#### Modeling and Simulating STP vs RSTP on Ring and Mesh Network Topologies

ENSC 835: Communication Networks Spring 2011

**Final Project Presentation** 

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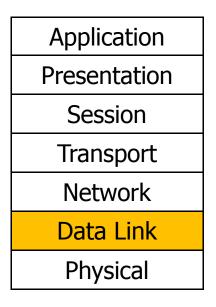
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## Roadmap

- Introduction
- Motivation and Objective
- Overview of related work
- Formation of Spanning Tree
- Bridge Protocol Data Unit
- Simulation and Results in OPNET14.0
- Summary and Conclusions
- Organization and Timelines
- Future Work
- References

#### Introduction

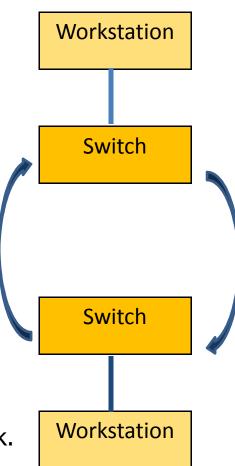
- Spanning Tree Protocol (STP) is a Data Link Layer protocol
- STP is standardized as IEEE 802.1D
- Rapid Spanning Tree Protocol (RSTP) is the refinement of STP
- RSTP is standardized as IEEE 802.1W
- Both, STP and RSTP, ensure a loop-free topology for a bridged Ethernet Local Area Network



OSI Model

#### **Motivation and Objective**

- When two endpoints have more than one Layer-2 path, a switching loop is formed
- Switching loop creates broadcast radiation, i.e. repeatedly rebroadcasting the broadcasted messages
- Objectives:
  - create a switching-loop free network using spanning tree protocols
  - analyze STP vs RSTP performance with a failure and recovery of network link
  - compare tree convergence behaviors by increasing the links and nodes in the network.



#### Overview of related work

Raniwala and T. Chiueh, "Architecture and Algorithms for an IEEE 802.11-Based Multi-Channel Wireless Mesh Network," *Proc. of IEEE Infocom 2005*, vol. 3, 13-17, pp. 2223-2234, March 2005

used traffic load to dynamically assign STP forwarding path

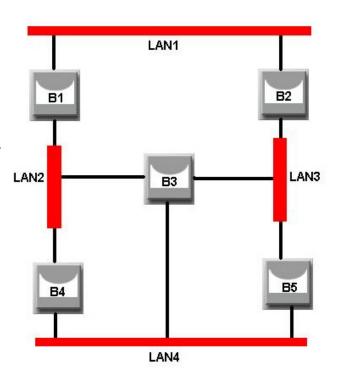
H. S. Chiu, B. Wu, K. L. Yeung, and K.-S. Lui, "Widest Spanning Tree for Multi-Channel Multi-Interface Wireless Mesh Networks," *Proc. Of IEEE WCNC*, pp. 2194-2199, April 2008

solved wide STP problem using mathematical formulation

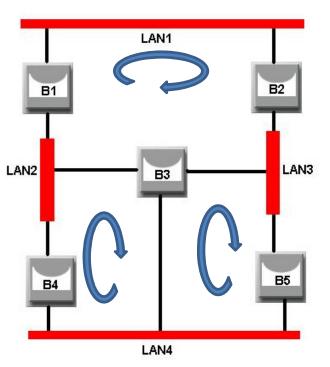
K. Lui, W.C. Lee and K. Nahrstedt, "STAR: A transparent spanning tree bridge protocol with alternate routing," ACM SIGCOMM Computer Communications, Review 32, July 2002

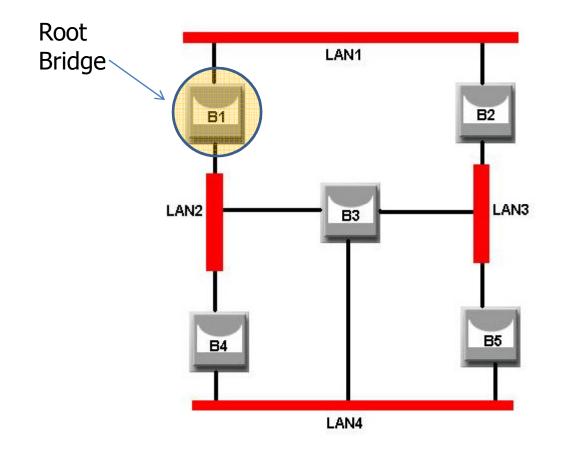
presented shortest forwarding path to improve QoS

Five bridges
B1, B2, B3, B4 and B5
Four LANs
LAN1, LAN2, LAN3 and LAN4
LANs connected to each other through Bridges

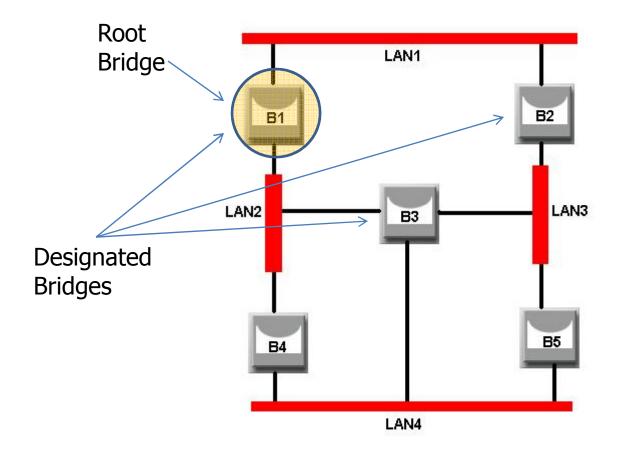


- Five bridges
  - B1, B2, B3, B4 and B5
- Four LANs
  - LAN1, LAN2, LAN3 and LAN4
- LANs connected to each other through Bridges
- Loops are created due to physical connections

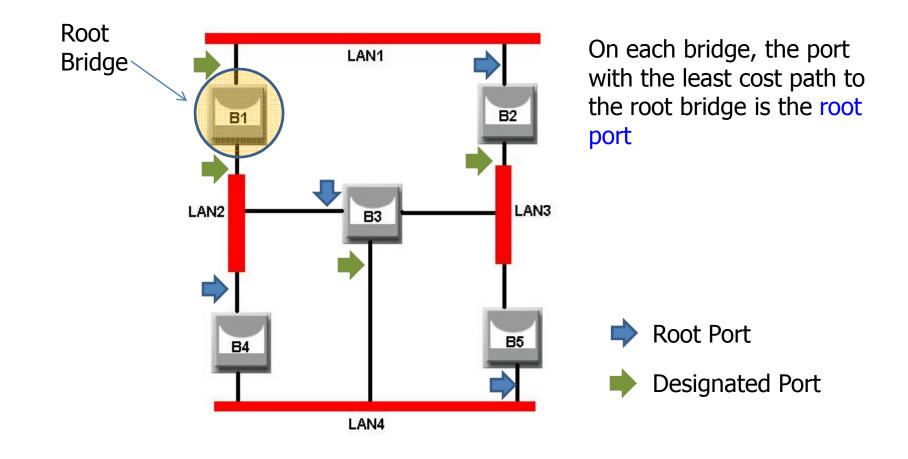




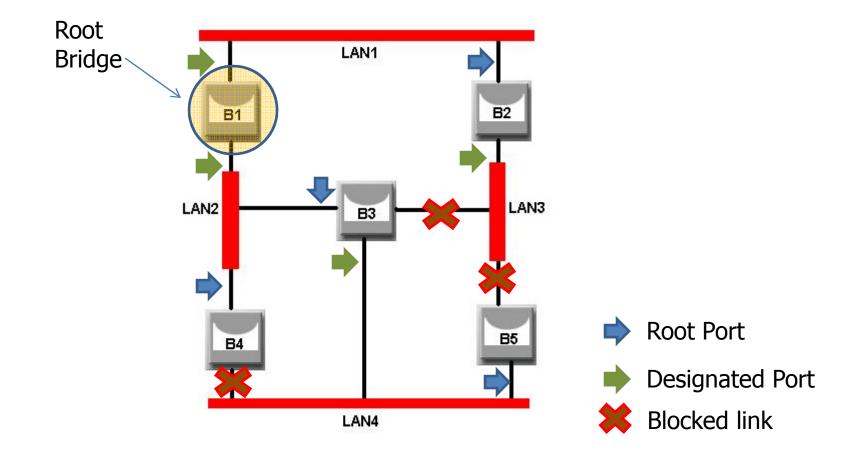
The bridge with the smallest ID is the root bridge



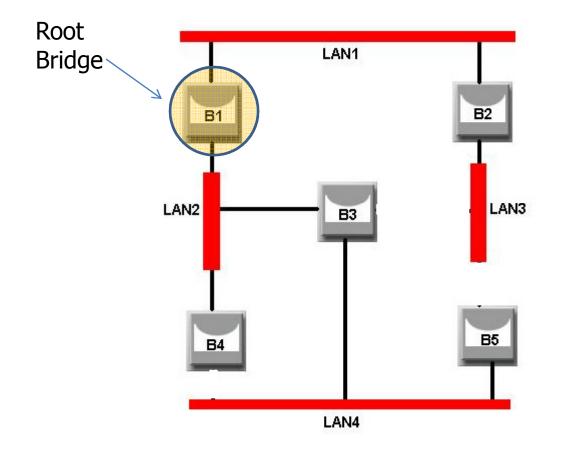
On each LAN, the bridge with the least cost path to the root bridge is the designated bridge











#### Bridge Protocol Data Unit (BPDU)

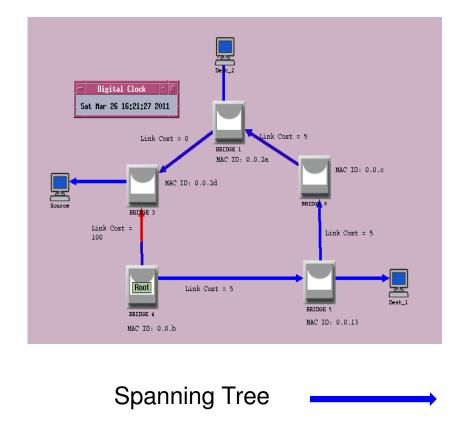
2	1	1	1	8	4	8	2	2	2	2	2
Proto ID	Ver	BPDU Type	Flags	Root ID	Root Path Cost	Bridge ID	Port ID	Msg Age	Max Age	Hello Time	Forward Delay

- BPDU is a 802.1 data frame used to communicate between bridges
- Contains timing parameters,
   MAC Address and priority of the root bridge and designated bridge

Timing Parameter	Description	Simulatio n Value
Hello Time	How often bridge broadcasts BPDU frame to the network	2 seconds
Maximum Age	Time before the received BPDU expires in the bridge	20 seconds
Forwarding Delay	Time the bridge takes to converge from blocking state to forwarding state	15 seconds

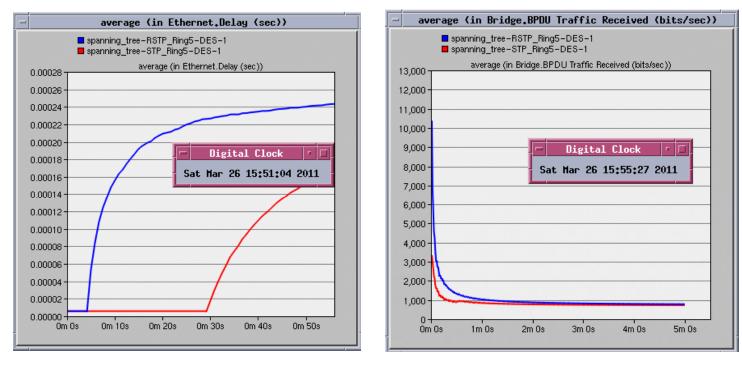
## Simulation – OPNET14.0

- Ring Topology with 5 bridges
- Server connected to Bridge 3
- Workstation connected to Bridge 5
- Bridge 3-4 link virtually blocked by STP formation



**Blocked Link** 

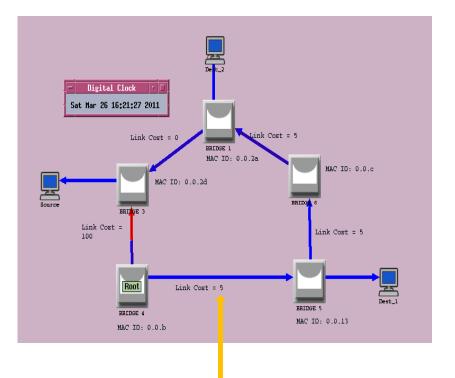
## STP vs RSTP in Ring Topology



- STP formation delay ~ 30s
   RSTP formation delay ~ 6s
- BPDU traffic in RSTP > STP
- Results in quick formation of spanning tree in RSTP

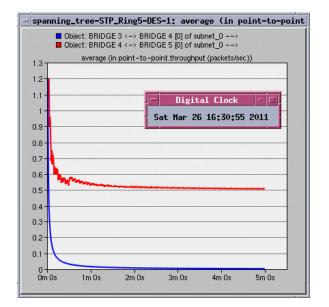
#### Failure/Recovery Analysis in STP

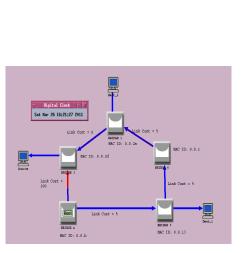
- At 2 minutes, link
   between Bridge 4 and
   Bridge 5 fails
- Blocked link between
   Bridge 3 and Bridge 4
   reactivates
- At 3 minutes, link
   between Bridge 4 and
   Bridge 5 recovers
- Link between Bridge 3 and Bridge 4 gets blocked

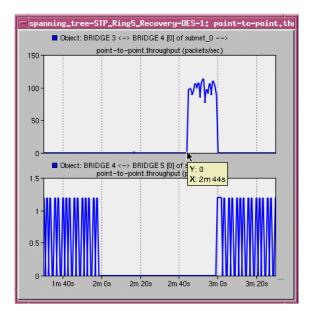


Link fails at 2 min. Link recovers at 3 min.

## Failure/Recovery Analysis in STP



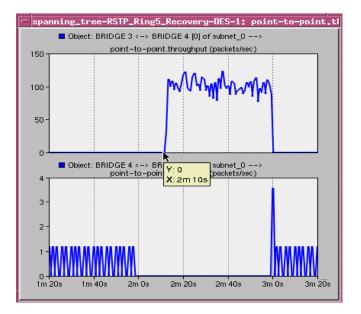




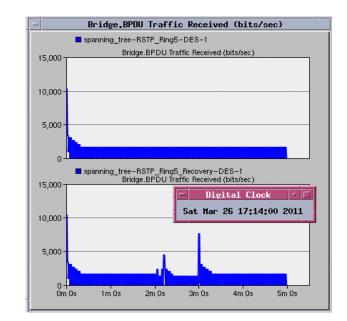
 Without failure/recovery
 Bridge4<->Bridge5 has continuous throughput
 Bridge3<->Bridge4 has zero throughput

- With failure/recovery
- Bridge4<->Bridge5 fails from 2 minutes up to 3 minutes
- Bridge3<->Bridge4 activates after 2 minutes + STP formation time

#### Failure/Recovery Analysis in RSTP



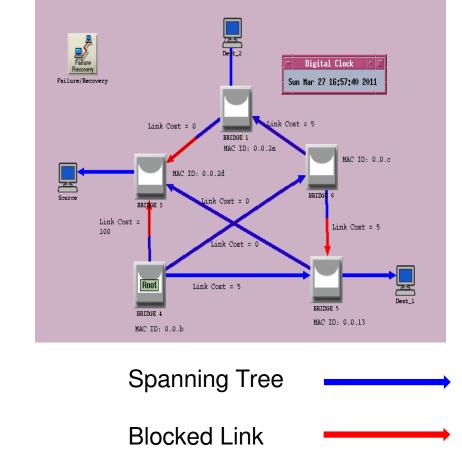
- Bridge4<->Bridge5 fails from 2 minutes up to 3 minutes
- Bridge3<->Bridge4 activates after 2 minutes + RSTP formation time



 BPDU Traffic increases at 2 minutes and 3 minutes

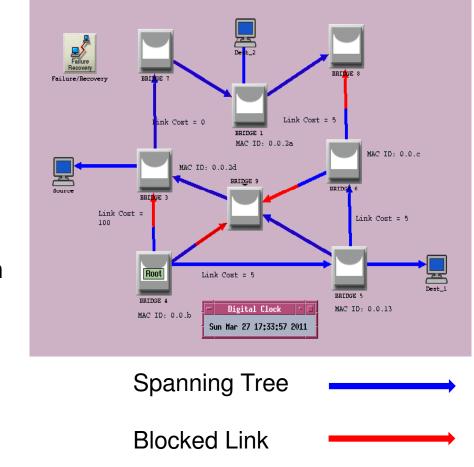
## Simulation – OPNET14.0

- Mesh Topology with 5 bridges
- Server connected to Bridge 3
- Workstation connected to Bridge 5
- Bridge3<->Bridge4,
   Bridge1<->Bridge3 and
   Bridge5<->Bridge6
   blocked by STP formation

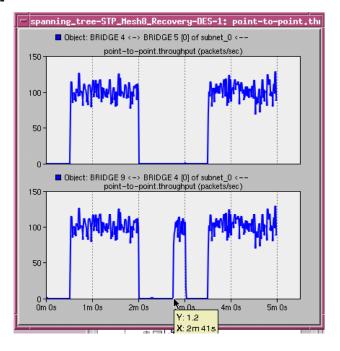


#### Simulation – OPNET14.0

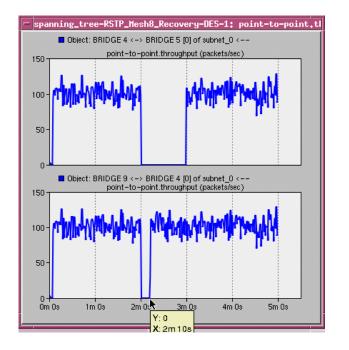
- Mesh Topology with 8 bridges
- Bridge3<->Bridge4,
   Bridge4<->Bridge9,
   Bridge9<->Bridge6 and
   Bridge9<->Bridge8
   blocked by STP formation



#### STP vs RSTP with 8 bridges in Mesh



STP takes 41 seconds to reform a tree



 RSTP takes 10 seconds to reform a tree

#### Conclusion

- STP and RSTP create virtual network and prevents switching loop
- Ring Topology
  - STP takes ~30s to form a spanning tree
  - RSTP takes ~6s to form a spanning tree
- Mesh Topology
  - Increasing links did not change the spanning tree formation time
  - Increasing nodes increased STP formation time by 11s and RSTP formation time by 4s
- RSTP uses more BPDU traffic than STP

# **Organization and Timelines**

Task	Completed by	Completion Time
Understand Spanning Tree protocols	Simran and Manjur	4 weeks
Familiarize with OPNET14.0	Simran and Manjur	12 weeks
Create Ring Model	Simran and Manjur	4 weeks
Create Mesh Model with 5 bridges	Simran and Manjur	1 week
Create Mesh Model with 8 bridges	Simran and Manjur	1 week
Analyze results	Simran and Manjur	2 weeks

# Future Work

- Determine the maximum number of nodes supported by a tree to ensure the BPDU travels to all the nodes before its maximum age expires
- Reduce Spanning tree re-formation time by storing pre-calculated spanning trees with every possible link failure

#### References

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