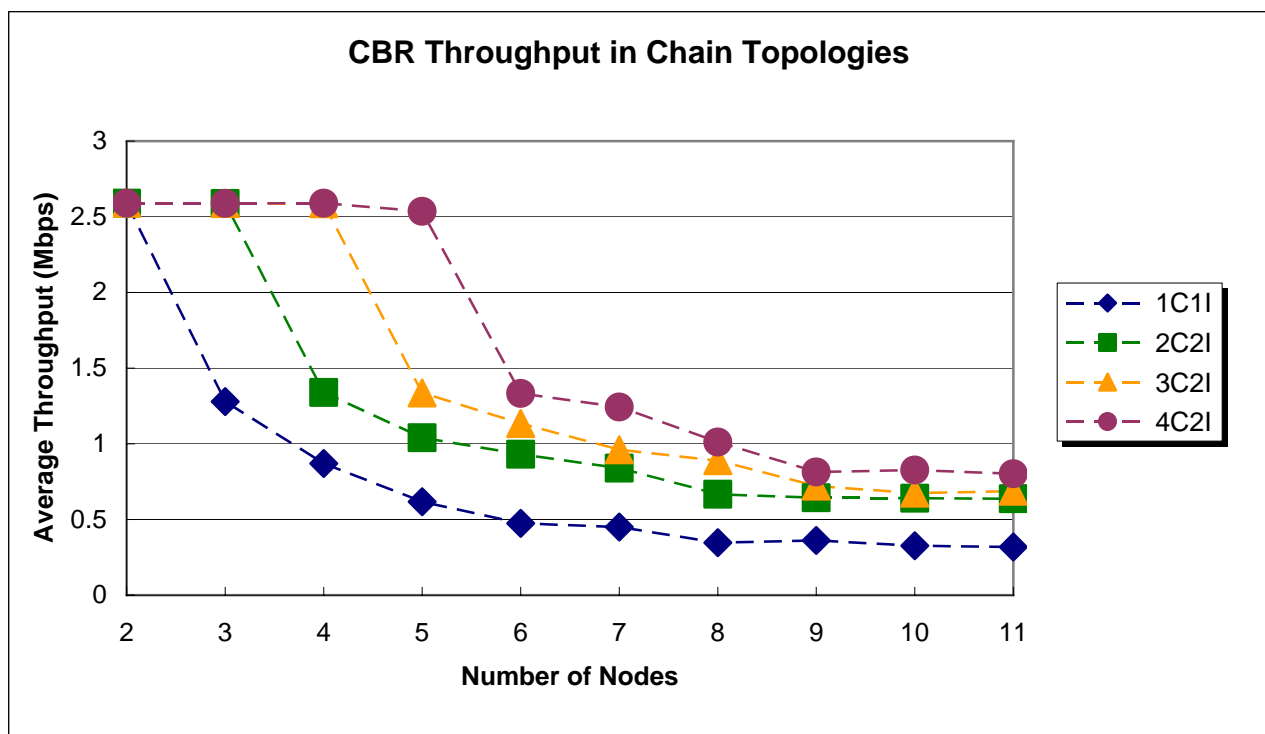


Figure 22: Throughputs of single-channel and multi-channel networks in chain topologies

As can be seen from the graph above, the throughput of the  $1\text{C}1\text{I}$  scenario degrades as the number of nodes along the chain increases by 1 each time. One reason is that intermediate nodes



cannot send and receive data at the same time, reducing the achievable throughput by about half. Another factor is that the interference within the carrier sense range inhibits transmissions on adjacent nodes; thus, the resultant throughput is also impaired.

By contrast, we observe higher throughputs with multiple channels and multiple interfaces on each node, such as 2C2I, 3C2I and 4C2I. Due to the fact that *interface switching* assigns the fixed channel of successive nodes to different channels, intermediate nodes can send data to the next node using its switchable interface, while receiving data on its fixed interface.

However, the throughput improvements are smaller when the number of nodes is greater than the number of channels + 1. When the chain length goes beyond this threshold, some adjacent nodes will be likely on some common channels; thus, interference will degrade the throughput. Nevertheless, the overall throughputs of these multi-channel multi-interface scenarios are generally still higher than the case of 1C1I.

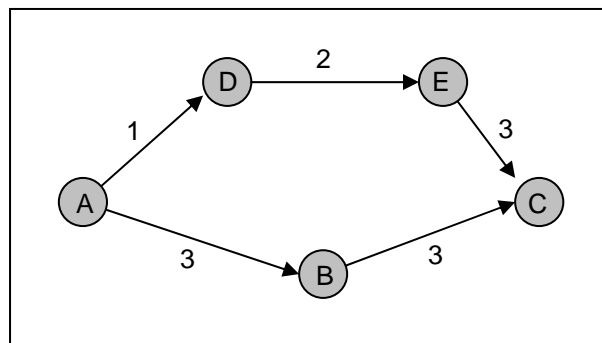
In addition, more channels are useful with longer chains. Over a chain of two nodes, at most two channels can be utilized (one channel for each node) despite more available channels may be available. Therefore, the performance with three or more channels is similar as the performance of two channels over a two-node chain (though higher than that of 1C1I). By the same idea, at most three channels can be used over a chain of three nodes; the performance with more channels over a three-node chain is also similar. On the other hand, over a chain of 5 nodes, more channels can be utilized over different nodes; therefore, having 4 channels is better than having 3 or less number of channels.

In words, the simulation results have demonstrated that multiple channels and multiple interfaces can noticeably improve the throughput in ad-hoc wireless networks. Furthermore, the number of wireless interfaces can be less than the number of available channels. In our simulation, even two interfaces can utilize multiple channels.

## 7. Future Work and Improvement

Due to the time constraint of the project, we do not implement the Multi-Channel Routing (MCR) metric proposed by [2] in the ns-2 simulator. Therefore, routing is not tested during our simulation. As pointed out in the paper, merely the *interface switching* solution implemented at the data-link layer may not be sufficient for effectively utilizing multiple channels, as the routing protocol may select routes wherein successive hops interfere with each other.

In this project, we have enhanced the existing implementation of the AODV routing agent in ns-2 to support multiple interfaces. AODV selects the route that has the least number of hops through other nodes [2]. The shortest path does not distinguish between a route that uses many channels and a route that uses fewer channels. The following diagram illustrates the need for a specialized routing protocol. Suppose all nodes in the ad-hoc wireless network in the figure have already chosen their fixed interfaces. AODV would favour routes A-B-C as it is the shortest. However, routes A-D-E-C use different channels between two adjacent nodes. This path allows all links to be active simultaneously; thus, the end-to-end throughput can be potentially higher. On the other hand, both links A-B and B-C use channel 3; only one of the links can be active at a time, implying the resulting throughput is halved.



**Figure 23: The scenario that illustrates the need for selecting channel deviated routes [2]**

As a result, the more ideal implementation of a multi-channel multi-interface ad-hoc wireless network is to incorporate the proposed MCR as the routing agent in ns-2. With this capability, we are able to effectively evaluate the throughput performance in random topologies with multiple traffic flows. However, the underlying work required for implementing the MCR may be beyond the scope of a regular course project.

## 8. Conclusion

In this project, we have extended the Network Simulator (ns-2.32) to support multiple channels and multiple interfaces by modifying ns-2's *MobileNode* library to support multiple interfaces. By referring to the guideline written by R. A. Calvo and J. P. Campo in [1], the legacy operations of IEEE 802.11 wireless interfaces in ns-2 can be preserved. Not only have we gained valuable exposure to ns-2's wireless node model, we have an opportunity to acquire a good understanding of ns-2's architecture.

We have also studied and explored the use of multiple channels and multiple interfaces in wireless ad-hoc networks. We have learned about the interface assignment strategy proposed by P. Kyasanur and N. H. Vaidya in [2], which uses the notion of fixed and switchable interfaces. This assignment approach allows effective utilization of all the available channels even when the number of interfaces is smaller than the number of channels. Although quite a few amount of times have been spent to study the Ad-hoc On-demand Distance Vector (AODV) routing in ns-2, we have successfully integrated this algorithm in the AODV routing agent in ns-2. With the implementation of the multi-channel multi-interface support in ns-2 as well as the development of the simulation script, we have expanded our horizons in C++ and TCL programming.

Next, using the modified ns-2 with the multi-interface and *interface switching* capabilities, we have simulated a multi-channel multi-interface ad-hoc wireless network in simple chain topologies. By varying the number of available channels and the number of nodes in the chain while keeping each node to have two interfaces (one fixed and one switchable), we have generated various *interface switching* interactions and end-to-end throughputs using the modified ns-2. Simulation results validate the effectiveness of our implemented *interface switching* algorithm in ns-2's AODV routing protocol. More importantly, the results demonstrate that network throughput can significantly improve in ad-hoc wireless networks with multiple channels and multiple interfaces in each wireless node.

## Glossary

|       |                                                   |
|-------|---------------------------------------------------|
| AODV  | Ad-hoc On-demand Distance Vector                  |
| ARP   | Address Resolution Protocol                       |
| CBR   | Constant Bit Rate                                 |
| CUL   | ChannelUsageList                                  |
| DSDV  | Destination-Sequenced Distance-Vector Routing     |
| DSR   | Dynamic Source Routing                            |
| IEEE  | Institute of Electrical and Electronics Engineers |
| IP    | Internet Protocol                                 |
| MAC   | Media Access Control                              |
| MCMCI | Multi-Channel Multi-Interface                     |
| nam   | Network Animator                                  |
| ns-2  | Network Simulator 2                               |
| NT    | NeighbourTable                                    |
| OTCL  | Object TCL                                        |
| RFC   | Request for Comments                              |
| RERR  | Route Error                                       |
| RREP  | Route Reply                                       |
| RREQ  | Route Request                                     |
| SSCH  | Slotted Seeded Channel Hopping                    |
| TCL   | Tool Command Language                             |
| TORA  | Temporally Ordered Routing Algorithm              |
| UDP   | User Datagram Protocol                            |
| WCETT | Weighted Cumulative Expected Transmission Time    |

## References

- [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2," [User Guide], University of Cantabria, Jan. 2007.
- [2] P. Kyasanur and N. H. Vaidya, "Routing and Link-layer Protocols for Multi-Channel Multi-Interface Ad Hoc Wireless Networks," *SIGMOBILE Mobile Computing and Communications Review*, vol. 10, no. 1, pp. 31-43, Jan. 2006.
- [3] A. Raniwala, K. Gopalan, and T. Chiueh, "Centralized Channel Assignment and Routing Algorithms for Multi-Channel Wireless Mesh Networks," *Mobile Computing and Communications Review (MC2R) 2004*, vol. 8, no. 2, pp. 50-65, Apr. 2004.
- [4] S. Wu, C. Lin, Y. Tseng, and J. Sheu, "A New Multi-Channel MAC Protocol with On-Demand Channel Assignment for Multi-Hop Mobile Ad Hoc Networks," *International Symposium on Parallel Architectures, Algorithms and Networks (ISPAN), 2000*, Dallas, TX, pp. 232-237, Dec. 2000.
- [5] R. Maheshwari, H. Gupta, and S. R. Das, "Multichannel MAC Protocols for Wireless Networks," *Sensor and Ad Hoc Communications and Networks (SECON) 2006*, Reston, VA, vol. 2, pp. 393-401, Sept. 2006.
- [6] P. Bahl, R. Chandra, and J. Dunagan, "SSCH: Slotted Seeded Channel Hopping for Capacity Improvement in IEEE 802.11 Ad-Hoc Wireless Networks," *ACM Annual International Conference on Mobile Computing and Networking (MobiCom) 2004*, Philadelphia, PA, pp. 216-230, Oct. 2004.
- [7] R. Draves, J. Padhye, and B. Zill, "Routing in Multi- Radio, Multi-Hop Wireless Mesh Networks," *ACM Annual International Conference on Mobile Computing and Networking (MobiCom) 2004*, Philadelphia, PA, pp. 114-128, Oct. 2004.
- [8] C. Perkins, et. al., "Ad hoc On-Demand Distance Vector (AODV) Routing," RFC 3561, *Network Working Group*, The Internet Society, Jul. 2003.
- [9] K. Fall and K. Varadhan, The ns Manual [Online]. Available: <http://www.isi.edu/nsnam/ns>.

## Appendix

### A.1. Code Listing

This section provides the ns-2 source code that we have modified in implementing the multi-channel multi-interface ad-hoc wireless network, as well as the scenario (test) scripts we use to demonstrate the effectiveness of exploiting multiple channels and multiple interfaces. Please note that due to lack of space, only the modified lines of code (marked with a “MCMI” tag) along with the entire function body that they belong to are presented in this section.

#### Test Script

- wireless\_chain\_mcmi.tcl
- avgStats.awk

#### Source Code

##### /ns-2.32

- **/aodv**
  - /aodv.cc
  - /aodv
  - /aodv\_packet.h
  - /aodv\_rtable.h
- **/common**
  - /mobilenode.cc
  - /mobilenode.h
- **/mac**
  - /channel.cc
  - /mac.h
  - /mac-802\_11.cc
- **/tcl**
  - **/lib**
    - /ns-lib.tcl
    - /ns-mobilenode.tcl

```

1 #=====
2 # ENSC 835: High-Performance Networks
3 # Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Networks
4 #
5 # Student: Chih-Hao Howard Chang
6 # 20007-2192
7 # howardc@sfu.ca
8 #
9 # Description: Simulation script for multi-Channel multi-Interface Ad-Hoc
10 # wireless network in chain topologies
11 #
12 # File: wireless_chain_mcmi.tcl
13 #=====
14
15
16 #=====
17 # MAC Layer Setup
18 #=====
19 Mac/802_11 set dataRate_ 5.4e6 ;# rate for data frames in Mbps
20
21
22 #=====
23 # Simulation Parameters Setup
24 #=====
25 set val(chan) Channel/WirelessChannel ;# channel type
26 set val(prop) Propagation/TwoRayGround ;# radio-propagation model
27 set val(netif) Phy/WirelessPhy ;# network interface type
28 set val(mac) Mac/802_11 ;# MAC type
29 set val(ifq) Queue/DropTail/PriQueue ;# interface queue type
30 set val(ll) LL ;# link layer type
31 set val(ant) Antenna/OmniAntenna ;# antenna model
32 set val(ifqlen) 50 ;# max packet in ifq
33 set val(rp) AODV ;# routing protocol
34 set val(x) 1000 ;# X dimension of topography
35 set val(y) 1000 ;# Y dimension of topography
36 set val(stop) 70 ;# nam stop time
37
38 set val(nn) 4 ;# number of mobilenodes
39 set val(nc) 3 ;# number of channels
40 set val(ni) 2 ;# number of interfaces, <= number of channels
41 set pktsize 1000 ;# packet size in bytes
42 set pktrate 0.0014 ;# packet rate in seconds
43
44 set filename wireless_chain_n$val(nn)_c$val(nc)_i$val(ni)
45 puts "Ad-Hoc Wireless Network in Chain Topologies - $val(nn) Nodes, $val(nc) Channels,
46 $val(ni) Interfaces"
47
48 #=====
49 # Initialization
50 #=====
51 # Create a ns simulator
52 set ns_ [new Simulator]
53
54 # Setup topography object
55 set topo [new Topography]
56 $topo load_flatgrid $val(x) $val(y)
57 set god_ [create-god [expr $val(nn)*$val(nc)]]
58
59 # Open the NS trace file
60 set tracefd [open $filename.tr w]
61 $ns_ trace-all $tracefd
62 $ns_ use-newtrace
63
64 # Open the NAM trace file
65 set namfile [open $filename.nam w]
66 $ns_ namtrace-all $namfile
67 $ns_ namtrace-all-wireless $namfile $val(x) $val(y)
68
69 # Create wireless channels
70 for {set i 0} {$i < $val(nc)} {incr i} {

```

```
71 set chan($i) [new $val(chan)]
72 }
73
74
75 #=====
76 # Mobile Node Parameter Setup
77 #=====
78 $ns_ node-config -adhocRouting $val(rp) \
79 -llType $val(ll) \
80 -macType $val(mac) \
81 -ifqType $val(ifq) \
82 -ifqLen $val(ifqlen) \
83 -antType $val(ant) \
84 -propType $val(prop) \
85 -phyType $val(netif) \
86 -channel $chan(0) \
87 -topoInstance $topo \
88 -agentTrace ON \
89 -routerTrace ON \
90 -macTrace OFF \
91 -movementTrace OFF \
92 -ifNum $val(ni)
93
94
95 #=====
96 # Nodes Definition
97 #=====
98 $ns_ change-numifs $val(nc)
99 for {set i 0} {$i < $val(nc)} {incr i} {
100 $ns_ add-channel $i $chan($i)
101 }
102
103 # Create nodes
104 for {set i 0} {$i < $val(nn)} {incr i} {
105 set n($i) [$ns_ node]
106 $god_ new_node $n($i)
107 }
108
109 # Set node positions in horizontal chain topology
110 set nodedist 250
111 for {set i 0} {$i < $val(nn)} {incr i} {
112 $n($i) set X_ [expr $i * $nodedist + 20]
113 $n($i) set Y_ 50
114 $n($i) set Z_ 0.0
115 $ns_ initial_node_pos $n($i) 40
116 $n($i) random-motion 0
117 }
118
119
120 #=====
121 # Agents Definition
122 #=====
123 set udp0 [new Agent/UDP]
124 $ns_ attach-agent $n(0) $udp0
125
126 set sink0 [new Agent/Null]
127 set last_node_id [expr $val(nn)-1]
128 $ns_ attach-agent $n($last_node_id) $sink0
129 $ns_ connect $udp0 $sink0
130
131
132 #=====
133 # Applications Definition
134 #=====
135 # Setup a CBR Application over UDP connection
136 set cbr0 [new Application/Traffic/CBR]
137 $cbr0 attach-agent $udp0
138 $cbr0 set packetSize_ $pktsize
139 $cbr0 set interval_ $pktrate
140 $ns_ at 1.0 "$cbr0 start"
141 $ns_ at 61.0 "$cbr0 stop"
```



```
142
143
144 #=====
145 # Simulation Termination
146 #=====
147 # Define a finish procedure
148 proc finish {} {
149 global ns_ tracefd filename pktsize last_node_id
150 global namfile
151 $ns_ flush-trace
152 close $tracefd
153 close $namfile
154 exec nam $filename.nam &
155
156 # Call throughput analyzer (AWK scripts written by Marco Fiore, marco.fiore@polito.it)
157 exec awk -f avgStats.awk src=0 dst=$last_node_id flow=0 pkt=$pktsize $filename.tr &
158 exit 0
159 }
160 for {set i 0} {$i < $val(nn)} {incr i} {
161 $ns_ at $val(stop) "\$n($i) reset"
162 }
163 $ns_ at $val(stop) "$ns_ nam-end-wireless $val(stop)"
164 $ns_ at $val(stop) "finish"
165 $ns_ at $val(stop) "puts \"done\" ; $ns_ halt"
166 $ns_ run
167
```

```
1 #=====
2 # ENSC 835: High-Performance Networks
3 # Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Networks
4 #
5 # Student: Chih-Hao Howard Chang
6 # 20007-2192
7 # howardc@sfu.ca
8 #
9 # Description: Post-processing script for analyzing the average throughput based
10 # on the wireless traffic trace produced by a ns-2 wireless
11 # simulation, written by Marco Fiore, marco.fiore@polito.it
12 #
13 # File: avgStats.awk
14 #=====
15
16
17 BEGIN {
18 recvdSize = 0
19 startTime = 1e6
20 stopTime = 0
21 }
22
23 {
24 # Trace line format: normal
25 if ($2 != "-t") {
26 event = $1
27 time = $2
28 if (event == "+" || event == "-") node_id = $3
29 if (event == "r" || event == "d") node_id = $4
30 flow_id = $8
31 pkt_id = $12
32 pkt_size = $6
33 flow_t = $5
34 level = "AGT"
35 }
36 # Trace line format: new
37 if ($2 == "-t") {
38 event = $1
39 time = $3
40 node_id = $5
41 flow_id = $39
42 pkt_id = $41
43 pkt_size = $37
44 flow_t = $45
45 level = $19
46 }
47
48 # Store packets send time
49 if (level == "AGT" && flow_id == flow && node_id == src &&
50 sendTime[pkt_id] == 0 && (event == "+" || event == "s") && pkt_size >= pkt) {
51 if (time < startTime) {
52 startTime = time
53 }
54 sendTime[pkt_id] = time
55 this_flow = flow_t
56 }
57
58 # Update total received packets' size and store packets arrival time
59 if (level == "AGT" && flow_id == flow && node_id == dst &&
60 event == "r" && pkt_size >= pkt) {
61 if (time > stopTime) {
62 stopTime = time
63 }
64 # Rip off the header
65 hdr_size = pkt_size % pkt
66 pkt_size -= hdr_size
67 # Store received packet's size
68 recvdSize += pkt_size
69 # Store packet's reception time
70 rcvTime[pkt_id] = time
71 }
72 }
```

```

72
73 }
74
75 END {
76 # Compute average delay
77 delay = avg_delay = recvdNum = 0
78 for (i in recvTime) {
79 if (sendTime[i] == 0) {
80 printf("\nError in delay.awk: receiving a packet that wasn't sent
%g\n",i)
81 }
82 delay += recvTime[i] - sendTime[i]
83 recvdNum ++
84 }
85 if (recvdNum != 0) {
86 avg_delay = delay / recvdNum
87 } else {
88 avg_delay = 0
89 }
90
91 # Compute average jitters
92 jitter1 = jitter2 = jitter3 = jitter4 = jitter5 = 0
93 prev_time = delay = prev_delay = processed = deviation = 0
94 prev_delay = -1
95 for (i=0; processed<recvdNum; i++) {
96 if(recvTime[i] != 0) {
97 if(prev_time != 0) {
98 delay = recvTime[i] - prev_time
99 e2eDelay = recvTime[i] - sendTime[i]
100 if(delay < 0) delay = 0
101 if(prev_delay != -1) {
102 jitter1 += abs(e2eDelay - prev_e2eDelay)
103 jitter2 += abs(delay-prev_delay)
104 jitter3 += (abs(e2eDelay-prev_e2eDelay) - jitter3) /
16
105 jitter4 += (abs(delay-prev_delay) - jitter4) / 16
106 }
107 # deviation += (e2eDelay-avg_delay)*(e2eDelay-avg_delay)
108 prev_delay = delay
109 prev_e2eDelay = e2eDelay
110 }
111 prev_time = recvTime[i]
112 processed++
113 }
114 }
115 if (recvdNum != 0) {
116 jitter1 = jitter1*1000/recvdNum
117 jitter2 = jitter2*1000/recvdNum
118 }
119 # if (recvdNum > 1) {
120 # jitter5 = sqrt(deviation/(recvdNum-1))
121 # }
122
123 # Output
124 if (recvdNum == 0) {
125 printf("#####\n
n" \
126 "# Warning: no packets were received, simulation may be too short
#\n" \
127 "#####\n
n\n")
128 }
129 printf("\n")
130 printf(" %15s: %g\n", "flowID", flow)
131 printf(" %15s: %s\n", "trafficType", this_flow)
132 printf(" %15s: %d\n", "srcNode", src)
133 printf(" %15s: %d\n", "destNode", dst)
134 printf(" %15s: %d\n", "startTime[s]", startTime)
135 printf(" %15s: %d\n", "stopTime[s]", stopTime)
136 printf(" %15s: %d\n", "duration[s]", stopTime-startTime)
137 printf(" %15s: %g\n", "receivedPkts", recvdNum)

```

```
138 printf(" %15s: %g\n", "avgTput[kbps]", (recvdSize/(stopTime-startTime))*(8/1000))
139 }
140
141 function abs(value) {
142 if (value < 0) value = 0-value
143 return value
144 }
145
```

```
1 /*
2 #=====
3 # ENSC 835: High-Performance Networks
4 # Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Networks
5 #
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9 #
10 # Description: Modified AODV routing agent with the implementation of the
11 # multi-interface approach developed by [1] and the interface
12 # switching protocol proposed by [2], which support multi-channel
13 # multi-interface ad-hoc wireless network simulation
14 #
15 # File: aadv.cc
16 #
17 # [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2,"
18 # University of Cantabria, Jan. 2007 (User Guide).
19 #
20 # [2] P. Kyasanur and N. H. Vaidya, "Routing and Link-layer Protocols for
21 # Multi-Channel Multi-Interface Ad Hoc Wireless Networks," SIGMOBILE Mobile
22 # Computing and Communications Review, vol. 10, no. 1, pp. 31-43, Jan. 2006.
23 #=====
24 */
25
26 ...
27
28 // MCFI: Toggle MCFI debugging (used only when required)
29 //#define MCFI_DEBUG
30
31 ...
32
33 int
34 AODV::command(int argc, const char*const* argv) {
35 if(argc == 2) {
36 Tcl& tcl = Tcl::instance();
37
38 if(strncasecmp(argv[1], "id", 2) == 0) {
39 tcl.resultf("%d", index);
40 return TCL_OK;
41 }
42
43 if(strncasecmp(argv[1], "start", 2) == 0) {
44 btimer.handle((Event*) 0);
45
46 #ifndef AODV_LINK_LAYER_DETECTION
47 htimer.handle((Event*) 0);
48 ntimer.handle((Event*) 0);
49 #endif // LINK LAYER DETECTION
50
51 rtimer.handle((Event*) 0);
52 return TCL_OK;
53 }
54 }
55
56 else if(argc == 3) {
57 if(strcmp(argv[1], "index") == 0) {
58 index = atoi(argv[2]);
59 return TCL_OK;
60 }
61
62 else if(strcmp(argv[1], "log-target") == 0 || strcmp(argv[1], "tracetarget") == 0) {
63 logtarget = (Trace*) TclObject::lookup(argv[2]);
64 if(logtarget == 0)
65 return TCL_ERROR;
66 return TCL_OK;
67 }
68 else if(strcmp(argv[1], "drop-target") == 0) {
69 int stat = rqueue.command(argc, argv);
70 if (stat != TCL_OK) return stat;
71 return Agent::command(argc, argv);

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72 }
73 else if(strcmp(argv[1], "if-queue") == 0) {
74 ifqueue = (PriQueue*) TclObject::lookup(argv[2]);
75
76 if(ifqueue == 0)
77 return TCL_ERROR;
78 return TCL_OK;
79 }
80 else if (strcmp(argv[1], "port-dmux") == 0) {
81 dmux_ = (PortClassifier *)TclObject::lookup(argv[2]);
82 if (dmux_ == 0) {
83 fprintf (stderr, "%s: %s lookup of %s failed\n", __FILE__,
84 argv[1], argv[2]);
85 return TCL_ERROR;
86 }
87 return TCL_OK;
88 }
89 }
90
91 // MCMI: Populate ifqueuelist and targetlist, take the corresponding values
92 // while the interfaces are being created in the nodes.
93 else if (argc == 4) {
94 if (strcmp(argv[1], "if-queue") == 0) {
95 int temp_ = atoi(argv[2]);
96 if (temp_ == nIfaces) {
97 nIfaces++;
98 }
99 ifqueuelist[temp_] = (PriQueue *)TclObject::lookup(argv[3]);
100 if (ifqueuelist[temp_]) {
101 // MCMI: Initially, choose a random interface index in the range
102 // 0 to temp_. Have to do it everytime since the node does
103 // not know how many channels it is going to have.
104 fixed_interface = rand() % (temp_+1);
105 #ifdef MCMI_DEBUG
106 printf("n%d\tcommand - fixed_interface (initial): %d\n", (int)index,
107 fixed_interface);
108 #endif
109 if (nIfaces > 1) {
110 do {
111 switchable_interface = rand() % (temp_+1);
112 } while (switchable_interface == fixed_interface);
113 #ifdef MCMI_DEBUG
114 printf("n%d\tcommand - switchable_interface (initial): %d\n",
115 (int)index, switchable_interface);
116 #endif
117 }
118
119 // MCMI: Initialize NeighbourTable and ChannelUsageList
120 for (int i = 0; i < MAX_NT_CUL_ENTRIES; i++) {
121 neighbour_table[i] = -1;
122 channel_usage_list[i] = 0;
123 }
124
125 // MCMI: Add the fixed channel used by this node to its NeighbourTable
126 int fixed_channel = ((Mac
127 *)ifqueuelist[fixed_interface]->target()->netif()->channel()->index());
128 neighbour_table[(int)index] = fixed_channel; // fixed channel being used by
129 the node itself
130 channel_usage_list[fixed_channel]++;
131 #ifdef MCMI_DEBUG
132 printf("n%d\tcommand - fixed_channel: %d\n", (int)index, fixed_channel);
133 rt_print_NT_CUL();
134 #endif
135 return TCL_OK;
136 } else {
137 return TCL_ERROR;
138 }
139 } // if (strcmp(argv[1], "if-queue") == 0)
140 if (strcmp(argv[1], "target") == 0) {
141 int temp_ = atoi(argv[2]);
142 if (temp_ == nIfaces) {

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139 nIfaces++;
140 }
141 targetlist[temp_] = (NsObject *)TclObject::lookup(argv[3]);
142 if (targetlist[temp_] {
143 return TCL_OK;
144 } else {
145 return TCL_ERROR;
146 }
147 } // if (strcmp(argv[1], "target") == 0)
148 }
149
150 return Agent::command(argc, argv);
151 }
152
153
154 /*
155 * Constructor
156 */
157 AODV::AODV(nsaddr_t id) : Agent(PT_AODV),
158 btimer(this), htimer(this), ntimer(this),
159 rtimer(this), lrtimer(this), rqueue() {
160
161 index = id;
162 seqno = 2;
163 bid = 1;
164
165 LIST_INIT(&nbhead);
166 LIST_INIT(&bihead);
167
168 logtarget = 0;
169 ifqueue = 0;
170 nIfaces = 0; // MCMI: Initialize the number of interfaces used by this node
171 }
172
173 ...
174
175 // MCMI: Print the routing table of a node
176 void AODV::rt_print(nsaddr_t node_id) {
177 FILE *out_file;
178 char nowfile[50];
179 sprintf(nowfile, "rtable_node%d.txt", (int)node_id);
180 out_file = fopen(nowfile, "a");
181 aodv_rt_entry *rt;
182 fprintf(out_file, "-----\n");
183 for (rt = rtable.head(); rt; rt = rt->rt_link.le_next) {
184 fprintf(out_file, "NODE: %i\tTIME: %.4lf\tDST: %i\tNXTHOP: %i\tHOPS: %i\tSEQ:
185 %i\tEXP: %.4lf\tFLAG: n%d\tIF: %d\n",
186 (int)node_id, CURRENT_TIME, rt->rt_dst, rt->rt_nexthop,
187 rt->rt_hops, rt->rt_seqno, rt->rt_expire, rt->rt_flags,
188 (int)rt->rt_interface);
189 }
190 fclose(out_file);
191 }
192
193 // MCMI: Print NeighbourTable and ChannelUsageList of this node
194 void AODV::rt_print_NT_CUL() {
195 printf("n%d\t\tNT\t\tCUL\n", (int)index);
196 for (int i = 0; i < MAX_NT_CUL_ENTRIES; i++) {
197 printf("\tn/c%d\t%d\t%d\n", i, neighbour_table[i], channel_usage_list[i]);
198 }
199 printf("\n");
200 }
201
202 ...
203
204 void
205 AODV::recvRequest(Packet *p) {
206 struct hdr_cmh *ch = HDR_CMH(p); // MCMI: Get the common header of this packet
207 struct hdr_ip *ih = HDR_IP(p);
208 struct hdr_aodv_request *rq = HDR_AODV_REQUEST(p);

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```
209 aadv_rt_entry *rt;
210
211 /*
212 * Drop if:
213 * - I'm the source
214 * - I recently heard this request.
215 */
216 if(rq->rq_src == index) {
217 #ifdef DEBUG
218 fprintf(stderr, "%s: got my own REQUEST\n", __FUNCTION__);
219 #endif // DEBUG
220 Packet::free(p);
221 return;
222 }
223 if (id_lookup(rq->rq_src, rq->rq_bcast_id)) {
224 #ifdef DEBUG
225 fprintf(stderr, "%s: discarding request\n", __FUNCTION__);
226 #endif // DEBUG
227 Packet::free(p);
228 return;
229 }
230
231 /*
232 * Cache the broadcast ID
233 */
234 id_insert(rq->rq_src, rq->rq_bcast_id);
235
236
237 if (nIfaces) {
238 // MCMI: When the node receives a "Hello" (RREQ) packet from a neighbour,
239 // update its NeighbourTable with the fixed channel of that neighbour.
240 neighbour_table[(int)(rq->rq_sender_node_ip)] = rq->rq_fixed_channel_used;
241
242 // MCMI: Update the node's ChannelUsageList using the NeighbourTable of its
243 // neighbour. Doing so ensures ChannelUsageList will contain two-hop
244 // channel usage information.
245 for (int i = 0; i < MAX_NT_CUL_ENTRIES; i++) {
246 if (i == (int)index) {
247 continue; // MCMI: No need to deal with the fixed channel used
248 // by itself again; avoid repetition
249 }
250 if (rq->rq_neighbour_table[i] != -1) {
251 channel_usage_list[rq->rq_neighbour_table[i]]++;
252 }
253 }
254 #ifdef MCMI_DEBUG
255 printf("n%d\trecvRequest - rq->rq_sender_node_ip: %d\n", (int)index,
256 rq->rq_sender_node_ip);
257 printf("n%d\trecvRequest - rq->rq_fixed_channel_used: %d\n", (int)index,
258 rq->rq_fixed_channel_used);
259 rt_print_NT_CUL();
260 #endif
261 } // if (nIfaces)
262
263 /*
264 * We are either going to forward the REQUEST or generate a
265 * REPLY. Before we do anything, we make sure that the REVERSE
266 * route is in the route table.
267 */
268 aadv_rt_entry *rt0; // rt0 is the reverse route
269
270 rt0 = rtable.rt_lookup(rq->rq_src);
271 if(rt0 == 0) { /* if not in the route table */
272 // create an entry for the reverse route.
273 rt0 = rtable.rt_add(rq->rq_src);
274
275 // MCMI: Get the interface index
276 if (nIfaces) {
277 #ifdef MCMI_DEBUG
```



```

278 printf("n%d\trecvRequest - ch->iface() (rt0 = 0): %d\n", (int)index,
ch->iface());
279 printf("n%d\trecvRequest - ((Mac *)ifqueuelist[0]->target()->addr() (rt0 = 0):
%d\n", (int)index, ((Mac *)ifqueuelist[0]->target()->addr());
280 #endif
281 rt0->rt_interface = ch->iface()-((Mac *)ifqueuelist[0]->target()->addr();
282 } else {
283 rt0->rt_interface = -1;
284 }
285 #ifdef MCMCI_DEBUG
286 printf("n%d\trecvRequest - rt0->rt_interface (rt0 = 0): %d\n\n", (int)index,
rt0->rt_interface);
287 #endif
288 }
289
290 rt0->rt_expire = max(rt0->rt_expire, (CURRENT_TIME + REV_ROUTE_LIFE));
291
292 if ((rq->rq_src_seqno > rt0->rt_seqno) ||
293 ((rq->rq_src_seqno == rt0->rt_seqno) &&
294 (rq->rq_hop_count < rt0->rt_hops))) {
295 // If we have a fresher seq no. or lesser #hops for the
296 // same seq no., update the rt entry. Else don't bother.
297 rt_update(rt0, rq->rq_src_seqno, rq->rq_hop_count, ih->saddr(),
298 max(rt0->rt_expire, (CURRENT_TIME + REV_ROUTE_LIFE)));
299
300
301 // MCMCI: Get the interface index
302 if (nIfaces) {
303 #ifdef MCMCI_DEBUG
304 printf("n%d\trecvRequest - ch->iface(): %d\n", (int)index, ch->iface());
305 printf("n%d\trecvRequest - ((Mac *)ifqueuelist[0]->target()->addr(): %d\n",
(int)index, ((Mac *)ifqueuelist[0]->target()->addr());
306 #endif
307 rt0->rt_interface = ch->iface()-((Mac *)ifqueuelist[0]->target()->addr();
308 } else {
309 rt0->rt_interface = -1;
310 }
311 #ifdef MCMCI_DEBUG
312 printf("n%d\trecvRequest - rt0->rt_interface: %d\n\n", (int)index,
rt0->rt_interface);
313 #endif
314
315
316 if (rt0->rt_req_timeout > 0.0) {
317 // Reset the soft state and
318 // Set expiry time to CURRENT_TIME + ACTIVE_ROUTE_TIMEOUT
319 // This is because route is used in the forward direction,
320 // but only sources get benefited by this change
321 rt0->rt_req_cnt = 0;
322 rt0->rt_req_timeout = 0.0;
323 rt0->rt_req_last_ttl = rq->rq_hop_count;
324 rt0->rt_expire = CURRENT_TIME + ACTIVE_ROUTE_TIMEOUT;
325 }
326
327 /* Find out whether any buffered packet can benefit from the
328 * reverse route.
329 * May need some change in the following code - Mahesh 09/11/99
330 */
331 assert (rt0->rt_flags == RTF_UP);
332 Packet *buffered_pkt;
333 while ((buffered_pkt = rqueue.deque(rt0->rt_dst))) {
334 if (rt0 && (rt0->rt_flags == RTF_UP)) {
335 assert(rt0->rt_hops != INFINITY2);
336 forward(rt0, buffered_pkt, NO_DELAY);
337 }
338 }
339 }
340 // End for putting reverse route in rt table
341
342
343 /*

```

```
344 * We have taken care of the reverse route stuff.
345 * Now see whether we can send a route reply.
346 */
347 rt = rtable.rt_lookup(rq->rq_dst);
348
349 // First check if I am the destination ..
350 if(rq->rq_dst == index) {
351
352 #ifdef DEBUG
353 fprintf(stderr, "%d - %s: destination sending reply\n", index, __FUNCTION__);
354 #endif // DEBUG
355
356 // Just to be safe, I use the max. Somebody may have
357 // incremented the dst seqno.
358 seqno = max(seqno, rq->rq_dst_seqno)+1;
359 if (seqno%2) seqno++;
360
361 destination_ip = (int)index;
362 sendReply(rq->rq_src, // IP Destination
363 1, // Hop Count
364 index, // Dest IP Address
365 seqno, // Dest Sequence Num
366 MY_ROUTE_TIMEOUT, // Lifetime
367 rq->rq_timestamp); // timestamp
368
369 Packet::free(p);
370 }
371
372 // I am not the destination, but I may have a fresh enough route.
373 else if (rt && (rt->rt_hops != INFINITY2) &&
374 (rt->rt_seqno >= rq->rq_dst_seqno)) {
375
376 //assert (rt->rt_flags == RTF_UP);
377 assert(rq->rq_dst == rt->rt_dst);
378 //assert ((rt->rt_seqno%2) == 0); // is the seqno even?
379 sendReply(rq->rq_src,
380 rt->rt_hops + 1,
381 rq->rq_dst,
382 rt->rt_seqno,
383 (u_int32_t) (rt->rt_expire - CURRENT_TIME),
384 //
385 rt->rt_expire - CURRENT_TIME,
386 rq->rq_timestamp);
387 // Insert nexthops to RREQ source and RREQ destination in the
388 // precursor lists of destination and source respectively
389 rt->pc_insert(rt0->rt_nexthop); // nexthop to RREQ source
390 rt0->pc_insert(rt->rt_nexthop); // nexthop to RREQ destination
391
392 #ifdef RREQ_GRAT_RREP
393 sendReply(rq->rq_dst,
394 rq->rq_hop_count,
395 rq->rq_src,
396 rq->rq_src_seqno,
397 (u_int32_t) (rt->rt_expire - CURRENT_TIME),
398 //
399 rt->rt_expire - CURRENT_TIME,
400 rq->rq_timestamp);
401 #endif
402
403 // TODO: send grat RREP to dst if G flag set in RREQ using rq->rq_src_seqno,
404 // rq->rq_hop_count
405
406 // DONE: Included gratuitous replies to be sent as per IETF aadv draft
407 // specification. As of now, G flag has not been dynamically used and is always set or
408 // reset in aadv-packet.h --- Anant Utgikar, 09/16/02.
409
410 Packet::free(p);
411 }
412
413 /*
414 * Can't reply. So forward the Route Request
415 */
416 else {
```

```
412 ih->saddr() = index;
413 ih->daddr() = IP_BROADCAST;
414 rq->rq_hop_count += 1;
415 // Maximum sequence number seen en route
416 if (rt) rq->rq_dst_segno = max(rt->rt_segno, rq->rq_dst_segno);
417 forward((aodv_rt_entry*) 0, p, DELAY);
418 }
419 }
420
421
422 void
423 AODV::recvReply(Packet *p) {
424 struct hdr_cmn *ch = HDR_CMN(p); // MCMCI: Get the common header of this packet
425 struct hdr_ip *ih = HDR_IP(p);
426 struct hdr_aodv_reply *rp = HDR_AODV_REPLY(p);
427 aodv_rt_entry *rt;
428 char suppress_reply = 0;
429 double delay = 0.0;
430
431 #ifdef DEBUG
432 fprintf(stderr, "%d - %s: received a REPLY\n", index, __FUNCTION__);
433 #endif // DEBUG
434
435
436 if (nIfaces) {
437 // MCMCI: When the node receives a "Hello" (RREP) packet from a neighbour,
438 // update its NeighbourTable with the fixed channel of that neighbour
439 neighbour_table[(int)(rp->rp_sender_node_ip)] = rp->rp_fixed_channel_used;
440
441 // MCMCI: Update the node's ChannelUsageList using the NeighbourTable of its
442 // neighbour. Doing so ensures ChannelUsageList will contain two-hop
443 // channel usage information.
444 for (int i = 0; i < MAX_NT_CUL_ENTRIES; i++) {
445 if (i == (int)index) {
446 continue; // MCMCI: No need to deal with the fixed channel used
447 // by itself again; avoid repetition
448 }
449 if (rp->rp_neighbour_table[i] != -1) {
450 //if ((rp->rp_neighbour_table[i] != -1) &&
451 !(channel_usage_list[rp->rp_neighbour_table[i]])) {
452 channel_usage_list[rp->rp_neighbour_table[i]]++;
453 }
454 }
455 #ifdef MCMCI_DEBUG
456 printf("n%d\trecvReply - rp->rp_sender_node_ip: %d\n", (int)index,
457 rp->rp_sender_node_ip);
458 printf("n%d\trecvReply - rp->rp_fixed_channel_used: %d\n", (int)index,
459 rp->rp_fixed_channel_used);
460 rt_print_NT_CUL();
461 #endif
462 } // if (nIfaces)
463
464 /*
465 * Got a reply. So reset the "soft state" maintained for
466 * route requests in the request table. We don't really have
467 * have a separate request table. It is just a part of the
468 * routing table itself.
469 */
470 // Note that rp_dst is the dest of the data packets, not the
471 // the dest of the reply, which is the src of the data packets.
472 rt = rtable.rt_lookup(rp->rp_dst);
473
474 /*
475 * If I don't have a rt entry to this host... adding
476 */
477 if (rt == 0) {
478 rt = rtable.rt_add(rp->rp_dst);
479
480 // MCMCI: Get the interface index
```

```

480 if (nIfaces) {
481 #ifdef MCMI_DEBUG
482 printf("n%d\trecvReply - ch->iface() (rt = 0): %d\n", (int)index, ch->iface());
483 printf("n%d\trecvReply - ((Mac *)ifqueuelist[0]->target()->addr() (rt = 0):
484 %d\n", (int)index, ((Mac *)ifqueuelist[0]->target()->addr());
485 #endif
486 rt->rt_interface = ch->iface()-((Mac *)ifqueuelist[0]->target()->addr());
487 } else {
488 rt->rt_interface = -1;
489 #ifdef MCMI_DEBUG
490 printf("n%d\trecvReply - rt->rt_interface (rt = 0): %d\n\n", (int)index,
491 rt->rt_interface);
492 #endif
493 }
494 /*
495 * Add a forward route table entry... here I am following
496 * Perkins-Royer AODV paper almost literally - SRD 5/99
497 */
498
499 if ((rt->rt_seqno < rp->rp_dst_seqno) || // newer route
500 ((rt->rt_seqno == rp->rp_dst_seqno) &&
501 (rt->rt_hops > rp->rp_hop_count))) { // shorter or better route
502
503 // Update the rt entry
504 rt_update(rt, rp->rp_dst_seqno, rp->rp_hop_count,
505 rp->rp_src, CURRENT_TIME + rp->rp_lifetime);
506
507
508 // MCMI: Get the interface index
509 if (nIfaces) {
510 #ifdef MCMI_DEBUG
511 printf("n%d\trecvReply - ch->iface(): %d\n", (int)index, ch->iface());
512 printf("n%d\trecvReply - ((Mac *)ifqueuelist[0]->target()->addr(): %d\n",
513 (int)index, ((Mac *)ifqueuelist[0]->target()->addr());
514 #endif
515 rt->rt_interface = ch->iface()-((Mac *)ifqueuelist[0]->target()->addr());
516 } else {
517 rt->rt_interface = -1;
518 #ifdef MCMI_DEBUG
519 printf("n%d\trecvReply - rt->rt_interface: %d\n\n", (int)index, rt->rt_interface);
520 #endif
521
522
523 // reset the soft state
524 rt->rt_req_cnt = 0;
525 rt->rt_req_timeout = 0.0;
526 rt->rt_req_last_ttl = rp->rp_hop_count;
527
528 if (ih->daddr() == index) { // If I am the original source
529 // Update the route discovery latency statistics
530 // rp->rp_timestamp is the time of request origination
531
532 rt->rt_disc_latency[(unsigned char)rt->hist_idx] = (CURRENT_TIME -
rp->rp_timestamp)
533
534 / (double) rp->rp_hop_count;
535 // increment idx for next time
536 rt->hist_idx = (rt->hist_idx + 1) % MAX_HISTORY;
537 }
538 /*
539 * Send all packets queued in the sendbuffer destined for
540 * this destination.
541 * XXX - observe the "second" use of p.
542 */
543 Packet *buf_pkt;
544 while((buf_pkt = rqueue.deque(rt->rt_dst))) {
545 if(rt->rt_hops != INFINITY2) {
546 assert (rt->rt_flags == RTF_UP);

```

```

547 // Delay them a little to help ARP. Otherwise ARP
548 // may drop packets. -SRD 5/23/99
549 forward(rt, buf_pkt, delay);
550 delay += ARP_DELAY;
551 }
552 }
553 }
554 else {
555 suppress_reply = 1;
556 }
557
558 /*
559 * If reply is for me, discard it.
560 */
561 if(ih->daddr() == index || suppress_reply) {
562 Packet::free(p);
563 }
564
565 /*
566 * Otherwise, forward the Route Reply.
567 */
568 else {
569 // Find the rt entry
570 aadv_rt_entry *rt0 = rtable.rt_lookup(ih->daddr());
571 // If the rt is up, forward
572 if(rt0 && (rt0->rt_hops != INFINITY2)) {
573 assert (rt0->rt_flags == RTF_UP);
574 rp->rp_hop_count += 1;
575 rp->rp_src = index;
576 forward(rt0, p, NO_DELAY);
577 // Insert the nexthop towards the RREQ source to
578 // the precursor list of the RREQ destination
579 rt->pc_insert(rt0->rt_nexthop); // nexthop to RREQ source
580 }
581 else {
582 // I don't know how to forward .. drop the reply.
583 #ifdef DEBUG
584 fprintf(stderr, "%s: dropping Route Reply\n", __FUNCTION__);
585 #endif // DEBUG
586 drop(p, DROP_RTR_NO_ROUTE);
587 }
588 }
589 }
590
591 void
592 AODV::recvError(Packet *p) {
593 struct hdr_ip *ih = HDR_IP(p);
594 struct hdr_aadv_error *re = HDR_AADV_ERROR(p);
595 aadv_rt_entry *rt;
596 u_int8_t i;
597 Packet *rerr = Packet::alloc();
598 struct hdr_aadv_error *nre = HDR_AADV_ERROR(rerr);
599
600 nre->DestCount = 0;
601
602 for (i=0; i<re->DestCount; i++) {
603 // For each unreachable destination
604 rt = rtable.rt_lookup(re->unreachable_dst[i]);
605 if (rt && (rt->rt_hops != INFINITY2) &&
606 (rt->rt_nexthop == ih->saddr()) &&
607 (rt->rt_seqno <= re->unreachable_dst_seqno[i])) {
608 assert(rt->rt_flags == RTF_UP);
609 assert((rt->rt_seqno%2) == 0); // is the seqno even?
610 #ifdef DEBUG
611 fprintf(stderr, "%s(%f): n%d\t(n%d\t%u\t%d)\t(n%d\t%u\t%d)\n",
612 __FUNCTION__, CURRENT_TIME,
613 index, rt->rt_dst, rt->rt_seqno, rt->rt_nexthop,
614 re->unreachable_dst[i], re->unreachable_dst_seqno[i],
615 ih->saddr());
616 #endif // DEBUG

```

```

617 rt->rt_seqno = re->unreachable_dst_seqno[i];
618 rt_down(rt);
619
620 // Not sure whether this is the right thing to do
621 Packet *pkt;
622
623
624 // MCMI: Get the corresponding interface queue
625 for (int i = 0; i < nIfaces; i++)
626 if (rt->rt_interface == i) {
627 ifqueue = ifqueueelist[i];
628 break;
629 }
630
631
632 while((pkt = ifqueue->filter(ih->saddr())) {
633 drop(pkt, DROP_RTR_MAC_CALLBACK);
634 }
635
636 // if precursor list non-empty add to RERR and delete the precursor list
637 if (!rt->pc_empty()) {
638 nre->unreachable_dst[nre->DestCount] = rt->rt_dst;
639 nre->unreachable_dst_seqno[nre->DestCount] = rt->rt_seqno;
640 nre->DestCount += 1;
641 rt->pc_delete();
642 }
643 }
644 }
645
646 if (nre->DestCount > 0) {
647 #ifdef DEBUG
648 fprintf(stderr, "%s(%f): n%d\t sending RERR...\n", __FUNCTION__, CURRENT_TIME,
649 index);
650 #endif // DEBUG
651 sendError(rerr);
652 }
653 else {
654 Packet::free(rerr);
655 }
656
657 Packet::free(p);
658 }
659
660 /*
661 Packet Transmission Routines
662 */
663
664 void
665 AODV::forward(aadv_rt_entry *rt, Packet *p, double delay) {
666 struct hdr_cmh *ch = HDR_CMH(p);
667 struct hdr_ip *ih = HDR_IP(p);
668 struct hdr_aadv_request *rq = HDR_AADV_REQUEST(p); // MCMI: RREQ packet
669 struct hdr_aadv_reply *rp = HDR_AADV_REPLY(p); // MCMI: RREP packet
670
671 if (ih->tll_ == 0) {
672 #ifdef DEBUG
673 fprintf(stderr, "%s: calling drop()\n", __PRETTY_FUNCTION__);
674 #endif // DEBUG
675
676 drop(p, DROP_RTR_TTL);
677 return;
678 }
679
680 if (ch->ptype() != PT_AADV && ch->direction() == hdr_cmh::UP
681 && ((u_int32_t)ih->daddr() == IP_BROADCAST)
682 || (ih->daddr() == here_.addr_)) {
683 dmux_->recv(p, 0);
684 return;
685 }
686 }

```

```

687 if (rt) {
688 assert(rt->rt_flags == RTF_UP);
689 rt->rt_expire = CURRENT_TIME + ACTIVE_ROUTE_TIMEOUT;
690 ch->next_hop_ = rt->rtnexthop;
691 ch->addr_type() = NS_AF_INET;
692 ch->direction() = hdr_cmn::DOWN; //important: change the packet's direction
693 }
694 else { // if it is a broadcast packet
695 // assert(ch->ptype() == PT_AODV); // maybe a diff pkt type like gaf
696 assert(ih->daddr() == (nsaddr_t) IP_BROADCAST);
697 ch->addr_type() = NS_AF_NONE;
698 ch->direction() = hdr_cmn::DOWN; //important: change the packet's direction
699 }
700
701 if (nIfaces) {
702 #ifdef MCMC_DEBUG
703 printf("n%d\tforward - rq_type: 0x%x\n", (int)index, rq->rq_type);
704 printf("n%d\tforward - rp_type: 0x%x\n", (int)index, rp->rp_type);
705 #endif
706 // MCMC: For the case of forwarding a RREQ...
707 if (rq->rq_type == AODVTYPE_RREQ) {
708 // MCMC: Add the fixed channel used by this node to the RREQ packet
709 rq->rq_fixed_channel_used = neighbour_table[(int)index];
710 rq->rq_sender_node_ip = (int)index;
711 rq->rq_neighbour_table = &neighbour_table[0];
712 #ifdef MCMC_DEBUG
713 printf("n%d\tforward - rq_fixed_channel_used: %d\n", (int)index,
714 rq->rq_fixed_channel_used);
715 printf("n%d\tforward - rq_sender_node_ip: %d\n", (int)index,
716 rq->rq_sender_node_ip);
717 #endif
718 }
719 // MCMC: For the case of forwarding a RREP...
720 if (rp->rp_type == AODVTYPE_RREP || rp->rp_type == AODVTYPE_HELLO) {
721 // MCMC: Add the fixed channel used by this node to the RREP packet
722 rp->rp_fixed_channel_used = neighbour_table[(int)index];
723 rp->rp_sender_node_ip = (int)index;
724 rp->rp_neighbour_table = &neighbour_table[0];
725 #ifdef MCMC_DEBUG
726 printf("n%d\tforward - rp_fixed_channel_used: %d\n", (int)index,
727 rp->rp_fixed_channel_used);
728 printf("n%d\tforward - rp_sender_node_ip: %d\n", (int)index,
729 rp->rp_sender_node_ip);
730 #endif
731 }
732 // MCMC: For the case of forwarding data packets...
733 if (rq->rq_type == 0 && rp->rp_type == 0) {
734 // MCMC: Consult the node's ChannelUsageList, find the largest usage
735 int channel_largest_usage = 0;
736 int channel_lowest_usage = 0;
737 int found = 0;
738 static int skip = 0;
739 for (int i = 0; i < nIfaces; i++) {
740 if (channel_usage_list[i] == 0) {
741 channel_lowest_usage = i;
742 continue;
743 }
744 if (channel_usage_list[i] >= channel_usage_list[channel_largest_usage]) {
745 channel_largest_usage = i;
746 found = 1;
747 }
748 if (channel_usage_list[i] <= channel_usage_list[channel_lowest_usage]) {
749 channel_lowest_usage = i;
750 }
751 }
752 #ifdef MCMC_DEBUG
753 printf("n%d\tforward - channel_largest_usage: %d, #nodes: %d\n", (int)index,
754 channel_largest_usage, channel_usage_list[channel_largest_usage]);

```

```

753 printf("n%d\tforward - channel_lowest_usage: %d, #nodes: %d\n", (int)index,
channel_lowest_usage, channel_usage_list[channel_lowest_usage]);
754 #endif
755 if (found && !skip) {
756 // MCMI: If the node's fixed channel is the one with the largest usage...
757 if (channel_usage_list[neighbour_table[(int)index]] ==
channel_usage_list[channel_largest_usage]) {
758 int switch_probability = rand() % 101; // switching probability
759 #ifdef MCMI_DEBUG
760 printf("n%d\tforward - switch_probability: %d\n", (int)index,
switch_probability);
761 #endif
762 if (switch_probability <= 40) {
763 // MCMI: Reverse the ChannelUsageList about the fixed channel
previously used by the node
764 int fixed_channel = ((Mac
*)ifqueueelist[fixed_interface]->target()->netif()->channel()->index
());
765 channel_usage_list[fixed_channel]--;
766
767 // MCMI: Change the node's fixed channel to a less used channel
768 fixed_channel = channel_lowest_usage; // new fixed channel
769 channel_usage_list[fixed_channel]++;
770 neighbour_table[(int)index] = fixed_channel;
771 for (int i = 0; i < nIfaces; i++) {
772 if ((Mac
*)ifqueueelist[i]->target()->netif()->channel()->index() ==
fixed_channel) {
773 fixed_interface = i;
774 break;
775 }
776 }
777 #ifdef MCMI_DEBUG
778 printf("n%d\tforward - \tnew fixed_channel: %d\n", (int)index,
fixed_channel);
779 printf("n%d\tforward - \tnew fixed_interface: %d\n", (int)index,
fixed_interface);
780 rt_print_NT_CUL();
781 printf("n%d\tforward - \tready to sendHello\n", (int)index);
782 #endif
783 // MCMI: Transmit a Hello packet informing neighbours of its new
fixed channel
784 sendHello();
785 } // if (switch_probability <= 40)
786 } // if (channel_usage_list[neighbour_table[(int)index]] >= 2)
787 } // if (found)
788 skip = 1;
789
790 // MCMI: Look up the fixed channel of the next node in its
791 // NeighbourTable and assign it as the switchable channel
792 int switchable_channel = neighbour_table[(int)(rt->rt_nexthop)];
793 #ifdef MCMI_DEBUG
794 printf("n%d\tforward - back to forward\n", (int)index);
795 printf("n%d\tforward - rt->rt_nexthop: %d\n", (int)index,
(int)(rt->rt_nexthop));
796 printf("n%d\tforward - switchable_channel: %d\n", (int)index,
switchable_channel);
797 #endif
798 // MCMI: Given the switchable channel, find the corresponding interface
799 for (int i = 0; i < MAX_IF; i++) {
800 if (((Mac *)ifqueueelist[i]->target()->netif()->channel()->index() ==
switchable_channel) {
801 switchable_interface = i;
802 break;
803 }
804 }
805 #ifdef MCMI_DEBUG
806 printf("n%d\tforward - switchable_interface: %d\n", (int)index,
switchable_interface);
807 #endif
808 } // if (rq->rq_type == 0 && rp->rp_type == 0)

```



```

809 } // if (nIfaces)
810
811
812 if (ih->daddr() == (nsaddr_t) IP_BROADCAST) {
813 // If it is a broadcast packet
814 assert(rt == 0);
815 /*
816 * Jitter the sending of broadcast packets by 10ms
817 */
818 if (nIfaces) { // MCMI: Send a broadcast packet
819 for (int i = 0; i < nIfaces; i++) {
820 Packet *p_copy = p->copy();
821 Scheduler::instance().schedule(targetlist[i], p_copy, 0.01 *
Random::uniform());
822 #ifdef MCMI_DEBUG
823 printf("n%d\tforward - broadcast\n\n", (int)index);
824 #endif
825 }
826 Packet::free((Packet *)p);
827 }
828 else {
829 Scheduler::instance().schedule(target_, p, 0.01 * Random::uniform());
830 }
831 }
832 else { // Not a broadcast packet
833 if (delay > 0.0) {
834 if (nIfaces) {
835 if (rq->rq_type == 0 && rp->rp_type == 0) { // data packet
836 #ifdef MCMI_DEBUG
837 printf("n%d\tforward - switchable_interface (delay > 0): %d\n\n",
(int)index, switchable_interface);
838 #endif
839 Scheduler::instance().schedule(targetlist[switchable_interface], p,
delay);
840 } else {
841 // MCMI
842 #ifdef MCMI_DEBUG
843 printf("n%d\tforward - rt->rt_interface (delay > 0): %d\n\n",
(int)index, rt->rt_interface);
844 #endif
845 Scheduler::instance().schedule(targetlist[rt->rt_interface], p, delay);
846 }
847 } else {
848 Scheduler::instance().schedule(target_, p, delay);
849 }
850 }
851 else {
852 // Not a broadcast packet, no delay, send immediately
853 if (nIfaces) { // MCMI: Send a unicast packet
854 if (rq->rq_type == 0 && rp->rp_type == 0) { // data packet
855 #ifdef MCMI_DEBUG
856 printf("n%d\tforward - switchable_interface (else): %d\n\n", (int)index,
switchable_interface);
857 #endif
858 Scheduler::instance().schedule(targetlist[switchable_interface], p, 0);
859 } else {
860 #ifdef MCMI_DEBUG
861 printf("n%d\tforward - rt->rt_interface (else): %d\n\n", (int)index,
rt->rt_interface);
862 #endif
863 Scheduler::instance().schedule(targetlist[rt->rt_interface], p, 0);
864 }
865 } else {
866 Scheduler::instance().schedule(target_, p, 0);
867 }
868 }
869 }
870 }
871
872
873 void

```

```
874 AADV::sendRequest(nsaddr_t dst) {
875 // Allocate a RREQ packet
876 Packet *p = Packet::alloc();
877 struct hdr_cmn *ch = HDR_CMN(p);
878 struct hdr_ip *ih = HDR_IP(p);
879 struct hdr_aadv_request *rq = HDR_AADV_REQUEST(p);
880 aadv_rt_entry *rt = rtable.rt_lookup(dst);
881
882 assert(rt);
883
884 /*
885 * Rate limit sending of Route Requests. We are very conservative
886 * about sending out route requests.
887 */
888
889 if (rt->rt_flags == RTF_UP) {
890 assert(rt->rt_hops != INFINITY2);
891 Packet::free((Packet *)p);
892 return;
893 }
894
895 if (rt->rt_req_timeout > CURRENT_TIME) {
896 Packet::free((Packet *)p);
897 return;
898 }
899
900 // rt_req_cnt is the no. of times we did network-wide broadcast
901 // RREQ_RETRIES is the maximum number we will allow before
902 // going to a long timeout.
903
904 if (rt->rt_req_cnt > RREQ_RETRIES) {
905 rt->rt_req_timeout = CURRENT_TIME + MAX_RREQ_TIMEOUT;
906 rt->rt_req_cnt = 0;
907 Packet *buf_pkt;
908 while ((buf_pkt = rqueue.deque(rt->rt_dst))) {
909 drop(buf_pkt, DROP_RTR_NO_ROUTE);
910 }
911 Packet::free((Packet *)p);
912 return;
913 }
914
915 #ifdef DEBUG
916 fprintf(stderr, "(%2d) - %2d sending Route Request, dst: %d\n",
917 ++route_request, index, rt->rt_dst);
918 #endif // DEBUG
919
920 // Determine the TTL to be used this time.
921 // Dynamic TTL evaluation - SRD
922
923 rt->rt_req_last_ttl = max(rt->rt_req_last_ttl, rt->rt_last_hop_count);
924
925 if (0 == rt->rt_req_last_ttl) {
926 // first time query broadcast
927 ih->ttl_ = TTL_START;
928 }
929 else {
930 // Expanding ring search.
931 if (rt->rt_req_last_ttl < TTL_THRESHOLD)
932 ih->ttl_ = rt->rt_req_last_ttl + TTL_INCREMENT;
933 else {
934 // network-wide broadcast
935 ih->ttl_ = NETWORK_DIAMETER;
936 rt->rt_req_cnt += 1;
937 }
938 }
939
940 // remember the TTL used for the next time
941 rt->rt_req_last_ttl = ih->ttl_;
942
943 // PerHopTime is the roundtrip time per hop for route requests.
944 // The factor 2.0 is just to be safe .. SRD 5/22/99
```

```

945 // Also note that we are making timeouts to be larger if we have
946 // done network wide broadcast before.
947
948 rt->rt_req_timeout = 2.0 * (double) ih->tTL_ * PerHopTime(rt);
949 if (rt->rt_req_cnt > 0)
950 rt->rt_req_timeout *= rt->rt_req_cnt;
951 rt->rt_req_timeout += CURRENT_TIME;
952
953 // Don't let the timeout to be too large, however .. SRD 6/8/99
954 if (rt->rt_req_timeout > CURRENT_TIME + MAX_RREQ_TIMEOUT)
955 rt->rt_req_timeout = CURRENT_TIME + MAX_RREQ_TIMEOUT;
956 rt->rt_expire = 0;
957
958 #ifdef DEBUG
959 fprintf(stderr, "(%2d) - %2d sending Route Request, dst: %d, tout %f ms\n",
960 ++route_request,
961 index, rt->rt_dst,
962 rt->rt_req_timeout - CURRENT_TIME);
963 #endif // DEBUG
964
965
966 // Fill out the RREQ packet
967 // ch->uid() = 0;
968 ch->ptype() = PT_AODV;
969 ch->size() = IP_HDR_LEN + rq->size();
970 ch->iface() = -2;
971 ch->error() = 0;
972 ch->addr_type() = NS_AF_NONE;
973 ch->prev_hop_ = index; // AODV hack
974
975 ih->saddr() = index;
976 ih->daddr() = IP_BROADCAST;
977 ih->sport() = RT_PORT;
978 ih->dport() = RT_PORT;
979
980 // Fill up some more fields.
981 rq->rq_type = AODVTYPE_RREQ;
982 rq->rq_hop_count = 1;
983 rq->rq_bcst_id = bid++;
984 rq->rq_dst = dst;
985 rq->rq_dst_seqno = (rt ? rt->rt_seqno : 0);
986 rq->rq_src = index;
987 seqno += 2;
988 assert ((seqno%2) == 0);
989 rq->rq_src_seqno = seqno;
990 rq->rq_timestamp = CURRENT_TIME;
991
992
993 // MCMI: Add the fixed channel used by this node to the packet
994 rq->rq_fixed_channel_used = neighbour_table[(int)index];
995 rq->rq_sender_node_ip = (int)index;
996 rq->rq_neighbour_table = &neighbour_table[0];
997 #ifdef MCMI_DEBUG
998 printf("%d\tsendRequest - rq_fixed_channel_used: %d\n", (int)index,
999 rq->rq_fixed_channel_used);
1000 printf("%d\tsendRequest - rq_sender_node_ip: %d\n", (int)index, rq->rq_sender_node_ip);
1001 #endif
1002
1003 // MCMI: Send a broadcast packet
1004 if (nIfaces) {
1005 for (int i = 0; i < nIfaces; i++) {
1006 Packet *p_copy = p->copy();
1007 Scheduler::instance().schedule(targetlist[i], p_copy, 0.);
1008 #ifdef MCMI_DEBUG
1009 printf("%d\tsendRequest - broadcast\n\n", (int)index);
1010 #endif
1011 }
1012 Packet::free((Packet *)p);
1013 }
1014 else {
1015 Scheduler::instance().schedule(target_, p, 0.);

```

```
1015 }
1016 }
1017
1018
1019 void
1020 AODV::sendReply(nsaddr_t ipdst, u_int32_t hop_count, nsaddr_t rpdst,
1021 u_int32_t rpseq, u_int32_t lifetime, double timestamp) {
1022 Packet *p = Packet::alloc();
1023 struct hdr_cmn *ch = HDR_CMN(p);
1024 struct hdr_ip *ih = HDR_IP(p);
1025 struct hdr_aodv_reply *rp = HDR_AODV_REPLY(p);
1026 aodv_rt_entry *rt = rtable.rt_lookup(ipdst);
1027
1028 #ifdef DEBUG
1029 fprintf(stderr, "sending Reply from %d at %.2f\n", index,
1030 Scheduler::instance().clock());
1031 #endif // DEBUG
1032 assert(rt);
1033
1034 rp->rp_type = AODVTYPE_RREP;
1035 //rp->rp_flags = 0x00;
1036 rp->rp_hop_count = hop_count;
1037 rp->rp_dst = rpdst;
1038 rp->rp_dst_seqno = rpseq;
1039 rp->rp_src = index;
1040 rp->rp_lifetime = lifetime;
1041 rp->rp_timestamp = timestamp;
1042
1043 // ch->uid() = 0;
1044 ch->ptype() = PT_AODV;
1045 ch->size() = IP_HDR_LEN + rp->size();
1046 ch->iface() = -2;
1047 ch->error() = 0;
1048 ch->addr_type() = NS_AF_INET;
1049 ch->next_hop_ = rt->rtnexthop;
1050 ch->prev_hop_ = index; // AODV hack
1051 ch->direction() = hdr_cmn::DOWN;
1052
1053 ih->saddr() = index;
1054 ih->daddr() = ipdst;
1055 ih->sport() = RT_PORT;
1056 ih->dport() = RT_PORT;
1057 ih->ttl_ = NETWORK_DIAMETER;
1058
1059 // MCMI: Add the fixed channel used by this node to the packet
1060 rp->rp_fixed_channel_used = neighbour_table[(int)index];
1061 rp->rp_sender_node_ip = (int)index;
1062 rp->rp_neighbour_table = &neighbour_table[0];
1063 #ifdef MCMI_DEBUG
1064 printf("n%d\tsendReply - rp_fixed_channel_used: %d\n", (int)index,
1065 rp->rp_fixed_channel_used);
1066 printf("n%d\tsendReply - rp_sender_node_ip: %d\n", (int)index, rp->rp_sender_node_ip);
1067 #endif
1068
1069 // MCMI: Send a unicast packet
1070 if (nIfaces) {
1071 #ifdef MCMI_DEBUG
1072 printf("n%d\tsendReply - rt->rt_interface (unicast): %d\n\n", (int)index,
1073 rt->rt_interface);
1074 #endif
1075 Scheduler::instance().schedule(targetlist[rt->rt_interface], p, 0.);
1076 }
1077 else {
1078 Scheduler::instance().schedule(target_, p, 0.);
1079 }
1080 }
1081
1082 void
1083 AODV::sendError(Packet *p, bool jitter) {
```

```
1083 struct hdr_cmn *ch = HDR_CMN(p);
1084 struct hdr_ip *ih = HDR_IP(p);
1085 struct hdr_aadv_error *re = HDR_AADV_ERROR(p);
1086
1087 #ifdef ERROR
1088 fprintf(stderr, "sending Error from %d at %.2f\n", index,
1089 Scheduler::instance().clock());
1089 #endif // DEBUG
1090
1091 re->re_type = AADVTYPE_RERR;
1092 //re->reserved[0] = 0x00; re->reserved[1] = 0x00;
1093 // DestCount and list of unreachable destinations are already filled
1094
1095 // ch->uid() = 0;
1096 ch->ptype() = PT_AADV;
1097 ch->size() = IP_HDR_LEN + re->size();
1098 ch->iface() = -2;
1099 ch->error() = 0;
1100 ch->addr_type() = NS_AF_NONE;
1101 ch->next_hop_ = 0;
1102 ch->prev_hop_ = index; // AADV hack
1103 ch->direction() = hdr_cmn::DOWN; // important: change the packet's direction
1104
1105 ih->saddr() = index;
1106 ih->daddr() = IP_BROADCAST;
1107 ih->sport() = RT_PORT;
1108 ih->dport() = RT_PORT;
1109 ih->tttl_ = 1;
1110
1111 // Do we need any jitter? Yes
1112 if (jitter) {
1113 // MCFM: Send a broadcast packet
1114 if (nIfaces) {
1115 for (int i = 0; i < nIfaces; i++) {
1116 Packet *p_copy = p->copy();
1117 Scheduler::instance().schedule(targetlist[i], p_copy,
1118 0.01*Random::uniform());
1118 #ifdef MCFM_DEBUG
1119 printf("n%d\tsendMessage - broadcast (jitter)\n\n", (int)index);
1120 #endif
1121 }
1122 Packet::free((Packet *)p);
1123 }
1124 else {
1125 Scheduler::instance().schedule(target_, p, 0.01*Random::uniform());
1126 }
1127 }
1128 else {
1129 // MCFM: Send a broadcast packet
1130 if (nIfaces) {
1131 for (int i = 0; i < nIfaces; i++) {
1132 Packet *p_copy = p->copy();
1133 Scheduler::instance().schedule(targetlist[i], p_copy, 0.0);
1134 #ifdef MCFM_DEBUG
1135 printf("n%d\tsendMessage - broadcast (else)\n\n", (int)index);
1136 #endif
1137 }
1138 Packet::free((Packet *)p);
1139 }
1140 else {
1141 Scheduler::instance().schedule(target_, p, 0.0);
1142 }
1143 }
1144 }
1145
1146
1147 /*
1148 Neighbor Management Functions
1149 */
1150
1151 void
```

```
1152 AODV::sendHello() {
1153 Packet *p = Packet::alloc();
1154 struct hdr_cmn *ch = HDR_CMN(p);
1155 struct hdr_ip *ih = HDR_IP(p);
1156 struct hdr_aodv_reply *rh = HDR_AODV_REPLY(p);
1157
1158 #ifdef DEBUG
1159 fprintf(stderr, "sending Hello from %d at %.2f\n", index,
1160 Scheduler::instance().clock());
1161 #endif // DEBUG
1162
1163 rh->rp_type = AODVTYPE_HELLO;
1164 //rh->rp_flags = 0x00;
1165 rh->rp_hop_count = 1;
1166 rh->rp_dst = index;
1167 rh->rp_dst_seqno = seqno;
1168 rh->rp_lifetime = (1 + ALLOWED_HELLO_LOSS) * HELLO_INTERVAL;
1169
1170 // ch->uid() = 0;
1171 ch->ptype() = PT_AODV;
1172 ch->size() = IP_HDR_LEN + rh->size();
1173 ch->iface() = -2;
1174 ch->error() = 0;
1175 ch->addr_type() = NS_AF_NONE;
1176 ch->prev_hop_ = index; // AODV hack
1177
1178 ih->saddr() = index;
1179 ih->daddr() = IP_BROADCAST;
1180 ih->sport() = RT_PORT;
1181 ih->dport() = RT_PORT;
1182 ih->ttl_ = 1;
1183
1184 // MCMI: Send a broadcast packet
1185 if (nIfaces) {
1186 // MCMI: In order for the node to inform its neighbours of its new fixed
1187 // channel, add the fixed channel used by this node to the packet
1188 rh->rp_fixed_channel_used = neighbour_table[(int)index];
1189 rh->rp_sender_node_ip = (int)index;
1190 rh->rp_neighbour_table = &neighbour_table[0];
1191 #ifdef MCMI_DEBUG
1192 printf("n%d\tsendHello - rh_fixed_channel_used: %d\n", (int)index,
1193 rh->rp_fixed_channel_used);
1194 printf("n%d\tsendHello - rh_sender_node_ip: %d\n", (int)index,
1195 rh->rp_sender_node_ip);
1196 #endif
1197
1198 // MCMI: Send a broadcast packet
1199 for (int i = 0; i < nIfaces; i++) {
1200 Packet *p_copy = p->copy();
1201 Scheduler::instance().schedule(targetlist[i], p_copy, 0.0);
1202 #ifdef MCMI_DEBUG
1203 printf("n%d\tsendHello - broadcast\n\n", (int)index);
1204 #endif
1205 }
1206 Packet::free((Packet *)p);
1207 }
1208 else {
1209 Scheduler::instance().schedule(target_, p, 0.0);
1210 }
1211 }
1212
1213 void
1214 AODV::recvHello(Packet *p) {
1215 //struct hdr_ip *ih = HDR_IP(p);
1216 struct hdr_aodv_reply *rp = HDR_AODV_REPLY(p);
1217 AODV_Neighbor *nb;
1218
1219 /*
1220 * Whenever a node receives a Hello message from a neighbor, the node
```

```
1220 * SHOULD make sure that it has an active route to the neighbor, and
1221 * create one if necessary. If a route already exists, then the
1222 * Lifetime for the route should be increased, if necessary, to be at
1223 * least ALLOWED_HELLO_LOSS * HELLO_INTERVAL.
1224 */
1225 nb = nb_lookup(rp->rp_dst);
1226 if(nb == 0) {
1227 nb_insert(rp->rp_dst);
1228 }
1229 else {
1230 nb->nb_expire = CURRENT_TIME + (1.5 * ALLOWED_HELLO_LOSS * HELLO_INTERVAL);
1231 }
1232
1233
1234 if (nIfaces) {
1235 // MCFI: Update the node's ChannelUsageList using the NeighbourTable of
1236 // its neighbour. Doing so ensures ChannelUsageList will contain
1237 // two-hop channel usage information.
1238 for (int i = 0; i < MAX_NT_CUL_ENTRIES; i++) {
1239 // MCFI: No need to deal with the fixed channel used by itself again
1240 if ((rp->rp_neighbour_table[i] == -1) || i == (int)index
1241 || (neighbour_table[i] == rp->rp_neighbour_table[i])) {
1242 continue; // avoid repetition
1243
1244 // MCFI: If the fixed channel used by a neighbour is different
1245 // from the incoming update, the ChannelUsageList needs to
1246 // be reflected appropriately to handle the update
1247 } else if (neighbour_table[i] != rp->rp_neighbour_table[i]) {
1248 if (channel_usage_list[neighbour_table[i]] > 0) {
1249 channel_usage_list[neighbour_table[i]]--;
1250 }
1251 channel_usage_list[rp->rp_neighbour_table[i]]++;
1252
1253 // MCFI: Normally, just update the ChannelUsageList accordingly
1254 } else {
1255 channel_usage_list[rp->rp_neighbour_table[i]]++;
1256 }
1257 }
1258 // MCFI: When the node receives a "Hello" packet from a neighbour,
1259 // update its NeighbourTable with the fixed channel of that neighbour.
1260 neighbour_table[(int)(rp->rp_sender_node_ip)] = rp->rp_fixed_channel_used;
1261
1262 #ifdef MCFI_DEBUG
1263 printf("n%d\trecvHello - rp->rp_sender_node_ip: %d\n", (int)index,
1264 rp->rp_sender_node_ip);
1265 printf("n%d\trecvHello - rp->rp_fixed_channel_used: %d\n", (int)index,
1266 rp->rp_fixed_channel_used);
1267 rt_print_NT_CUL();
1268 #endif
1269 } // if (nIfaces)
1270
1271 Packet::free(p);
1272 ...
1273
```

```
1 /*
2 #=====
3 # ENSC 835: High-Performance Networks
4 # Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Networks
5 #
6 # Student: Chih-Hao Howard Chang
7 # 20007-2192
8 # howardc@sfu.ca
9 #
10 # Description: Modified AODV routing agent with the implementation of the
11 # multi-interface approach developed by [1] and the interface
12 # switching protocol proposed by [2], which support multi-channel
13 # multi-interface ad-hoc wireless network simulation
14 #
15 # File: aodv.h
16 #
17 # [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2,"
18 # University of Cantabria, Jan. 2007 (User Guide).
19 #
20 # [2] P. Kyasanur and N. H. Vaidya, "Routing and Link-layer Protocols for
21 # Multi-Channel Multi-Interface Ad Hoc Wireless Networks," SIGMOBILE Mobile
22 # Computing and Communications Review, vol. 10, no. 1, pp. 31-43, Jan. 2006.
23 #=====
24 */
25
26
27 ...
28
29 // MCFM: Define various constants
30 #define MAX_IF 12 // maximum number of interfaces per node
31 #define MAX_NT_CUL_ENTRIES 20 // maximum number of entries in NeighbourTable
32 // and ChannelUsageList
33
34 ...
35
36 /*
37 The Routing Agent
38 */
39 class AODV: public Agent {
40
41 /*
42 * make some friends first
43 */
44 friend class aodv_rt_entry;
45 friend class BroadcastTimer;
46 friend class HelloTimer;
47 friend class NeighborTimer;
48 friend class RouteCacheTimer;
49 friend class LocalRepairTimer;
50
51 public:
52 AODV(nsaddr_t id);
53 void recv(Packet *p, Handler *);
54
55
56 // MCFM: Keeps track of the number of interfaces that the agent is managing
57 int nIfaces;
58
59
60 protected:
61 int command(int, const char *const *);
62 int initialized() { return 1 && target_; }
63
64 /*
65 * Route Table Management
66 */
67 void rt_resolve(Packet *p);
68 void rt_update(aodv_rt_entry *rt, u_int32_t seqnum,
69 u_int16_t metric, nsaddr_t nexthop,
70 double expire_time);
71 void rt_down(aodv_rt_entry *rt);
```



```
72
73
74 // MCMI: Print the routing table of this node
75 void rt_print(nsaddr_t node_id);
76
77 // MCMI: Print NeighbourTable and ChannelUsageList of this node
78 void rt_print_NT_CUL();
79
80
81 void local_rt_repair(aadv_rt_entry *rt, Packet *p);
82 public:
83 void rt_ll_failed(Packet *p);
84 void handle_link_failure(nsaddr_t id);
85 protected:
86 void rt_purge(void);
87
88 void enqueue(aadv_rt_entry *rt, Packet *p);
89 Packet* deque(aadv_rt_entry *rt);
90
91 /*
92 * Neighbor Management
93 */
94 void nb_insert(nsaddr_t id);
95 AODV_Neighbor* nb_lookup(nsaddr_t id);
96 void nb_delete(nsaddr_t id);
97 void nb_purge(void);
98
99 /*
100 * Broadcast ID Management
101 */
102
103 void id_insert(nsaddr_t id, u_int32_t bid);
104 bool id_lookup(nsaddr_t id, u_int32_t bid);
105 void id_purge(void);
106
107 /*
108 * Packet TX Routines
109 */
110 void forward(aadv_rt_entry *rt, Packet *p, double delay);
111 void sendHello(void);
112 void sendRequest(nsaddr_t dst);
113
114 void sendReply(nsaddr_t ipdst, u_int32_t hop_count,
115 nsaddr_t rpdst, u_int32_t rpseq,
116 u_int32_t lifetime, double timestamp);
117 void sendError(Packet *p, bool jitter = true);
118
119 /*
120 * Packet RX Routines
121 */
122 void recvAODV(Packet *p);
123 void recvHello(Packet *p);
124 void recvRequest(Packet *p);
125 void recvReply(Packet *p);
126 void recvError(Packet *p);
127
128 /*
129 * History management
130 */
131
132 double PerHopTime(aadv_rt_entry *rt);
133
134
135 nsaddr_t index; // IP Address of this node
136 u_int32_t seqno; // Sequence Number
137 int bid; // Broadcast ID
138
139 aadv_rtable rthead; // routing table
140 aadv_ncache nbhead; // Neighbor Cache
141 aadv_bcache bihead; // Broadcast ID Cache
142
```

```
143 /*
144 * Timers
145 */
146 BroadcastTimer btimer;
147 HelloTimer htimer;
148 NeighborTimer ntimer;
149 RouteCacheTimer rtimer;
150 LocalRepairTimer lrtimer;
151
152 /*
153 * Routing Table
154 */
155 aadv_rtable rtable;
156 /*
157 * A "drop-front" queue used by the routing layer to buffer
158 * packets to which it does not have a route.
159 */
160 aadv_rqueue rqueue;
161
162 /*
163 * A mechanism for logging the contents of the routing
164 * table.
165 */
166 Trace *logtarget;
167
168 /*
169 * A pointer to the network interface queue that sits
170 * between the "classifier" and the "link layer".
171 */
172 PriQueue *ifqueue;
173
174 // MCMI: The routing agent needs to decide which one of the interfaces
175 // the outgoing packets should be routed to. Since there are now
176 // multiple interfaces. The originally used ifqueue and target
177 // pointers need to be modified as follows:
178
179 // MCMI: Store the LL modules for all the interfaces a node has
180 NsObject *targetlist[MAX_IF];
181
182 // MCMI: Keep the corresponding queues of all the interfaces
183 PriQueue *ifqueuelist[MAX_IF];
184
185 // MCMI: Contain the fixed channels used by the node's neighbours
186 int neighbour_table[MAX_NT_CUL_ENTRIES]; // index by node IP
187
188 // MCMI: Contain the number of nodes using each channel as their fixed channel
189 int channel_usage_list[MAX_NT_CUL_ENTRIES]; // index by channel ID
190
191 // MCMI: Fixed and switchable interfaces used by this node
192 int fixed_interface;
193 int switchable_interface;
194
195 /*
196 * Logging stuff
197 */
198 void log_link_del(nsaddr_t dst);
199 void log_link_broke(Packet *p);
200 void log_link_kept(nsaddr_t dst);
201
202 /* for passing packets up to agents */
203 PortClassifier *dmux_;
204 };
205
206 #endif /* __aadv_h__ */
207
208
209
210
```

```
1 /*
2 #=====
3 # ENSC 835: High-Performance Networks
4 # Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Networks
5 #
6 # Student: Chih-Hao Howard Chang
7 # 20007-2192
8 # howardc@sfu.ca
9 #
10 # Description: Modified AODV routing agent with the implementation of the
11 # multi-interface approach developed by [1] and the interface
12 # switching protocol proposed by [2], which support multi-channel
13 # multi-interface ad-hoc wireless network simulation
14 #
15 # File: aodv_packet.h
16 #
17 # [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2,"
18 # University of Cantabria, Jan. 2007 (User Guide).
19 #
20 # [2] P. Kyasanur and N. H. Vaidya, "Routing and Link-layer Protocols for
21 # Multi-Channel Multi-Interface Ad Hoc Wireless Networks," SIGMOBILE Mobile
22 # Computing and Communications Review, vol. 10, no. 1, pp. 31-43, Jan. 2006.
23 #=====
24 */
25
26
27 ...
28
29 struct hdr_aodv_request {
30 u_int8_t rq_type; // Packet Type
31 u_int8_t reserved[2];
32 u_int8_t rq_hop_count; // Hop Count
33 u_int32_t rq_bcast_id; // Broadcast ID
34
35 nsaddr_t rq_dst; // Destination IP Address
36 u_int32_t rq_dst_seqno; // Destination Sequence Number
37 nsaddr_t rq_src; // Source IP Address
38 u_int32_t rq_src_seqno; // Source Sequence Number
39
40 double rq_timestamp; // when REQUEST sent;
41 // used to compute route discovery latency
42
43
44 // MCFM
45 int rq_fixed_channel_used; // fixed channel used by this node
46 nsaddr_t rq_sender_node_ip; // node IP of the RREQ sender (not
47 int *rq_neighbour_table; // NeighbourTable of this node
48
49
50 // This define turns on gratuitous replies- see aodv.cc for implementation contributed by
51 // Anant Utgikar, 09/16/02.
52 // #define RREQ_GRAT_RREP 0x80
53
54 inline int size() {
55 int sz = 0;
56 /*
57 sz = sizeof(u_int8_t) // rq_type
58 + 2*sizeof(u_int8_t) // reserved
59 + sizeof(u_int8_t) // rq_hop_count
60 + sizeof(double) // rq_timestamp
61 + sizeof(u_int32_t) // rq_bcast_id
62 + sizeof(nsaddr_t) // rq_dst
63 + sizeof(u_int32_t) // rq_dst_seqno
64 + sizeof(nsaddr_t) // rq_src
65 + sizeof(u_int32_t); // rq_src_seqno
66 */
67 sz = 7*sizeof(u_int32_t);
68 assert (sz >= 0);
69 return sz;
70 }
```

```
71 };
72
73 struct hdr_aodv_reply {
74 u_int8_t rp_type; // Packet Type
75 u_int8_t reserved[2];
76 u_int8_t rp_hop_count; // Hop Count
77 nsaddr_t rp_dst; // Destination IP Address
78 u_int32_t rp_dst_seqno; // Destination Sequence Number
79 nsaddr_t rp_src; // Source IP Address
80 double rp_lifetime; // Lifetime
81
82 double rp_timestamp; // when corresponding REQ sent;
83 // used to compute route discovery latency
84
85
86 // MCMI
87 int rp_fixed_channel_used; // fixed channel used by this node
88 nsaddr_t rp_sender_node_ip; // node IP of the RREP sender (not
89 // originator)
90 int *rp_neighbour_table; // NeighbourTable of this node
91
92 inline int size() {
93 int sz = 0;
94 /*
95 * sz = sizeof(u_int8_t) // rp_type
96 * + 2*sizeof(u_int8_t) // rp_flags + reserved
97 * + sizeof(u_int8_t) // rp_hop_count
98 * + sizeof(double) // rp_timestamp
99 * + sizeof(nsaddr_t) // rp_dst
100 * + sizeof(u_int32_t) // rp_dst_seqno
101 * + sizeof(nsaddr_t) // rp_src
102 * + sizeof(u_int32_t); // rp_lifetime
103 */
104 sz = 6*sizeof(u_int32_t);
105 assert (sz >= 0);
106 return sz;
107 }
108 };
109 };
110 ...
111 ...
112
```

```
1 /*
2 #=====
3 # ENSC 835: High-Performance Networks
4 # Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Networks
5 #
6 # Student: Chih-Hao Howard Chang
7 # 20007-2192
8 # howardc@sfu.ca
9 #
10 # Description: Modified AODV routing agent with the implementation of the
11 # multi-interface approach developed by [1] and the interface
12 # switching protocol proposed by [2], which support multi-channel
13 # multi-interface ad-hoc wireless network simulation
14 #
15 # File: aodv_rtable.h
16 #
17 # [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2,"
18 # University of Cantabria, Jan. 2007 (User Guide).
19 #
20 # [2] P. Kyasanur and N. H. Vaidya, "Routing and Link-layer Protocols for
21 # Multi-Channel Multi-Interface Ad Hoc Wireless Networks," SIGMOBILE Mobile
22 # Computing and Communications Review, vol. 10, no. 1, pp. 31-43, Jan. 2006.
23 #=====
24 */
25
26
27 ...
28
29 /*
30 Route Table Entry
31 */
32
33 class aodv_rt_entry {
34 friend class aodv_rtable;
35 friend class AODV;
36 friend class LocalRepairTimer;
37 public:
38 aodv_rt_entry();
39 ~aodv_rt_entry();
40
41 void nb_insert(nsaddr_t id);
42 AODV_Neighbor* nb_lookup(nsaddr_t id);
43
44 void pc_insert(nsaddr_t id);
45 AODV_Precursor* pc_lookup(nsaddr_t id);
46 void pc_delete(nsaddr_t id);
47 void pc_delete(void);
48 bool pc_empty(void);
49
50 double rt_req_timeout; // when I can send another req
51 u_int8_t rt_req_cnt; // number of route requests
52
53 protected:
54 LIST_ENTRY(aodv_rt_entry) rt_link;
55
56 nsaddr_t rt_dst;
57 u_int32_t rt_seqno;
58 /* u_int8_t rt_interface; */
59 u_int16_t rt_hops; // hop count
60 int rt_last_hop_count; // last valid hop count
61 nsaddr_t rtnexthop; // next hop IP address
62
63
64 // MCFM: For routing the packet via a specific interface to a destination,
65 // unicast transmission of AODV is exercised. Just knowing the next
66 // hop is not enough; the routing agent must also consider which output
67 // interface to be used to reach the next hop.
68 int rt_interface; // interface index for outgoing route packets
69
70
71 /* list of precursors */
```

```
72 aadv_precursors rt_pclist;
73 double rt_expire; // when entry expires
74 u_int8_t rt_flags;
75
76 #define RTF_DOWN 0
77 #define RTF_UP 1
78 #define RTF_IN_REPAIR 2
79
80 /*
81 * Must receive 4 errors within 3 seconds in order to mark
82 * the route down.
83 u_int8_t rt_errors; // error count
84 double rt_error_time;
85 #define MAX_RT_ERROR 4 // errors
86 #define MAX_RT_ERROR_TIME 3 // seconds
87 */
88
89 #define MAX_HISTORY 3
90 double rt_disc_latency[MAX_HISTORY];
91 char hist_indx;
92 int rt_req_last_ttl; // last ttl value used
93 // last few route discovery latencies
94 // double rt_length [MAX_HISTORY];
95 // last few route lengths
96
97 /*
98 * a list of neighbors that are using this route.
99 */
100 aadv_ncache rt_nblast;
101 };
102
103 ...
104
```

```
1 /*
2 #=====
3 # ENSC 835: High-Performance Networks
4 # Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Networks
5 #
6 # Student: Chih-Hao Howard Chang
7 # 20007-2192
8 # howardc@sfu.ca
9 #
10 # Description: Modified MobileNode class with the implementation of the
11 # multi-interface approach developed by [1]
12 #
13 # File: mobilenode.cc
14 #
15 # [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2,"
16 # University of Cantabria, Jan. 2007 (User Guide).
17 #=====
18 */
19
20
21 ...
22
23 // MCMI: Retrieve the location of a node
24 void
25 MobileNode::getLoc(double *x, double *y, double *z) {
26 update_position();
27 *x = X_;
28 *y = Y_;
29 *z = Z_;
30 }
31
```

```
1 /*
2 #=====
3 # ENSC 835: High-Performance Networks
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5 #
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7 # 20007-2192
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9 #
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11 # multi-interface approach developed by [1]
12 #
13 # File: mobilenode.h
14 #
15 # [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2,"
16 # University of Cantabria, Jan. 2007 (User Guide).
17 #=====
18 */
19
20
21 ...
22
23 #define MAX_CHANNELS 12 // MCFI: Maximum number of channels per node
24
25 ...
26
27 class MobileNode : public Node
28 {
29 friend class PositionHandler;
30 public:
31 MobileNode();
32 virtual int command(int argc, const char*const* argv);
33
34 double distance(MobileNode*);
35 double propdelay(MobileNode*);
36 void start(void);
37 void getLoc(double *x, double *y, double *z); // MCFI
38
39 // MCFI: Remove the inline declaration of the getLoc() function. Due to the
40 // above changes on the MobileNode lists, the original declaration has been found
41 // to always
42 // return a zero distance, which leads to wrong packet receptions.
43 // inline void getLoc(double *x, double *y, double *z) {
44 // update_position(); *x = X_; *y = Y_; *z = Z_;
45 // }
46 inline void getVelo(double *dx, double *dy, double *dz) {
47 *dx = dX_ * speed_; *dy = dY_ * speed_; *dz = 0.0;
48 }
49 inline MobileNode* nextnode() { return link_.le_next; }
50 inline int base_stn() { return base_stn_; }
51 inline void set_base_stn(int addr) { base_stn_ = addr; }
52
53 void dump(void);
54
55 inline MobileNode*& next() { return next_; }
56 inline double X() { return X_; }
57 inline double Y() { return Y_; }
58 inline double Z() { return Z_; }
59 inline double speed() { return speed_; }
60 inline double dX() { return dX_; }
61 inline double dY() { return dY_; }
62 inline double dZ() { return dZ_; }
63 inline double destX() { return destX_; }
64 inline double destY() { return destY_; }
65 inline double radius() { return radius_; }
66 inline double getUpdateTime() { return position_update_time_; }
67 //inline double last_routingtime() { return last_rt_time_; }
68
69 void update_position();
70 void log_energy(int);
71 //void logrttime(double);
```



```
71 virtual void idle_energy_patch(float, float);
72
73 /* For list-keeper */
74 // MobileNode* nextX_;
75 // MobileNode* prevX_;
76 // MCFM: We define a new declaration of the MobileNode lists to replace the
77 // existing ones. ns-2 controls each instance of the MobileNode objects
78 // which are associated with a channel by means of a linked-list. Two
79 // lists are managed; one references the previous node, prevX_, while
80 // the other references the next node, nextX_. The original format of
81 // the list is simply a pointer to a node. In order to support multiple
82 // channels, the list is modified to be an array of pointers with the
83 // size of the array being the maximum of number of channels:
84 MobileNode* nextX_[MAX_CHANNELS];
85 MobileNode* prevX_[MAX_CHANNELS];
86
87 protected:
88 /*
89 * Last time the position of this node was updated.
90 */
91 double position_update_time_;
92 double position_update_interval_;
93
94 /*
95 * The following indicate the (x,y,z) position of the node on
96 * the "terrain" of the simulation.
97 */
98 double X_;
99 double Y_;
100 double Z_;
101 double speed_; // meters per second
102
103 /*
104 * The following is a unit vector that specifies the
105 * direction of the mobile node. It is used to update
106 * position
107 */
108 double dX_;
109 double dY_;
110 double dZ_;
111
112 /* where are we going? */
113 double destX_;
114 double destY_;
115
116 /*
117 * for gridkeeper use only
118 */
119 MobileNode* next_;
120 double radius_;
121
122 // Used to generate position updates
123 PositionHandler pos_handle_;
124 Event pos_intr_;
125
126 void log_movement();
127 void random_direction();
128 void random_speed();
129 void random_destination();
130 int set_destination(double x, double y, double speed);
131
132 private:
133 inline int initialized() {
134 return (T_ && log_target_ &&
135 X_ >= T_->lowerX() && X_ <= T_->upperX() &&
136 Y_ >= T_->lowerY() && Y_ <= T_->upperY());
137 }
138 void random_position();
139 void bound_position();
140 int random_motion_; // is mobile
141
```

```
142 /*
143 * A global list of mobile nodes
144 */
145 LIST_ENTRY(MobileNode) link_;
146
147
148 /*
149 * The topography over which the mobile node moves.
150 */
151 Topography *T_;
152 /*
153 * Trace Target
154 */
155 Trace* log_target_;
156 /*
157 * base_stn for mobilenodes communicating with wired nodes
158 */
159 int base_stn_;
160
161
162 //int last_rt_time_;
163 };
164
165 #endif // ns_mobilenode_h
166
```

```

1 /*
2 #=====
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5 #
6 # Student: Chih-Hao Howard Chang
7 # 20007-2192
8 # howardc@sfu.ca
9 #
10 # Description: Modified Channel class with the implementation of the
11 # multi-interface approach developed by [1]
12 #
13 # File: channel.cc
14 #
15 # [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2,"
16 # University of Cantabria, Jan. 2007 (User Guide).
17 #=====
18 */
19
20
21 ...
22
23 // MCMCI: Due to the changes on the MobileNode lists, we modify accessing each
24 // node entry when attaching, removing, and updating a new node to a channel
25 // to refer to the corresponding channel number. This number can be
26 // accessed by this->index(), where this is the current instance of the
27 // channel class object. In other words, whenever nextX_ and prevX_ appear
28 // in channel.cc, they need to be replaced by:
29 // nextX_[this->index()]
30 // prevX_[this->index()]
31
32 ...
33
34 void
35 WirelessChannel::sendUp(Packet* p, Phy *tifp)
36 {
37 Scheduler &s = Scheduler::instance();
38 Phy *rifp = ifhead_.lh_first;
39 Node *tnode = tifp->node();
40 Node *rnode = 0;
41 Packet *newp;
42 double propdelay = 0.0;
43 struct hdr_cmn *hdr = HDR_CMN(p);
44
45 /* list-based improvement */
46 if(highestAntennaZ_ == -1) {
47 calcHighestAntennaZ(tifp);
48 #ifdef DEBUG
49 fprintf(stdout, "channel.cc:sendUp - Calc highestAntennaZ_ and distCST_\n");
50 fprintf(stdout, "highestAntennaZ_ = %0.1f, distCST_ = %0.1f\n", highestAntennaZ_,
51 distCST_);
52 #endif
53 }
54
55 hdr->direction() = hdr_cmn::UP;
56
57 // still keep grid-keeper around ??
58 if (GridKeeper::instance()) {
59 int i;
60 GridKeeper* gk = GridKeeper::instance();
61 int size = gk->size_;
62
63 MobileNode **outlist = new MobileNode *[size];
64
65 int out_index = gk->get_neighbors((MobileNode*)tnode,
66 outlist);
67 for (i=0; i < out_index; i++) {
68
69 newp = p->copy();
70 rnode = outlist[i];
71 propdelay = get_pdelay(tnode, rnode);
72
73 rifp = (rnode->ifhead()).lh_first;
74 for(; rifp; rifp = rifp->nextnode()){
75 if (rifp->channel() == this){
76 s.schedule(rifp, newp, propdelay);
77 }
78 }
79 }
80 }
81 }

```

```

76 break;
77 }
78 }
79 }
80 delete [] outlist;
81
82 } else { // use list-based improvement
83
84 MobileNode *mtnode = (MobileNode *) tnode;
85 MobileNode **affectedNodes;// **aN;
86 int numAffectedNodes = -1, i;
87
88 if(!sorted_){
89 sortLists();
90 }
91
92 affectedNodes = getAffectedNodes(mtnode, distCST_ + /* safety */ 5,
&numAffectedNodes);
93 for (i=0; i < numAffectedNodes; i++) {
94 rnode = affectedNodes[i];
95
96 if(rnode == tnode)
97 continue;
98
99 newp = p->copy();
100
101 propdelay = get_pdelay(tnode, rnode);
102
103 rifp = (rnode->ifhead()).lh_first;
104 for(; rifp; rifp = rifp->nextnode()){
105 // MCMI: Checks which of the interfaces of the destination node
106 // is connected to the same channel
107 if (rifp->channel() == this) {
108 s.schedule(rifp, newp, propdelay);
109 }
110 }
111 }
112 delete [] affectedNodes;
113 }
114 Packet::free(p);
115 }
116
117
118 void
119 WirelessChannel::addNodeToList(MobileNode *mn)
120 {
121 MobileNode *tmp;
122
123 // create list of mobilenodes for this channel
124 if (xListHead_ == NULL) {
125 #ifdef DEBUG
126 fprintf(stderr, "INITIALIZE THE LIST xListHead\n");
127 #endif
128 xListHead_ = mn;
129 xListHead_->nextX_[this->index()] = NULL;
130 xListHead_->prevX_[this->index()] = NULL;
131 } else {
132 for (tmp = xListHead_; tmp->nextX_[this->index()] != NULL;
tmp=tmp->nextX_[this->index()]);
133 tmp->nextX_[this->index()] = mn;
134 mn->prevX_[this->index()] = tmp;
135 mn->nextX_[this->index()] = NULL;
136 }
137 numNodes_++;
138 }
139
140 void
141 WirelessChannel::removeNodeFromList(MobileNode *mn) {
142
143 MobileNode *tmp;
144 // Find node in list
145 for (tmp = xListHead_; tmp->nextX_[this->index()] != NULL;
tmp=tmp->nextX_[this->index()]) {
146 if (tmp == mn) {
147 if (tmp == xListHead_) {
148 xListHead_ = tmp->nextX_[this->index()];

```

```

149 if (tmp->nextX_[this->index()] != NULL)
150 tmp->nextX_[this->index()]->prevX_[this->index()] = NULL;
151 } else if (tmp->nextX_[this->index()] == NULL)
152 tmp->prevX_[this->index()]->nextX_[this->index()] = NULL;
153 else {
154 tmp->prevX_[this->index()]->nextX_[this->index()] =
tmp->nextX_[this->index()];
155 tmp->nextX_[this->index()]->prevX_[this->index()] =
tmp->prevX_[this->index()];
156 }
157 numNodes--;
158 return;
159 }
160 }
161 fprintf(stderr, "Channel: node not found in list\n");
162 }
163
164 void
165 WirelessChannel::sortLists(void) {
166 bool flag = true;
167 MobileNode *m, *q;
168 sorted_ = true;
169
170 #ifdef DEBUG
171 fprintf(stderr, "SORTING LISTS ...");
172 #endif
173
174 /* Buble sort algorithm */
175 // SORT x-list
176 while(flag) {
177 flag = false;
178 m = xListHead_;
179 while (m != NULL){
180 if(m->nextX_[this->index()] != NULL)
181 if (m->X() > m->nextX_[this->index()]->X()){
182 flag = true;
183 //delete_after m;
184 q = m->nextX_[this->index()];
185 m->nextX_[this->index()] = q->nextX_[this->index()];
186 if (q->nextX_[this->index()] != NULL)
187 q->nextX_[this->index()]->prevX_[this->index()] = m;
188
189 //insert_before m;
190 q->nextX_[this->index()] = m;
191 q->prevX_[this->index()] = m->prevX_[this->index()];
192 m->prevX_[this->index()] = q;
193 if (q->prevX_[this->index()] != NULL)
194 q->prevX_[this->index()]->nextX_[this->index()] = q;
195
196 // adjust Head of List
197 if(m == xListHead_)
198 xListHead_ = m->prevX_[this->index()];
199 }
200 m = m -> nextX_[this->index()];
201 }
202 }
203
204 #ifdef DEBUG
205 fprintf(stderr, "DONE!\n");
206 #endif
207 }
208
209 void
210 WirelessChannel::updateNodesList(class MobileNode *mn, double oldX) {
211
212 MobileNode* tmp;
213 double X = mn->X();
214 bool skipX=false;
215
216 if(!sorted_) {
217 sortLists();
218 return;
219 }
220
221 /* xListHead cannot be NULL here (they are created during creation of mobilenode) */
222

```

```

223 /*** DELETE ***/
224 // deleting mn from x-list
225 if(mn->nextX_[this->index()] != NULL) {
226 if(mn->prevX_[this->index()] != NULL){
227 if((mn->nextX_[this->index()]->X() >= X) && (mn->prevX_[this->index()]->X() <=
X)) skipX = true; // the node doesn't change its position in the list
228 else{
229 mn->nextX_[this->index()]->prevX_[this->index()] =
mn->prevX_[this->index()];
230 mn->prevX_[this->index()]->nextX_[this->index()] =
mn->nextX_[this->index()];
231 }
232 }
233 }
234 else{
235 if(mn->nextX_[this->index()]->X() >= X) skipX = true; // skip updating the
first element
236 else{
237 mn->nextX_[this->index()]->prevX_[this->index()] = NULL;
238 xListHead_ = mn->nextX_[this->index()];
239 }
240 }
241 }
242
243 else if(mn->prevX_[this->index()] != NULL){
244 if(mn->prevX_[this->index()]->X() <= X) skipX = true; // skip updating the last
element
245 else mn->prevX_[this->index()]->nextX_[this->index()] = NULL;
246 }
247
248 if ((mn->prevX_[this->index()] == NULL) && (mn->nextX_[this->index()] == NULL)) skipX
= true; //skip updating if only one element in list
249
250 /*** INSERT ***/
251 //inserting mn in x-list
252 if(!skipX){
253 if(X > oldX){
254 for(tmp = mn; tmp->nextX_[this->index()] != NULL &&
tmp->nextX_[this->index()]->X() < X; tmp = tmp->nextX_[this->index()]);
255 //fprintf(stdout, "Scanning the element addr %d X=%0.f\n",
tmp, tmp->X(), tmp->nextX_[this->index()], tmp->nextX_[this->index()]->X());
256 if(tmp->nextX_[this->index()] == NULL) {
257 //fprintf(stdout, "tmp->nextX_[this->index()] is NULL\n");
258 tmp->nextX_[this->index()] = mn;
259 mn->prevX_[this->index()] = tmp;
260 mn->nextX_[this->index()] = NULL;
261 }
262 else{
263 //fprintf(stdout, "tmp->nextX_[this->index()] is not NULL,
tmp->nextX_[this->index()]->X()=%0.f\n", tmp->nextX_[this->index()]->X());
264 mn->prevX_[this->index()] =
tmp->nextX_[this->index()]->prevX_[this->index()];
265 mn->nextX_[this->index()] = tmp->nextX_[this->index()];
266 tmp->nextX_[this->index()]->prevX_[this->index()] = mn;
267 tmp->nextX_[this->index()] = mn;
268 }
269 }
270 else{
271 for(tmp = mn; tmp->prevX_[this->index()] != NULL &&
tmp->prevX_[this->index()]->X() > X; tmp = tmp->prevX_[this->index()]);
272 //fprintf(stdout, "Scanning the element addr %d X=%0.f, prev addr %d
X=%0.f\n", tmp, tmp->X(), tmp->prevX_[this->index()],
tmp->prevX_[this->index()]->X());
273 if(tmp->prevX_[this->index()] == NULL) {
274 //fprintf(stdout, "tmp->prevX_[this->index()] is NULL\n");
275 tmp->prevX_[this->index()] = mn;
276 mn->nextX_[this->index()] = tmp;
277 mn->prevX_[this->index()] = NULL;
278 xListHead_ = mn;
279 }
280 else{
281 //fprintf(stdout, "tmp->prevX_[this->index()] is not NULL,
tmp->prevX_[this->index()]->X()=%0.f\n", tmp->prevX_[this->index()]->X());
282 mn->nextX_[this->index()] =
tmp->prevX_[this->index()]->nextX_[this->index()];
283 mn->prevX_[this->index()] = tmp->prevX_[this->index()];

```

```

284 tmp->prevX_[this->index()->nextX_[this->index()]] = mn;
285 tmp->prevX_[this->index()] = mn;
286 }
287 }
288 }
289 }
290
291
292 MobileNode **
293 WirelessChannel::getAffectedNodes(MobileNode *mn, double radius,
294 int *numAffectedNodes)
295 {
296 double xmin, xmax, ymin, ymax;
297 int n = 0;
298 MobileNode *tmp, **list, **tmpList;
299
300 if (xListHead_ == NULL) {
301 *numAffectedNodes=-1;
302 fprintf(stderr, "xListHead_ is NULL when trying to send!!!\n");
303 return NULL;
304 }
305
306 xmin = mn->X() - radius;
307 xmax = mn->X() + radius;
308 ymin = mn->Y() - radius;
309 ymax = mn->Y() + radius;
310
311 // First allocate as much as possibly needed
312 tmpList = new MobileNode*[numNodes_];
313
314 for(tmp = xListHead_; tmp != NULL; tmp = tmp->nextX_[this->index()]) tmpList[n++] =
315 tmp;
316 for(int i = 0; i < n; ++i)
317 if(tmpList[i]->speed()!=0.0 && (Scheduler::instance().clock() -
318 tmpList[i]->getUpdateTime()) > XLIST_POSITION_UPDATE_INTERVAL)
319 tmpList[i]->update_position();
320 n=0;
321 for(tmp = mn; tmp != NULL && tmp->X() >= xmin; tmp=tmp->prevX_[this->index()])
322 if(tmp->Y() >= ymin && tmp->Y() <= ymax){
323 tmpList[n++] = tmp;
324 }
325 for(tmp = mn->nextX_[this->index()]; tmp != NULL && tmp->X() <= xmax;
326 tmp=tmp->nextX_[this->index()]){
327 if(tmp->Y() >= ymin && tmp->Y() <= ymax){
328 tmpList[n++] = tmp;
329 }
330 }
331 list = new MobileNode*[n];
332 memcpy(list, tmpList, n * sizeof(MobileNode *));
333 delete [] tmpList;
334
335 *numAffectedNodes = n;
336 return list;
337 }
338
339
340
341 /* Only to be used with mobile nodes (WirelessPhy).
342 * NS-2 at its current state support only a flat (non 3D) movement of nodes,
343 * so we assume antenna heights do not change for the duration of
344 * a simulation.
345 * Another assumption - all nodes have the same wireless interface, so that
346 * the maximum distance, corresponding to CST (at max transmission power
347 * level) stays the same for all nodes.
348 */
349 void
350 WirelessChannel::calcHighestAntennaZ(Phy *tifp)
351 {
352 double highestZ = 0;
353 Phy *n;
354
355 for(n = ifhead_.lh_first; n; n = n->nextchnl()) {
356 if(((WirelessPhy *)n)->getAntennaZ() > highestZ)
357 highestZ = ((WirelessPhy *)n)->getAntennaZ();

```

```
358 }
359
360 highestAntennaZ_ = highestZ;
361
362 WirelessPhy *wifp = (WirelessPhy *)tifp;
363 distCST_ = wifp->getDist(wifp->getCSTthresh(), wifp->getPt(), 1.0, 1.0,
364 highestZ, highestZ, wifp->getL(),
365 wifp->getLambda());
366 }
367
368
369 double
370 WirelessChannel::get_pdelay(Node* tnode, Node* rnode)
371 {
372 // Scheduler &s = Scheduler::instance();
373 MobileNode* tmnode = (MobileNode*)tnode;
374 MobileNode* rmnode = (MobileNode*)rnode;
375 double propdelay = 0;
376
377 propdelay = tmnode->propdelay(rmnode);
378
379 assert(propdelay >= 0.0);
380 if (propdelay == 0.0) {
381 /* if the propdelay is 0 b/c two nodes are on top of
382 each other, move them slightly apart -dam 7/28/98 */
383 propdelay = 2 * DBL_EPSILON;
384 //printf ("propdelay 0: %d->%d at %f\n",
385 // tmnode->address(), rmnode->address(), s.clock());
386 }
387 return propdelay;
388 }
389
```



```
1 /*
2 #=====
3 # ENSC 835: High-Performance Networks
4 # Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Networks
5 #
6 # Student: Chih-Hao Howard Chang
7 # 20007-2192
8 # howardc@sfu.ca
9 #
10 # Description: Modified Mac class to support the implementation of the
11 # interface protocol proposed by [2]
12 #
13 # File: mac.h
14 #
15 # [2] P. Kyasanur and N. H. Vaidya, "Routing and Interface Assignment in
16 # Multi-Channel Multi-Interface Wireless Networks," Wireless Communications
17 # and Networking Conference (WCNC) 2005, vol. 4, pp. 2051-2056, Mar. 2005.
18 #=====
19 */
20
21
22 ...
23
24 /* =====
25 MAC data structure
26 =====*/
27
28 class Mac : public BiConnector {
29 public:
30 Mac();
31 virtual void recv(Packet* p, Handler* h);
32 virtual void sendDown(Packet* p);
33 virtual void sendUp(Packet *p);
34
35 virtual void resume(Packet* p = 0);
36 virtual void installTap(Tap *t) { tap_ = t; }
37
38 inline double txtime(int bytes) {
39 return 8. * bytes / bandwidth_;
40 }
41 inline double txtime(Packet* p) {
42 return 8. * (MAC_HDR_LEN + \
43 (HDR_CMN(p)->size()) / bandwidth_;
44 }
45 inline double bandwidth() const { return bandwidth_; }
46
47 inline int addr() { return index_; }
48 inline MacState state() { return state_; }
49 inline MacState state(int m) { return state_ = (MacState) m; }
50
51 // MCMI: So the network interface that this MAC is connecting to can be
52 // accessed externally. Before, it was a protected member variable.
53 inline Phy *netif() { return netif_; }
54
55 //mac methods to set dst, src and hdt_type in pkt hdrs.
56 // note: -1 is the broadcast mac addr.
57 virtual inline int hdr_dst(char* hdr, int dst = -2) {
58 struct hdr_mac *dh = (struct hdr_mac*) hdr;
59 if(dst > -2)
60 dh->macDA_ = dst;
61 return dh->macDA();
62 }
63 virtual inline int hdr_src(char* hdr, int src = -2) {
64 struct hdr_mac *dh = (struct hdr_mac*) hdr;
65 if(src > -2)
66 dh->macSA_ = src;
67 return dh->macSA();
68 }
69 virtual inline int hdr_type(char *hdr, u_int16_t type = 0) {
70 struct hdr_mac *dh = (struct hdr_mac*) hdr;
71 if (type)
72 dh->hdr_type_ = type;
73 return dh->hdr_type();
74 }
75
76 private:
```

```
77 void mac_log(Packet *p) {
78 logtarget_->recv(p, (Handler*) 0);
79 }
80 NSObject* logtarget_;
81
82 protected:
83 int command(int argc, const char*const* argv);
84 virtual int initialized() {
85 return (netif_ && uptarget_ && downtarget_);
86 }
87 int index_; // MAC address
88 double bandwidth_; // channel bitrate
89 double delay_; // MAC overhead
90 int abstract_; // MAC support for abstract LAN
91
92 Phy *netif_; // network interface
93 Tap *tap_; // tap agent
94 LL *ll_; // LL this MAC is connected to
95 Channel *channel_; // channel this MAC is connected to
96
97 Handler* callback_; // callback for end-of-transmission
98 MacHandlerResume hRes_; // resume handler
99 MacHandlerSend hSend_; // handle delay send due to busy channel
100 Event intr_;
101
102 /*
103 * Internal MAC State
104 */
105 MacState state_; // MAC's current state
106 Packet *pktRx_;
107 Packet *pktTx_;
108 };
109
110 #endif
111
```

```
1 /*
2 #=====
3 # ENSC 835: High-Performance Networks
4 # Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Networks
5 #
6 # Student: Chih-Hao Howard Chang
7 # 20007-2192
8 # howardc@sfu.ca
9 #
10 # Description: Modified Mac802_11 class with the implementation of the
11 # multi-interface approach developed by [1]
12 #
13 # File: mac-802_11.cc
14 #
15 # [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2,"
16 # University of Cantabria, Jan. 2007 (User Guide).
17 #=====
18 */
19
20
21 ...
22
23 void
24 Mac802_11::recv(Packet *p, Handler *h)
25 {
26 struct hdr_cmn *hdr = HDR_CMN(p);
27 /*
28 * Sanity Check
29 */
30 assert(initialized());
31
32 /*
33 * Handle outgoing packets.
34 */
35 if(hdr->direction() == hdr_cmn::DOWN) {
36 send(p, h);
37 return;
38 }
39 /*
40 * Handle incoming packets.
41 *
42 * We just received the 1st bit of a packet on the network
43 * interface.
44 *
45 */
46
47 /*
48 * If the interface is currently in transmit mode, then
49 * it probably won't even see this packet. However, the
50 * "air" around me is BUSY so I need to let the packet
51 * proceed. Just set the error flag in the common header
52 * to that the packet gets thrown away.
53 */
54 if(tx_active_ && hdr->error() == 0) {
55 hdr->error() = 1;
56 }
57
58 // MCMI: For correct handling of multiple interfaces by the routing agent,
59 // register the correct MAC receiving interface - which a message was
60 // received through
61 hdr->iface() = addr();
62
63 if(rx_state_ == MAC_IDLE) {
64 setRxState(MAC_RECV);
65 pktRx_ = p;
66 /*
67 * Schedule the reception of this packet, in
68 * txttime seconds.
69 */
70 mhRecv_.start(txttime(p));
71 } else {
72 /*
73 * If the power of the incoming packet is smaller than the
74 * power of the packet currently being received by at least
75 * the capture threshold, then we ignore the new packet.
76 */
77 }
```

```
77 if(pktRx_->txinfo_.RxPr / p->txinfo_.RxPr >= p->txinfo_.CPThresh) {
78 capture(p);
79 } else {
80 collision(p);
81 }
82 }
83 }
84
85 ...
86
```

```
1 #=====
2 # ENSC 835: High-Performance Networks
3 # Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Networks
4 #
5 # Student: Chih-Hao Howard Chang
6 # 20007-2192
7 # howardc@sfu.ca
8 #
9 # Description: Modified ns-lib TCL script with the implementation of the
10 # multi-interface approach developed by [1]
11 #
12 # File: ns-lib.tcl
13 #
14 # [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2,"
15 # University of Cantabria, Jan. 2007 (User Guide).
16 #=====
17
18
19 ...
20
21 # XXX This should be moved into the node initialization procedure instead
22 # of standing here in ns-lib.tcl.
23 Simulator instproc create-wireless-node args {
24 $self instvar routingAgent_ wiredRouting_ propInstance_ llType_ \
25 macType_ ifqType_ ifqlen_ phyType_ chan antType_ \
26 energyModel_ initialEnergy_ txPower_ rxPower_ \
27 idlePower_ sleepPower_ sleepTime_ transitionPower_ transitionTime_ \
28 topoInstance_ level1_ level2_ inerrProc_ outerrProc_ FECProc_ \
29 numifs_
30
31 Simulator set IMEPFlag_ OFF
32
33 # create node instance
34 set node [eval $self create-node-instance $args]
35
36 # basestation address setting
37 if { [info exist wiredRouting_] && $wiredRouting_ == "ON" } {
38 $node base-station [AddrParams addr2id [$node node-addr]]
39 }
40 switch -exact $routingAgent_ {
41 DSDV {
42 set ragent [$self create-dsdv-agent $node]
43 }
44 DSR {
45 $self at 0.0 "$node start-dsr"
46 }
47 AODV {
48 set ragent [$self create-aodv-agent $node]
49 }
50 TORA {
51 Simulator set IMEPFlag_ ON
52 set ragent [$self create-tora-agent $node]
53 }
54 DIFFUSION/RATE {
55 eval $node addr $args
56 set ragent [$self create-diffusion-rate-agent $node]
57 }
58 DIFFUSION/PROB {
59 eval $node addr $args
60 set ragent [$self create-diffusion-probability-agent $node]
61 }
62 Directed_Diffusion {
63 eval $node addr $args
64 set ragent [$self create-core-diffusion-rtg-agent $node]
65 }
66 FLOODING {
67 eval $node addr $args
68 set ragent [$self create-flooding-agent $node]
69 }
70 OMNIMCAST {
71 eval $node addr $args
72 set ragent [$self create-omnimcast-agent $node]
73 }
74 DumbAgent {
75 set ragent [$self create-dumb-agent $node]
76 }
77 }
```

```

77 ManualRtg {
78 set ragent [$self create-manual-rtg-agent $node]
79 }
80 # Manual Routing
81 MANUAL {
82 set ragent [$self create-manual-routing-agent $node]
83 }
84 default {
85 eval $node addr $args
86 puts "Wrong node routing agent!"
87 exit
88 }
89 }
90
91 # errProc_ and FECProc_ are an option unlike other
92 # parameters for node interface
93 if ![info exists inerrProc_] {
94 set inerrProc_ ""
95 }
96 if ![info exists outerrProc_] {
97 set outerrProc_ ""
98 }
99 if ![info exists FECProc_] {
100 set FECProc_ ""
101 }
102
103 # MCMI: Add main node interface
104 if {[info exists numifs_]} {
105 for {set i 0} {$i < $numifs_} {incr i} {
106 # Add one interface per channel
107 #puts "chan $i: $chan($i)"
108 $node add-interface $chan($i) $propInstance_ $llType_ $macType_ \
109 $ifqType_ $ifqlen_ $phyType_ $antType_ $topoInstance_ \
110 $inerrProc_ $outerrProc_ $FECProc_
111 }
112 } else {
113 $node add-interface $chan $propInstance_ $llType_ $macType_ \
114 $ifqType_ $ifqlen_ $phyType_ $antType_ $topoInstance_ \
115 $inerrProc_ $outerrProc_ $FECProc_
116 }
117
118 # Attach agent
119 if {$routingAgent_ != "DSR"} {
120 $node attach $ragent [Node set rtagent_port_]
121 }
122 if {$routingAgent_ == "DIFFUSION/RATE" ||
123 $routingAgent_ == "DIFFUSION/PROB" ||
124 $routingAgent_ == "FLOODING" ||
125 $routingAgent_ == "OMNIMCAST" ||
126 $routingAgent_ == "Directed_Diffusion" } {
127 $ragent port-dmux [$node demux]
128 $node instvar ll_
129 $ragent add-ll $ll_(0)
130 }
131 if { $routingAgent_ == "DumbAgent" } {
132 $ragent port-dmux [$node demux]
133 }
134
135
136 # Bind routing agent and mip agent if existing basestation
137 # address setting
138 if { [info exists wiredRouting_] && $wiredRouting_ == "ON" } {
139 if { $routingAgent_ != "DSR" } {
140 $node mip-call $ragent
141 }
142 }
143 #
144 # This Trace Target is used to log changes in direction
145 # and velocity for the mobile node.
146 #
147 set tracefd [$self get-ns-traceall]
148 if {$tracefd != ""} {
149 $node nodetrace $tracefd
150 $node agenttrace $tracefd
151 }
152 set namtracefd [$self get-nam-traceall]

```

```
153 if {$namtracefd != ""} {
154 $node namattach $namtracefd
155 }
156 if [info exists energyModel_] {
157 if [info exists level1_] {
158 set l1 $level1_
159 } else {
160 set l1 0.5
161 }
162 if [info exists level2_] {
163 set l2 $level2_
164 } else {
165 set l2 0.2
166 }
167 $node addenergymodel [new $energyModel_ $node \
168 $initialEnergy_ $l1 $l2]
169 }
170 if [info exists txPower_] {
171 $node setPt $txPower_
172 }
173 if [info exists rxPower_] {
174 $node setPr $rxPower_
175 }
176 if [info exists idlePower_] {
177 $node setPidle $idlePower_
178 }
179 #
180 if [info exists sleepPower_] {
181 $node setPsleep $sleepPower_
182 }
183 if [info exists sleepTime_] {
184 $node setTSleep $sleepTime_
185 }
186 if [info exists transitionPower_] {
187 $node setPtransition $transitionPower_
188 }
189 if [info exists transitionTime_] {
190 $node setTtransition $transitionTime_
191 }
192 #
193 $node topography $topoInstance_
194
195 return $node
196 }
197
198 ...
199
200 # MCMI: Procedure to change the number of interfaces
201 Simulator instproc change-numifs {newnumifs} {
202 $self instvar numifs_
203 set numifs_ $newnumifs
204 }
205
206 # MCMI: Procedure to add an interface on a node
207 Simulator instproc add-channel {indexch ch} {
208 $self instvar chan
209 set chan($indexch) $ch
210 }
211
212 # MCMI: Procedure to get the number of interfaces
213 Simulator instproc get-numifs { } {
214 $self instvar numifs_
215 if [info exists numifs_] {
216 return $numifs_
217 } else {
218 return ""
219 }
220 }
221
222 # Procedure to add multiple interfaces as an argument to node-config label
223 Simulator instproc ifNum {val} {
224 $self set numifs_ $val
225 }
226
```

```
1 #=====
2 # ENSC 835: High-Performance Networks
3 # Implementation of a Multi-Channel Multi-Interface Ad-Hoc Wireless Networks
4 #
5 # Student: Chih-Hao Howard Chang
6 # 20007-2192
7 # howardc@sfu.ca
8 #
9 # Description: Modified ns-mobilenode TCL library script with the implementation
10 # of the multi-interface approach developed by [1]
11 #
12 # File: ns-lib.tcl
13 #
14 # [1] R. A. Calvo and J. P. Campo, "Adding Multiple Interface Support in NS-2,"
15 # University of Cantabria, Jan. 2007 (User Guide).
16 #=====
17
18
19 ...
20
21 Node/MobileNode instproc reset {} {
22 $self instvar arptable_ nifs_ netif_ mac_ ifq_ ll_ imep_
23 for {set i 0} {$i < $nifs_} {incr i} {
24 $netif_($i) reset
25 $mac_($i) reset
26 $ll_($i) reset
27 $ifq_($i) reset
28 if { [info exists opt(imep)] && $opt(imep) == "ON" } {
29 $imep_($i) reset
30 }
31 # MCMI: Reset the ARP table of a MobileNode object per the number of
32 # interfaces defined.
33 if {$arptable_($i) != ""} {
34 $arptable_($i) reset
35 }
36 }
37 }
38
39 #
40 # Attach an agent to a node. Pick a port and
41 # bind the agent to the port number.
42 # if portnumber is 255, default target is set to the routing agent
43 #
44 Node/MobileNode instproc add-target { agent port } {
45 $self instvar dmux_ imep_ toraDebug_
46
47 set ns [Simulator instance]
48 set newapi [$ns imep-support]
49
50 $agent set sport_ $port
51
52 # MCMI: Get the number of interfaces from the simulator object
53 set numIfsSimulator [$ns get-numifs]
54
55 # special processing for TORA/IMEP node
56 set toraonly [string first "TORA" [$agent info class]]
57 if {$toraonly != -1} {
58 $agent if-queue [$self set ifq_(0)] ;# ifq between LL and MAC
59 #
60 # XXX: The routing protocol and the IMEP agents needs handles
61 # to each other.
62 #
63 $agent imep-agent [$self set imep_(0)]
64 [$self set imep_(0)] rtagent $agent
65 }
66
67 # Special processing for AODV
68 set aodvonly [string first "AODV" [$agent info class]]
69 if {$aodvonly != -1} {
70 $agent if-queue [$self set ifq_(0)] ;# ifq between LL and MAC
71 }
72
73 #<zheng: add>
74 # Special processing for ZBR
75 #set zbronly [string first "ZBR" [$agent info class]]
76 #if {$zbronly != -1} {
```



```

77 # $agent if-queue [$self set ifq_(0)] ;# ifq between LL and MAC
78 #}
79 #</zheng: add>
80
81 if { $port == [Node set rtagent_port_] } {
82 # MCMI: Special processing when multiple interfaces are supported
83 if { $numIfsSimulator != "" } {
84 for {set i 0} {$i < [$self set nifs_]} {incr i} {
85 $agent if-queue $i [$self set ifq_($i)]
86 }
87 }
88
89 # Ad hoc routing agent setup needs special handling
90 $self add-target-rtagent $agent $port
91 return
92 }
93
94 # Attaching a normal agent
95 set namfp [$ns get-nam-traceall]
96 if { [Simulator set AgentTrace_] == "ON" } {
97 #
98 # Send Target
99 #
100 if { $newapi != "" } {
101 set sndT [$self mobility-trace Send "AGT"]
102 } else {
103 set sndT [cmu-trace Send AGT $self]
104 }
105 if { $namfp != "" } {
106 $sndT namattach $namfp
107 }
108 $sndT target [$self entry]
109 $agent target $sndT
110 #
111 # Recv Target
112 #
113 if { $newapi != "" } {
114 set rcvT [$self mobility-trace Recv "AGT"]
115 } else {
116 set rcvT [cmu-trace Recv AGT $self]
117 }
118 if { $namfp != "" } {
119 $rcvT namattach $namfp
120 }
121 $rcvT target $agent
122 $dmux_ install $port $rcvT
123 } else {
124 #
125 # Send Target
126 #
127 $agent target [$self entry]
128 #
129 # Recv Target
130 #
131 $dmux_ install $port $agent
132 }
133 }
134
135 Node/MobileNode instproc add-target-rtagent { agent port } {
136 $self instvar imep_ toraDebug_
137
138 set ns [Simulator instance]
139 set newapi [$ns imep-support]
140 set namfp [$ns get-nam-traceall]
141
142 set dmux_ [$self demux]
143 set classifier_ [$self entry]
144
145 # MCMI: Whether multiple interfaces exist in the simulation
146 set numIfsSimulator [$ns get-numifs]
147
148 # let the routing agent know about the port dmux
149 $agent port-dmux $dmux_
150
151 if { [Simulator set RouterTrace_] == "ON" } {
152 #

```

```

153 # Send Target
154 #
155 if {$newapi != ""} {
156 set sndT [$self mobility-trace Send "RTR"]
157 } else {
158 set sndT [cmu-trace Send "RTR" $self]
159 }
160 if { $namfp != "" } {
161 $sndT namattach $namfp
162 }
163 if { $newapi == "ON" } {
164 $agent target $imep_(0)
165 $imep_(0) sendtarget $sndT
166 # second tracer to see the actual
167 # types of tora packets before imep packs them
168 if { [info exists toraDebug_] && $toraDebug_ == "ON" } {
169 set sndT2 [$self mobility-trace Send "TRP"]
170 $sndT2 target $imep_(0)
171 $agent target $sndT2
172 }
173 # MCMI
174 $sndT target [$self set ll_(0)]
175 } else { ;# no IMEP
176 # MCMI: If the number of interfaces is non-zero, the procedure
177 # associates the routing agent with the corresponding link
178 # layer target entity as many times as the number of
179 # interfaces for the receiving target.
180 if {$numIfsSimulator != ""} {
181 for {set i 0} {$i < [$self set nifs_]} {incr i} {
182 set sndT [cmu-trace Send "RTR" $self]
183 $agent target $i $sndT
184 $sndT target [$self set ll_($i)]
185 }
186 } else {
187 $agent target $sndT
188 $sndT target [$self set ll_(0)]
189 }
190 }
191 }
192 #
193 # Recv Target
194 #
195 if {$newapi != ""} {
196 set rcvT [$self mobility-trace Recv "RTR"]
197 } else {
198 set rcvT [cmu-trace Recv "RTR" $self]
199 }
200 if { $namfp != "" } {
201 $rcvT namattach $namfp
202 }
203 if {$newapi == "ON" } {
204 [$self set ll_(0)] up-target $imep_(0)
205 $classifier_ defaulttarget $agent
206 # need a second tracer to see the actual
207 # types of tora packets after imep unpacks them
208 # no need to support any hier node
209 if { [info exists toraDebug_] && $toraDebug_ == "ON" } {
210 set rcvT2 [$self mobility-trace Recv "TRP"]
211 $rcvT2 target $agent
212 $classifier_ defaulttarget $rcvT2
213 }
214 } else {
215 $rcvT target $agent
216 $classifier_ defaulttarget $rcvT
217 $dmux_ install $port $rcvT
218 }
219 } else {
220 #
221 # Send Target
222 #
223 # if tora is used
224 if { $newapi == "ON" } {
225 $agent target $imep_(0)
226 # second tracer to see the actual
227 # types of tora packets before imep packs them
228 if { [info exists toraDebug_] && $toraDebug_ == "ON" } {

```

```

229 set sndT2 [$self mobility-trace Send "TRP"]
230 $sndT2 target $imep_(0)
231 $agent target $sndT2
232 }
233 $imep_(0) sendtarget [$self set ll_(0)]
234
235 } else { ;# no IMEP
236 # MCFM: If the number of interfaces is non-zero, the procedure
237 # associates the routing agent with the corresponding link
238 # layer target entity as many times as the number of
239 # interfaces for the sending target.
240 if {$numIfsSimulator != ""} {
241 for {set i 0} {$i < [$self set nifs_]} {incr i} {
242 $agent target $i [$self set ll_($i)]
243 }
244 } else {
245 $agent target [$self set ll_(0)]
246 }
247 }
248
249 #
250 # Recv Target
251 #
252 if {$newapi == "ON"} {
253 [$self set ll_(0)] up-target $imep_(0)
254 $classifier_ defaulttarget $agent
255 # need a second tracer to see the actual
256 # types of tora packets after imep unpacks them
257 # no need to support any hier node
258 if {[info exists toraDebug_] && $toraDebug_ == "ON"} {
259 set rcvT2 [$self mobility-trace Recv "TRP"]
260 $rcvT2 target $agent
261 [$self set classifier_] defaulttarget $rcvT2
262 }
263 } else {
264 $classifier_ defaulttarget $agent
265 $dmux_ install $port $agent
266 }
267 }
268 }
269
270 #
271 # The following setups up link layer, mac layer, network interface
272 # and physical layer structures for the mobile node.
273 #
274 Node/MobileNode instproc add-interface { channel pmodel lltype mactype qtype qlen iftype
anttype topo inerrproc outerrproc fecproc } {
275 $self instvar arptable_ nifs_ netif_ mac_ ifq_ ll_ imep_ inerr_ outerr_ fec_
276
277 set ns [Simulator instance]
278 set imepflag [$ns imep-support]
279 set t $nifs_
280 incr nifs_
281
282 set netif_($t) [new $iftype] ;# interface
283 set mac_($t) [new $mactype] ;# mac layer
284 set ifq_($t) [new $qtype] ;# interface queue
285 set ll_($t) [new $lltype] ;# link layer
286 set ant_($t) [new $anttype]
287
288 $ns mac-type $mactype
289 set inerr_($t) ""
290 if {$inerrproc != ""} {
291 set inerr_($t) [$inerrproc]
292 }
293 set outerr_($t) ""
294 if {$outerrproc != ""} {
295 set outerr_($t) [$outerrproc]
296 }
297 set fec_($t) ""
298 if {$fecproc != ""} {
299 set fec_($t) [$fecproc]
300 }
301
302 set namfp [$ns get-nam-traceall]
303 if {$imepflag == "ON"} {

```

```
304 # IMEP layer
305 set imep_($t) [new Agent/IMEP [$self id]]
306 set imep $imep_($t)
307 set drpT [$self mobility-trace Drop "RTR"]
308 if { $namfp != "" } {
309 $drpT namattach $namfp
310 }
311 $imep drop-target $drpT
312 $ns at 0.[$self id] "$imep_($t) start" ;# start beacon timer
313 }
314
315 #
316 # Local Variables
317 #
318 set nullAgent_ [$ns set nullAgent_]
319 set netif $netif_($t)
320 set mac $mac_($t)
321 set ifq $ifq_($t)
322 set ll $ll_($t)
323
324 set inerr $inerr_($t)
325 set outerr $outerr_($t)
326 set fec $fec_($t)
327
328 # MCFI: Create one ARP table per interface. Originally, it creates one ARP
329 # table (for address resolution) per node. The reason for such a
330 # walk-around is that, if a node is using one interface to communicate
331 # with another one, the current design of MobileNode in ns-2 will not
332 # allow the node to use another interface since the request to the ARP
333 # entity will still be serving the previous interface.
334 set arptable_($t) [new ARPTable $self $mac]
335 set arptable $arptable_($t)
336
337 # FOR backward compatibility sake, hack only
338 if {$imepflag != ""} {
339 set drpT [$self mobility-trace Drop "IFQ"]
340 } else {
341 set drpT [cmu-trace Drop "IFQ" $self]
342 }
343 $arptable drop-target $drpT
344 if { $namfp != "" } {
345 $drpT namattach $namfp
346 }
347
348 #
349 # Link Layer
350 #
351 $ll arptable $arptable; # MCFI
352 $ll mac $mac
353 $ll down-target $ifq
354
355 if {$imepflag == "ON" } {
356 $imep recvtarget [$self entry]
357 $imep sendtarget $ll
358 $ll up-target $imep
359 } else {
360 $ll up-target [$self entry]
361 }
362
363 #
364 # Interface Queue
365 #
366 $ifq target $mac
367 $ifq set limit_ $qlen
368 if {$imepflag != ""} {
369 set drpT [$self mobility-trace Drop "IFQ"]
370 } else {
371 set drpT [cmu-trace Drop "IFQ" $self]
372 }
373 $ifq drop-target $drpT
374 if { $namfp != "" } {
375 $drpT namattach $namfp
376 }
377 if {[$ifq info class] == "Queue/XCP" } {
378 $mac set bandwidth_ [$ll set bandwidth_]
379 $mac set delay_ [$ll set delay_]
```

```

380 $ifq set-link-capacity [$mac set bandwidth_]
381 $ifq queue-limit $qlen
382 $ifq link $ll
383 $ifq reset
384
385 }
386
387 #
388 # Mac Layer
389 #
390
391 $mac netif $netif
392 $mac up-target $ll
393
394 if {$outerr == "" && $fec == ""} {
395 $mac down-target $netif
396 } elseif {$outerr != "" && $fec == ""} {
397 $mac down-target $outerr
398 $outerr target $netif
399 } elseif {$outerr == "" && $fec != ""} {
400 $mac down-target $fec
401 $fec down-target $netif
402 } else {
403 $mac down-target $fec
404 $fec down-target $outerr
405 $err target $netif
406 }
407
408 set god_ [God instance]
409 if {$mactype == "Mac/802_11"} {
410 $mac nodes [$god_ num_nodes]
411 }
412 #
413 # Network Interface
414 #
415 #if {$fec == ""} {
416 # $netif up-target $mac
417 #} else {
418 # $netif up-target $fec
419 # $fec up-target $mac
420 #}
421
422 $netif channel $channel
423 if {$inerr == "" && $fec == ""} {
424 $netif up-target $mac
425 } elseif {$inerr != "" && $fec == ""} {
426 $netif up-target $inerr
427 $inerr target $mac
428 } elseif {$err == "" && $fec != ""} {
429 $netif up-target $fec
430 $fec up-target $mac
431 } else {
432 $netif up-target $inerr
433 $inerr target $fec
434 $fec up-target $mac
435 }
436
437 $netif propagation $pmodel ;# Propagation Model
438 $netif node $self ;# Bind node <---> interface
439 $netif antenna $ant_($t)
440 #
441 # Physical Channel
442 #
443 $channel addif $netif
444
445 # List-based improvement
446 # For nodes talking to multiple channels this should
447 # be called multiple times for each channel
448 $channel add-node $self
449
450 # let topo keep handle of channel
451 $topo channel $channel
452 # =====
453
454 if { [Simulator set MacTrace_] == "ON" } {
455 #

```

```

456 # Trace RTS/CTS/ACK Packets
457 #
458 if {$imepflag != ""} {
459 set rcvT [$self mobility-trace Recv "MAC"]
460 } else {
461 set rcvT [cmu-trace Recv "MAC" $self]
462 }
463 $mac log-target $rcvT
464 if { $namfp != "" } {
465 $rcvT namattach $namfp
466 }
467 #
468 # Trace Sent Packets
469 #
470 if {$imepflag != ""} {
471 set sndT [$self mobility-trace Send "MAC"]
472 } else {
473 set sndT [cmu-trace Send "MAC" $self]
474 }
475 $sndT target [$mac down-target]
476 $mac down-target $sndT
477 if { $namfp != "" } {
478 $sndT namattach $namfp
479 }
480 #
481 # Trace Received Packets
482 #
483 if {$imepflag != ""} {
484 set rcvT [$self mobility-trace Recv "MAC"]
485 } else {
486 set rcvT [cmu-trace Recv "MAC" $self]
487 }
488 $rcvT target [$mac up-target]
489 $mac up-target $rcvT
490 if { $namfp != "" } {
491 $rcvT namattach $namfp
492 }
493 #
494 # Trace Dropped Packets
495 #
496 if {$imepflag != ""} {
497 set drpT [$self mobility-trace Drop "MAC"]
498 } else {
499 set drpT [cmu-trace Drop "MAC" $self]
500 }
501 $mac drop-target $drpT
502 if { $namfp != "" } {
503 $drpT namattach $namfp
504 }
505 } else {
506 $mac log-target [$ns set nullAgent_]
507 $mac drop-target [$ns set nullAgent_]
508 }
509
510 # change wrt Mike's code
511 if { [Simulator set EotTrace_] == "ON" } {
512 #
513 # Also trace end of transmission time for packets
514 #
515 if {$imepflag != ""} {
516 set eotT [$self mobility-trace EOT "MAC"]
517 } else {
518 set eotT [cmu-trace EOT "MAC" $self]
519 }
520 $mac eot-target $eotT
521 }
522
523 # =====
524 $self addif $netif
525 }
526 ...
527
528
529
530

```